# Operating Systems Lecture 9: Limited Direct Execution +

Memory Virtualisation

Nipun Batra Aug 23, 2018

#### Administrative

- 1. Next Wednesday answer sheets in lab session
- 2. Projects list would be available on Monday
  - 1. Project 5 -> 8% (3% reduced from homework)
  - 2. More details on Tuesday...

Time multiplexing: Share resource by dividing over time

1. CPU scheduling on single core

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- 2. Think more?!

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- 3. Class room scheduling single class runs in at any given point of time

- 1. CPU scheduling on single core
- 2. Think more?!
- 3. Class room scheduling single class runs in at any given point of time
- 4. TDMA??

Space multiplexing: Share resource by dividing into smaller pieces

1. CPU scheduling on multiple cores?

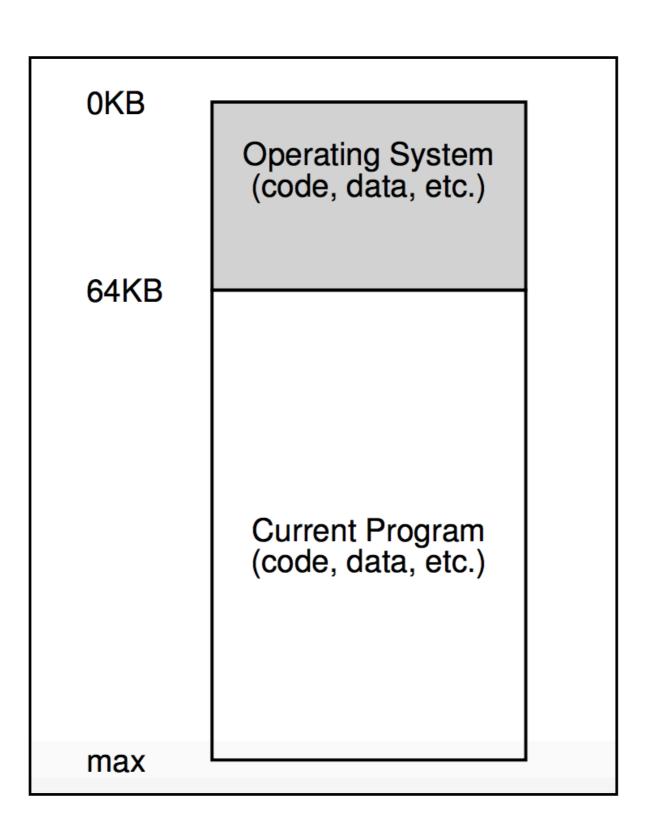
- 1. CPU scheduling on multiple cores?
- 2. Cake sharing

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- 3. Think more?

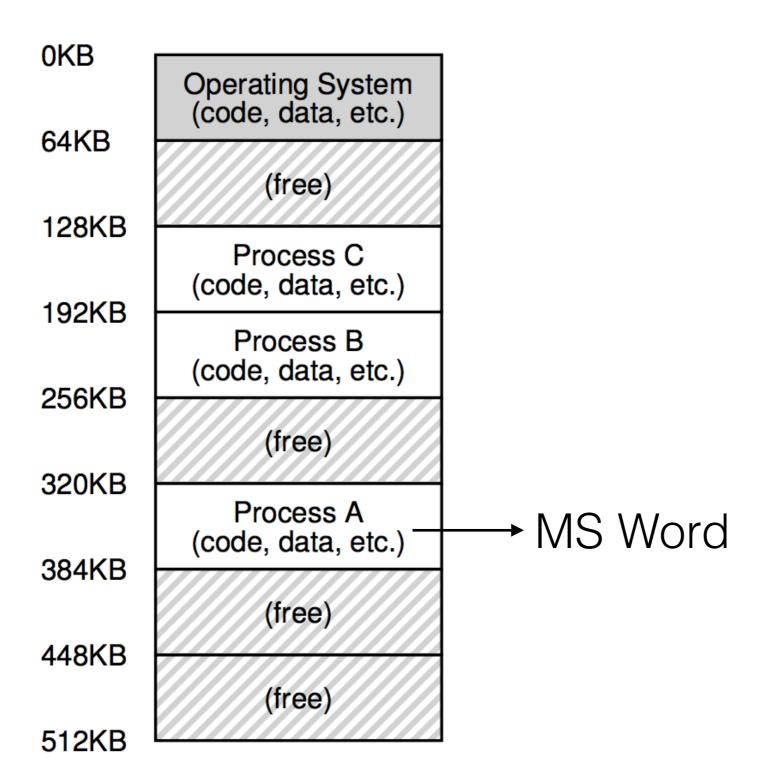
- 1. CPU scheduling on multiple cores?
- 2. Cake sharing
- 3. Think more?
- 4. Memory management

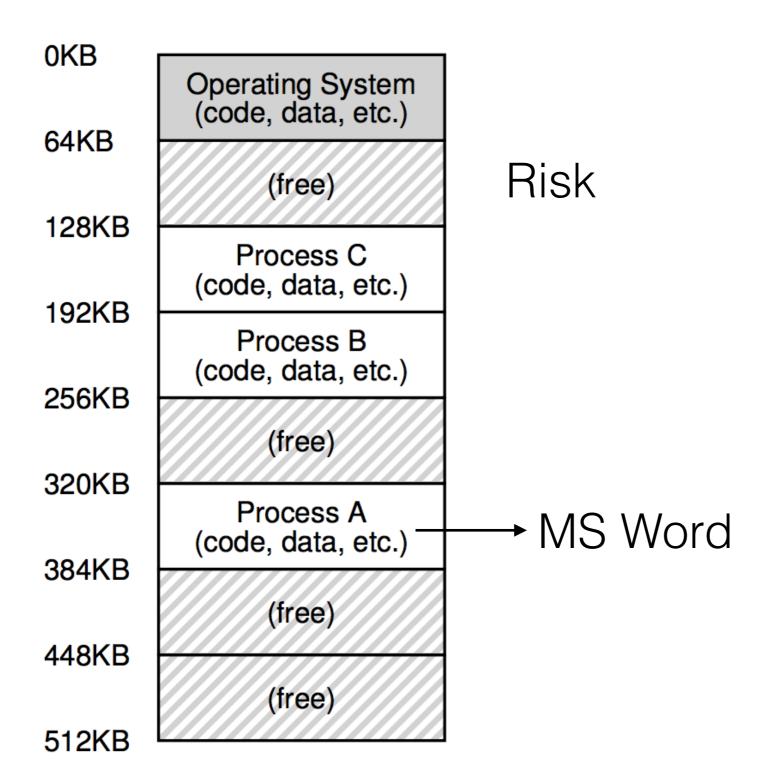
#### Memory Virtualisation

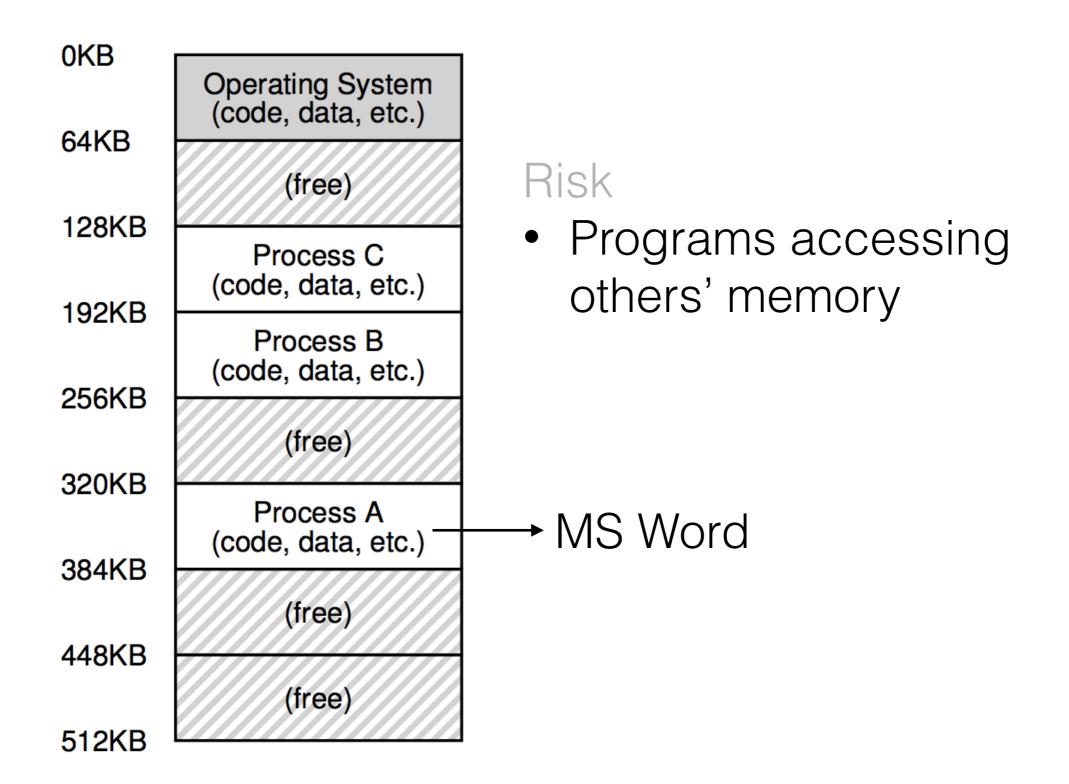
Early days
Single program



0KB	
64KB	Operating System (code, data, etc.)
	(free)
128KB	Process C (code, data, etc.)
192KB	Process B (code, data, etc.)
256KB	(free)
320KB	Process A (code, data, etc.)
384KB	(free)
448KB	(free)
512KB	

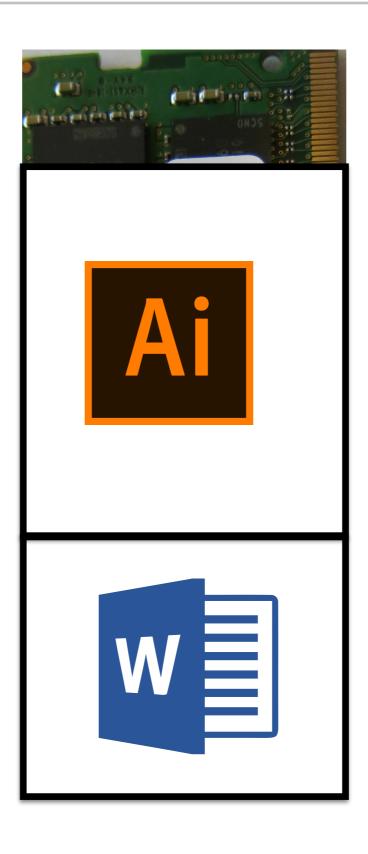


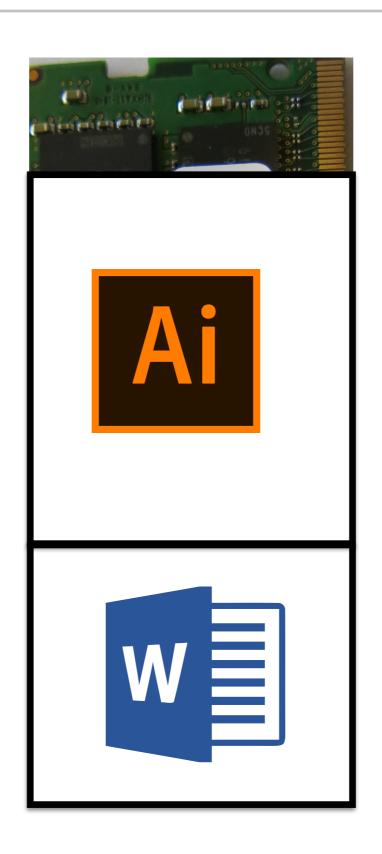


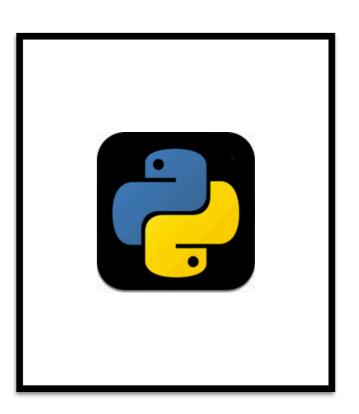




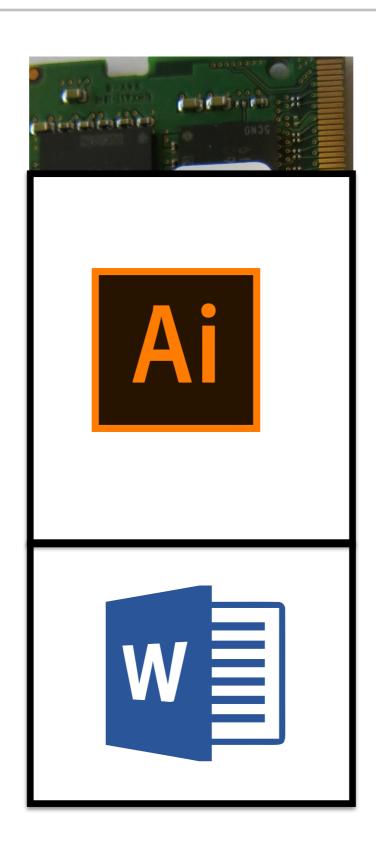


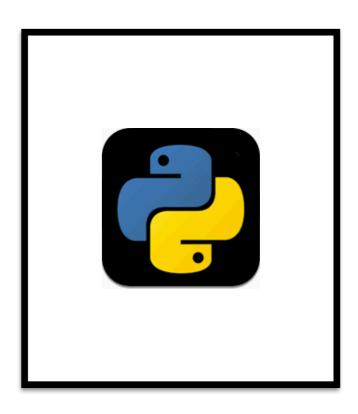






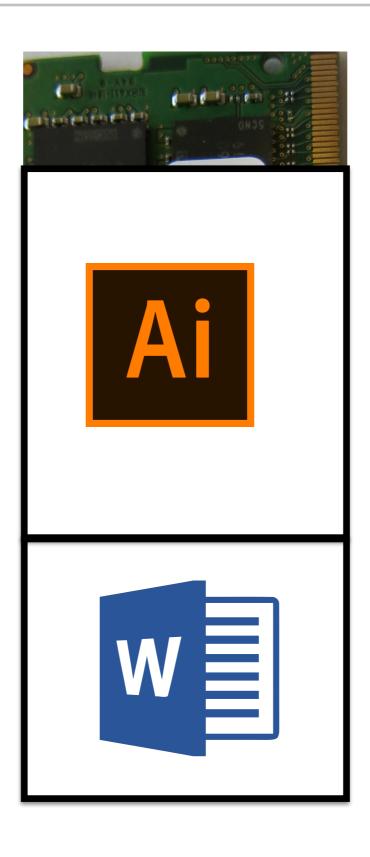
Limited to physical memory on the system

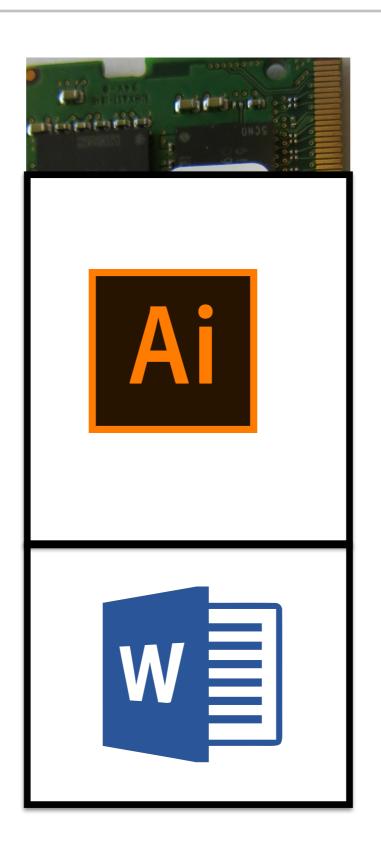


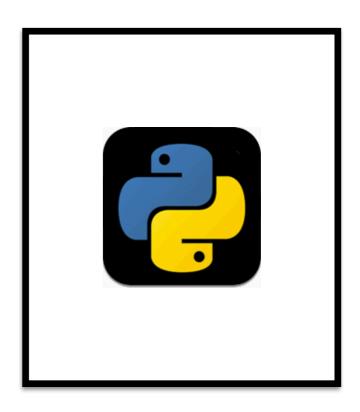




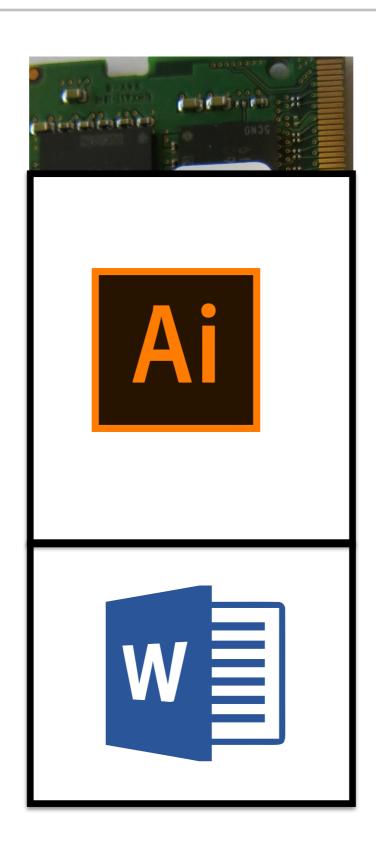


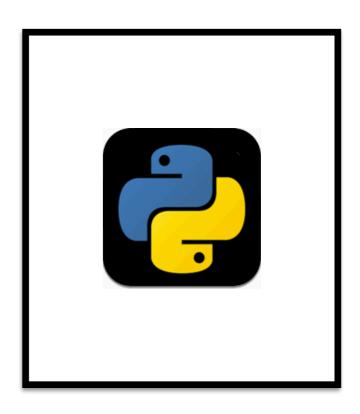


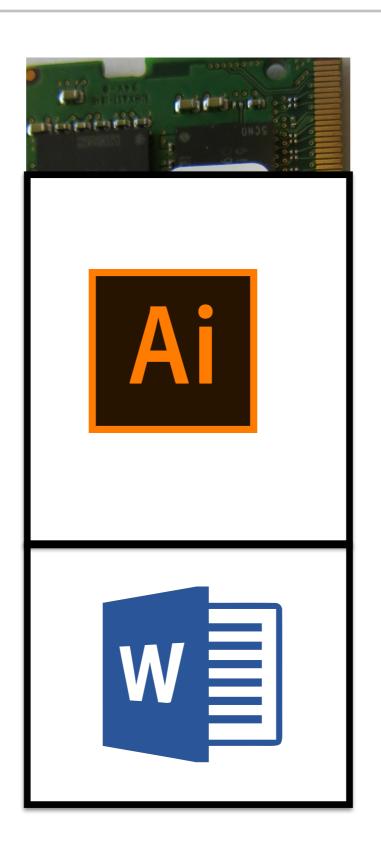


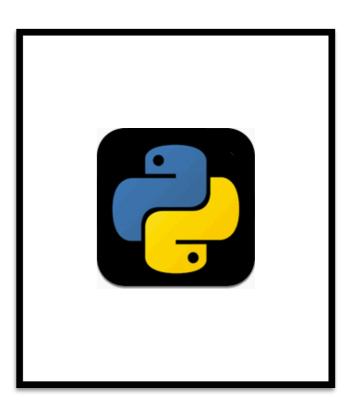


Limited to physical memory on the system

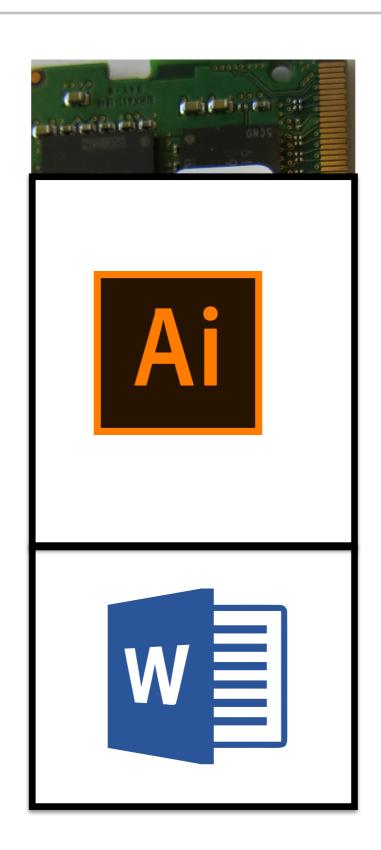


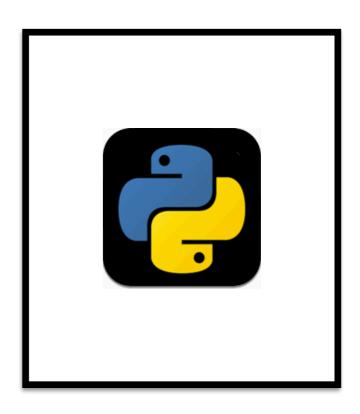


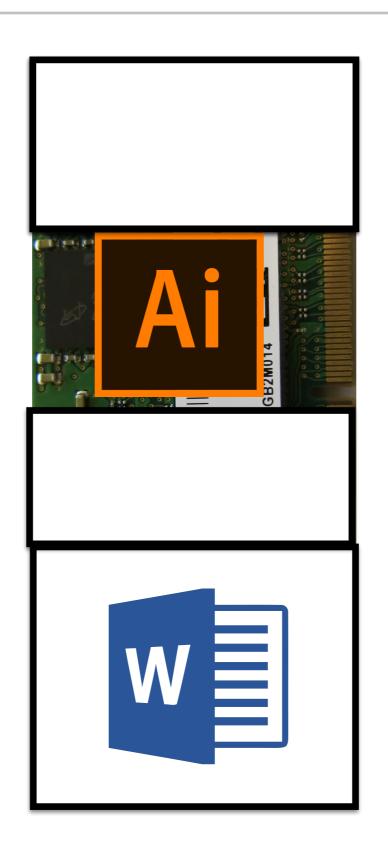


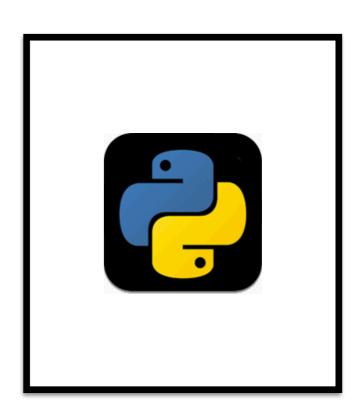


What if process says it wants full memory?

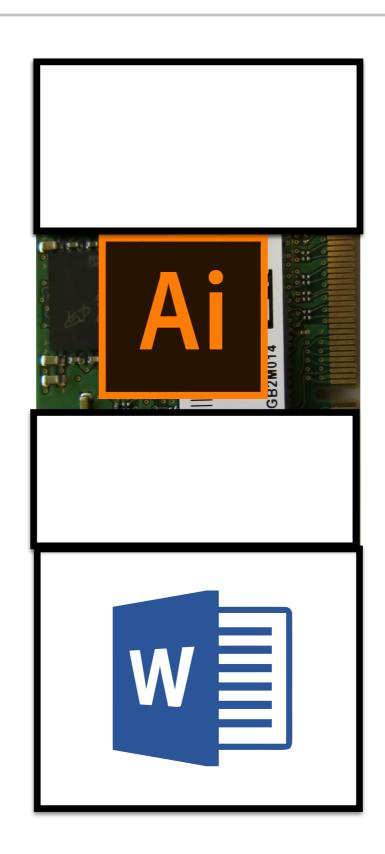


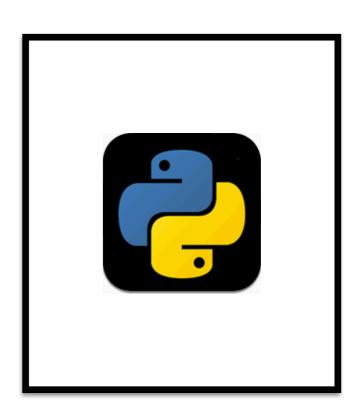






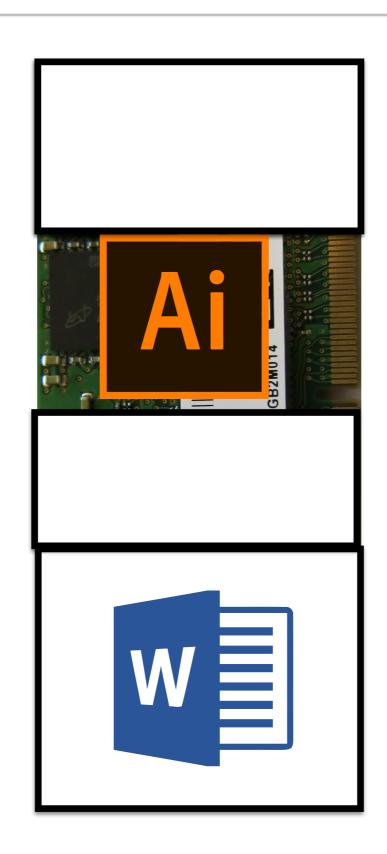
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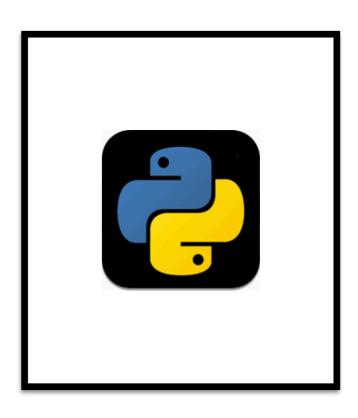


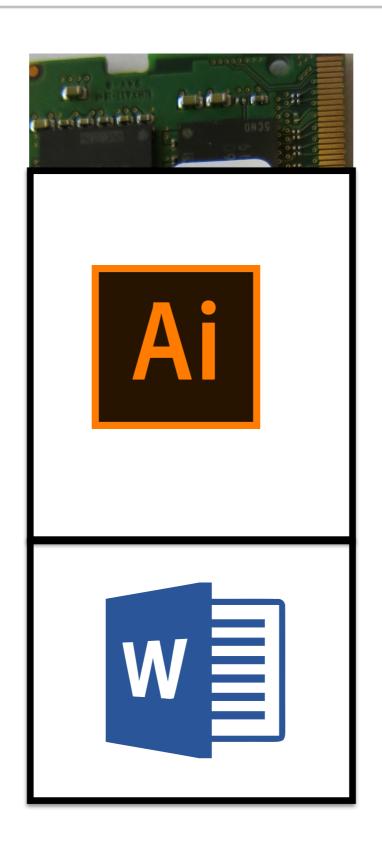


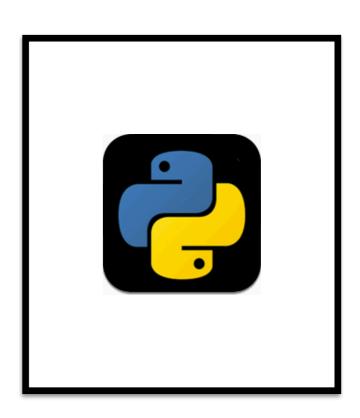
#### Internal Fragmentation

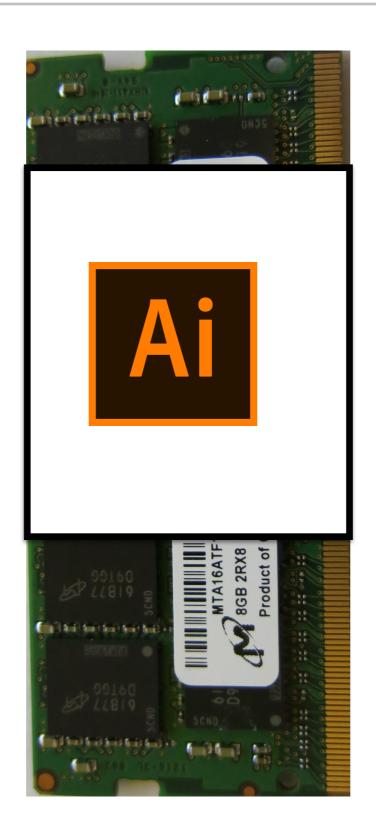
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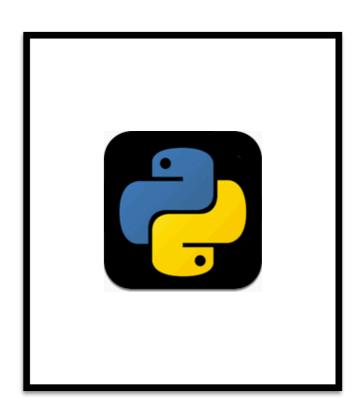








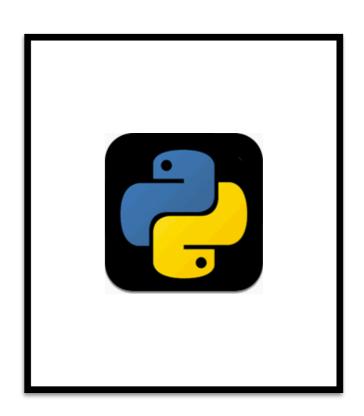




Can Python run now?

Total memory Memory req for
Illustrator > =
Memory req for
Python



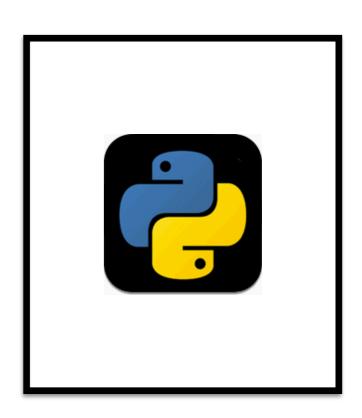


# External Fragmentation

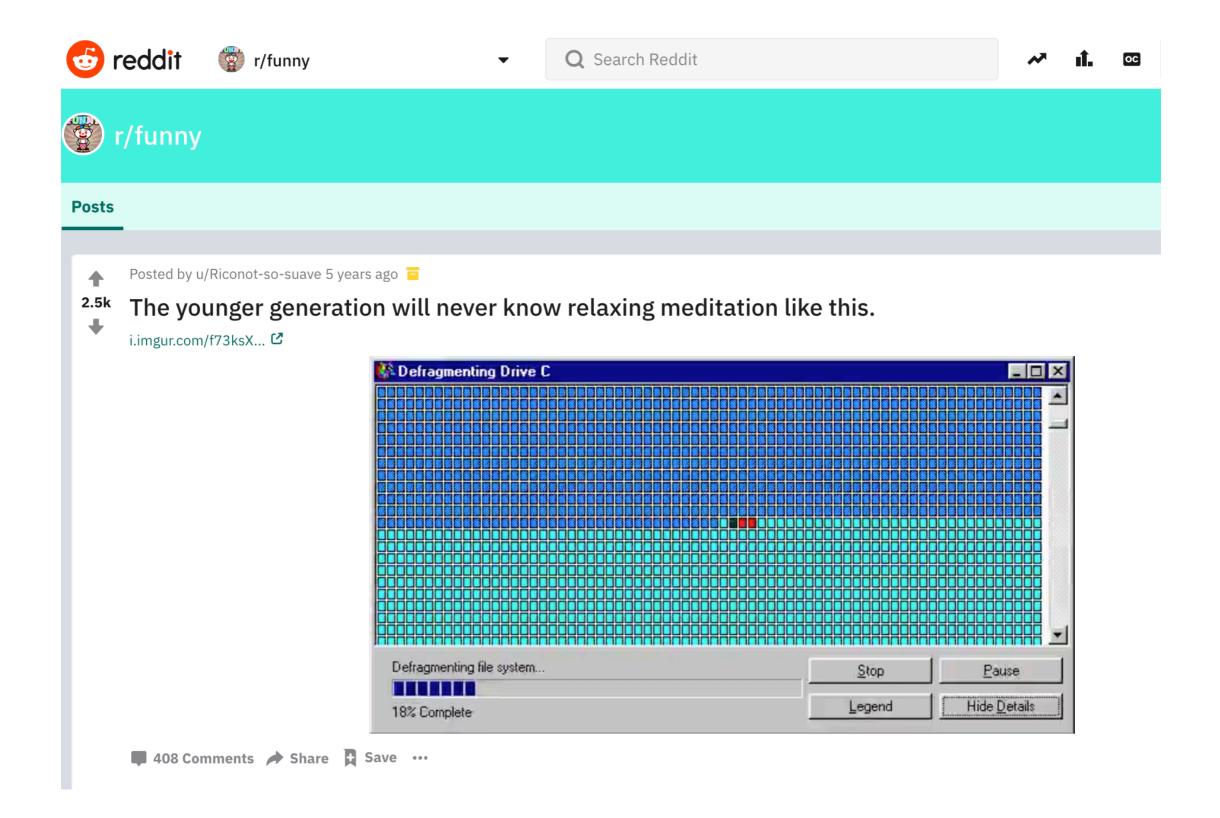
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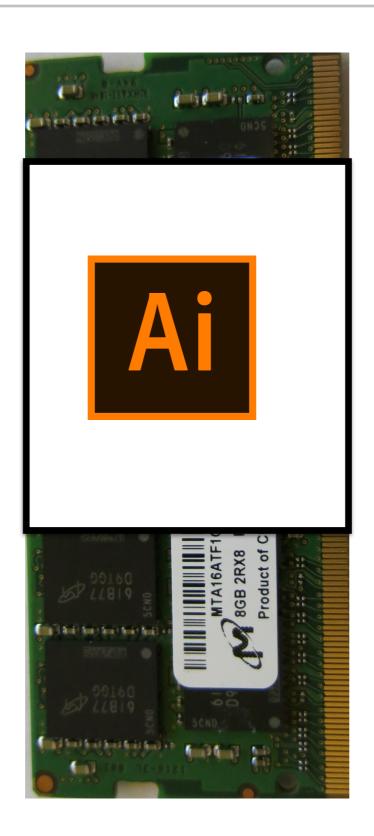


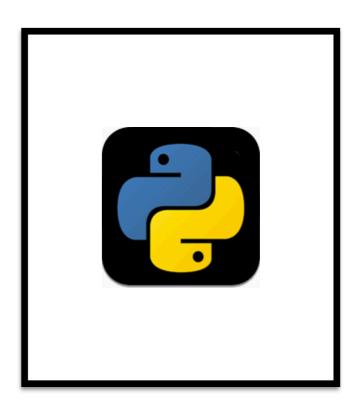


# Defragmentation Memories ...



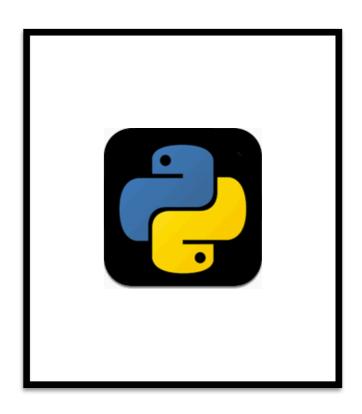
#### Defragmentation



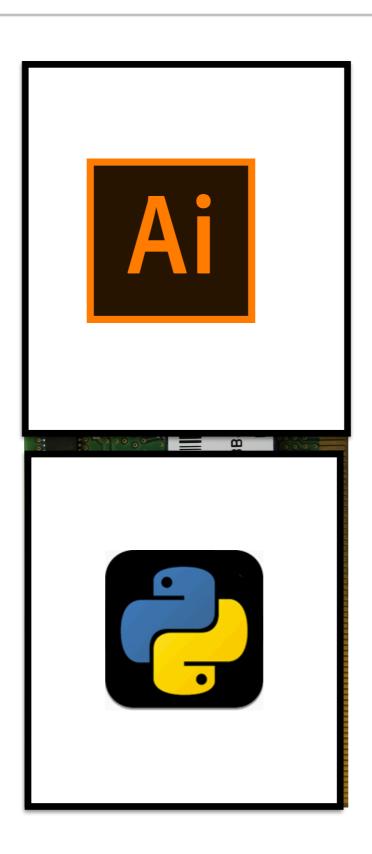


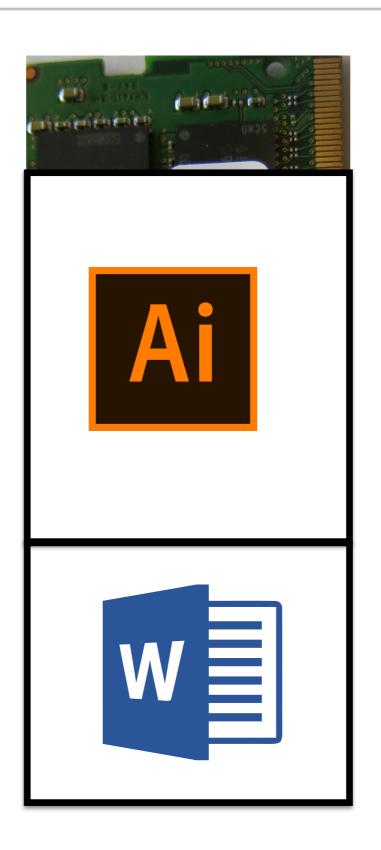
#### Defragmentation

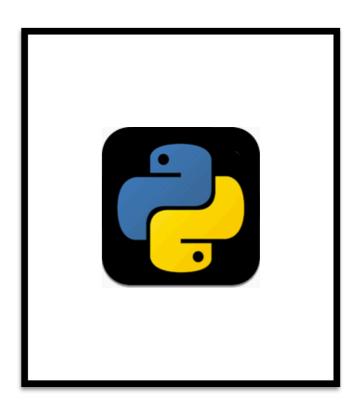


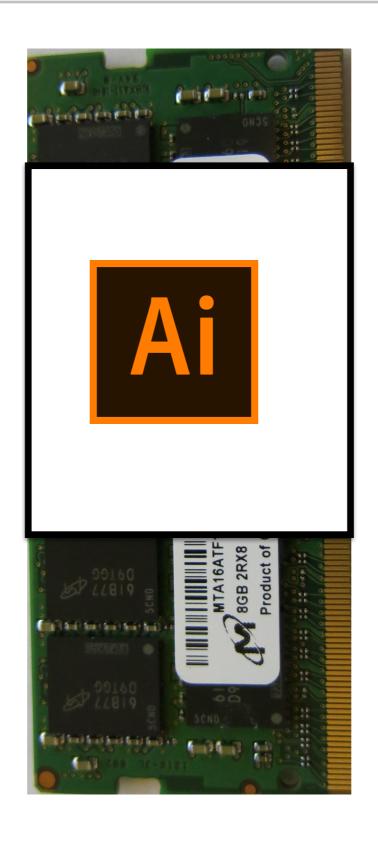


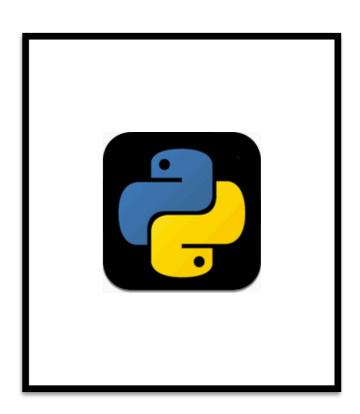
#### Defragmentation







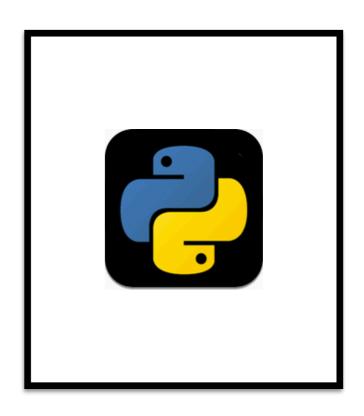




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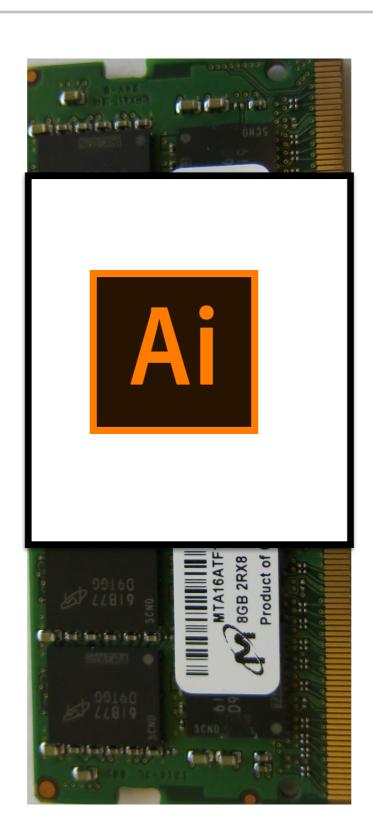


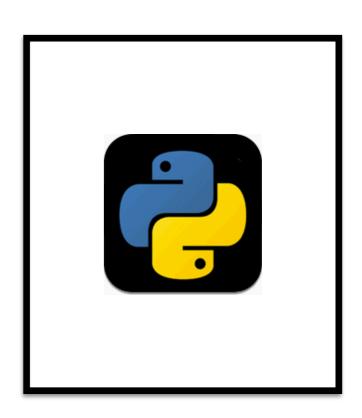


# External Fragmentation

Can Python run now?

Total memory Memory req for
Illustrator > =
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# Goals of OS for Memory Virtualisation/ Management

- 1. Transparency
  - 1. Physical memory is invisible to user program
  - 2. Program thinks it has own private large (contiguous + plentiful) memory
- 2. Efficiency
  - 1. Not taking very long
  - 2. Not taking too much space
- 3. Protection/Isolation
  - 1. Protect processes from each other

### Memory Interface

- 1. Load (address)
- 2. Store (address, value)

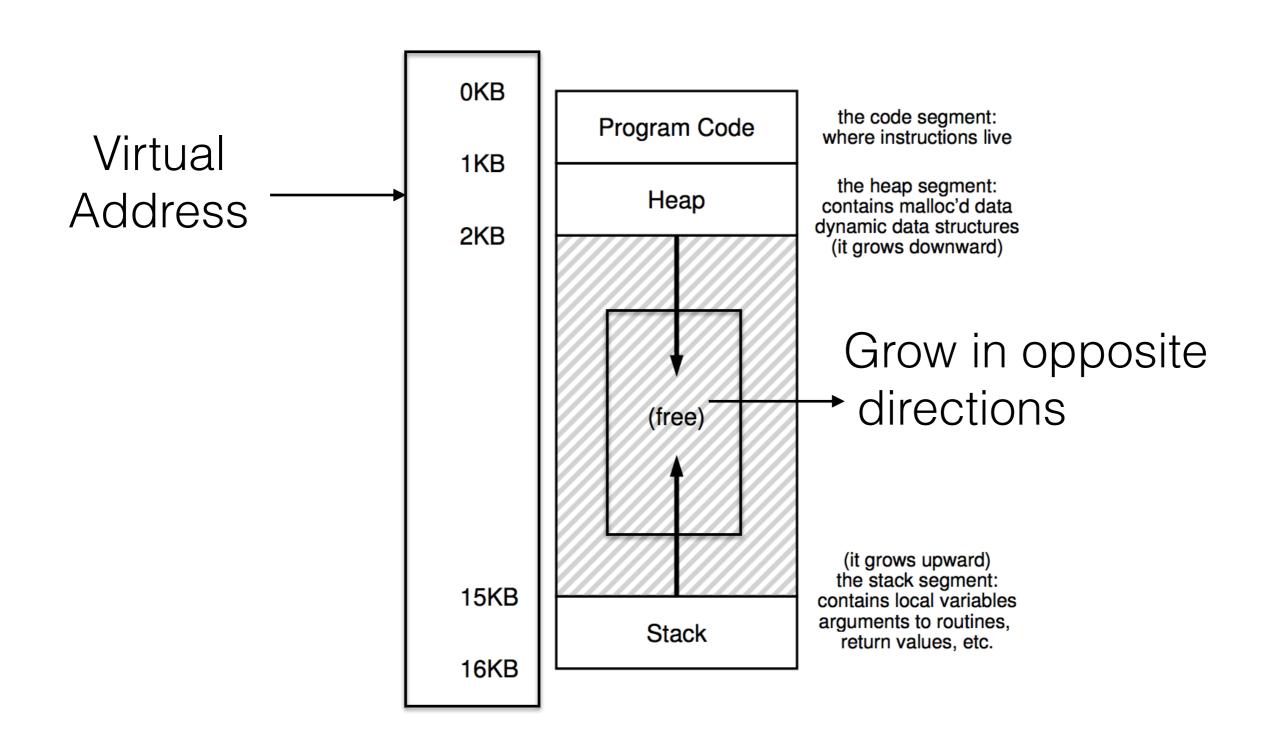
1. Abstraction: Break the connection between physical memory and an address

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- Data accessed using memory interface is virtual address

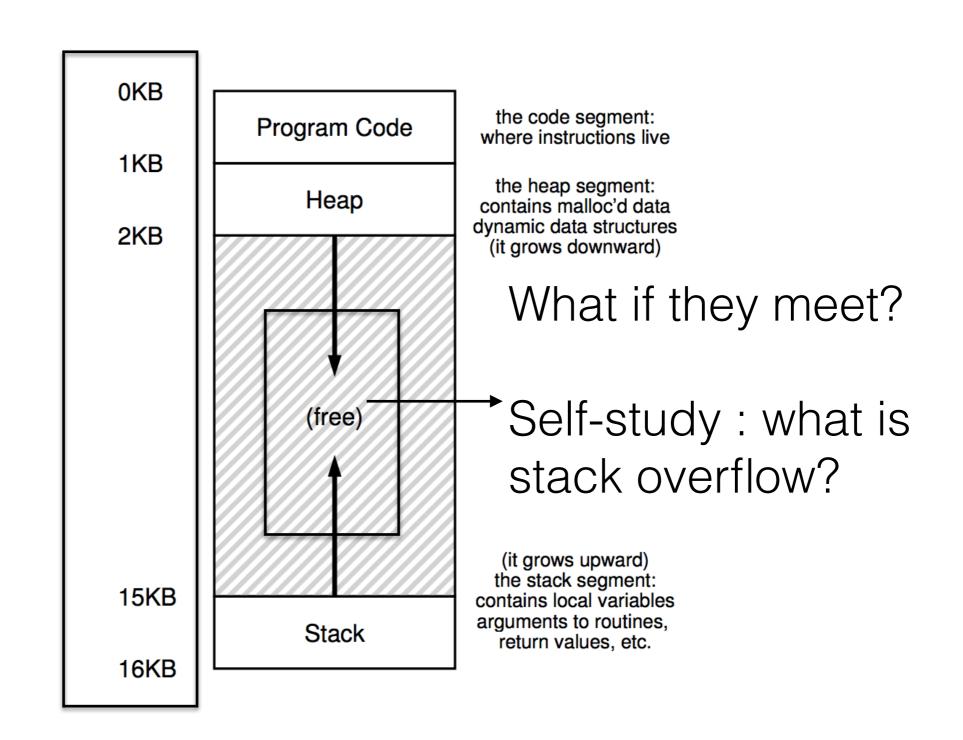
- 1. Abstraction: Break the connection between physical memory and an address
- 2. Data accessed using memory interface is virtual address
  - 1. Physical address points to memory

- 1. Abstraction: Break the connection between physical memory and an address
- 2. Data accessed using memory interface is virtual address
  - 1. Physical address points to memory
  - 2. Virtual address points to something *acting like memory*

# Address Space



#### Stack Overflow?!



#### Stack Overflow?!

```
def fib(n):
    if n==1 or n==0:
        return 1
    else:
        return n*fib(n-1)
```

### Exec Revisited

fork\_same\_address.c

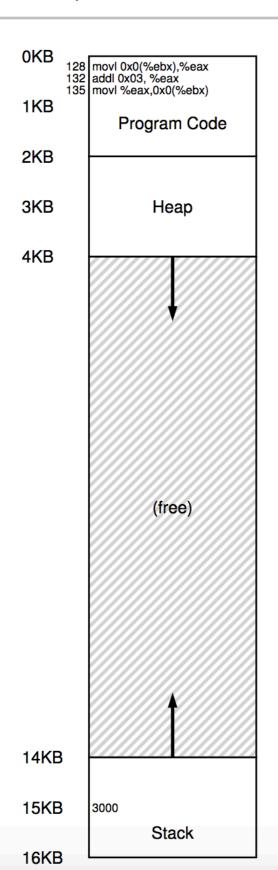
```
void func() {
  int x = 3000; //
  x = x + 3; //
...
```

```
void func() {
  int x = 3000; // Compiler
  x = x + 3; //
  ...
```

```
Compiler 128: movl 0x0(%ebx), %eax 132: addl $0x03, %eax 135: movl %eax, 0x0(%ebx)
```

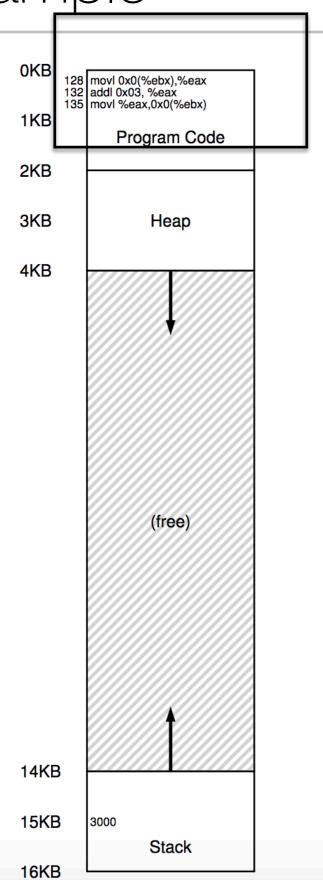
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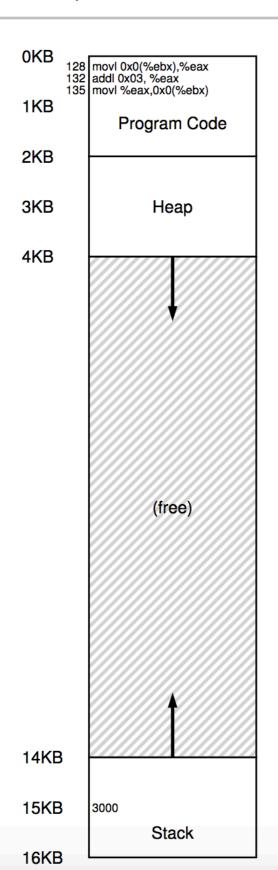
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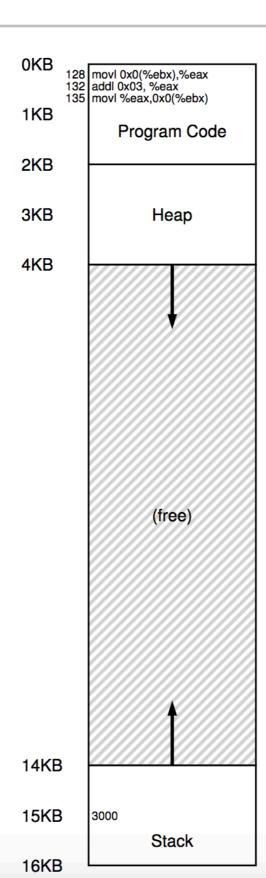
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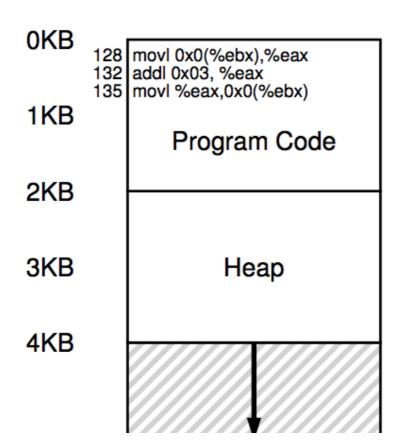
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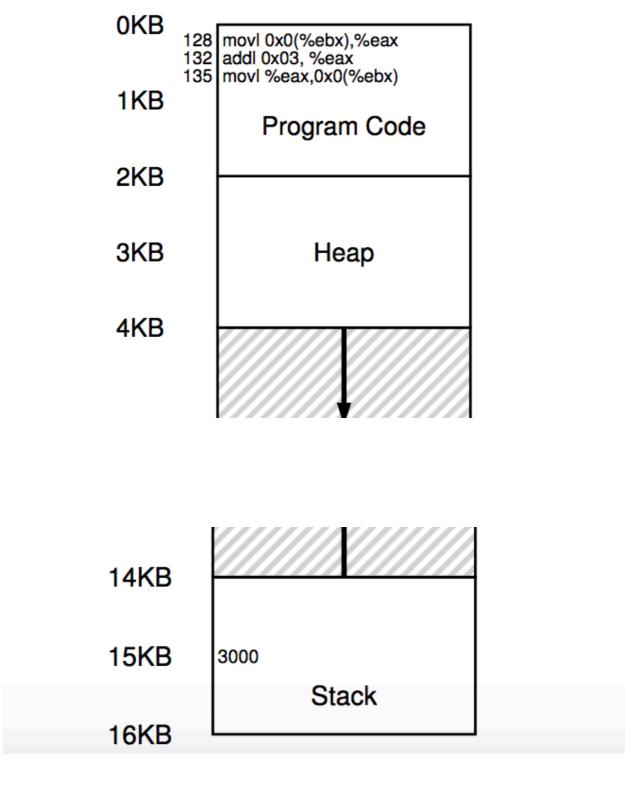
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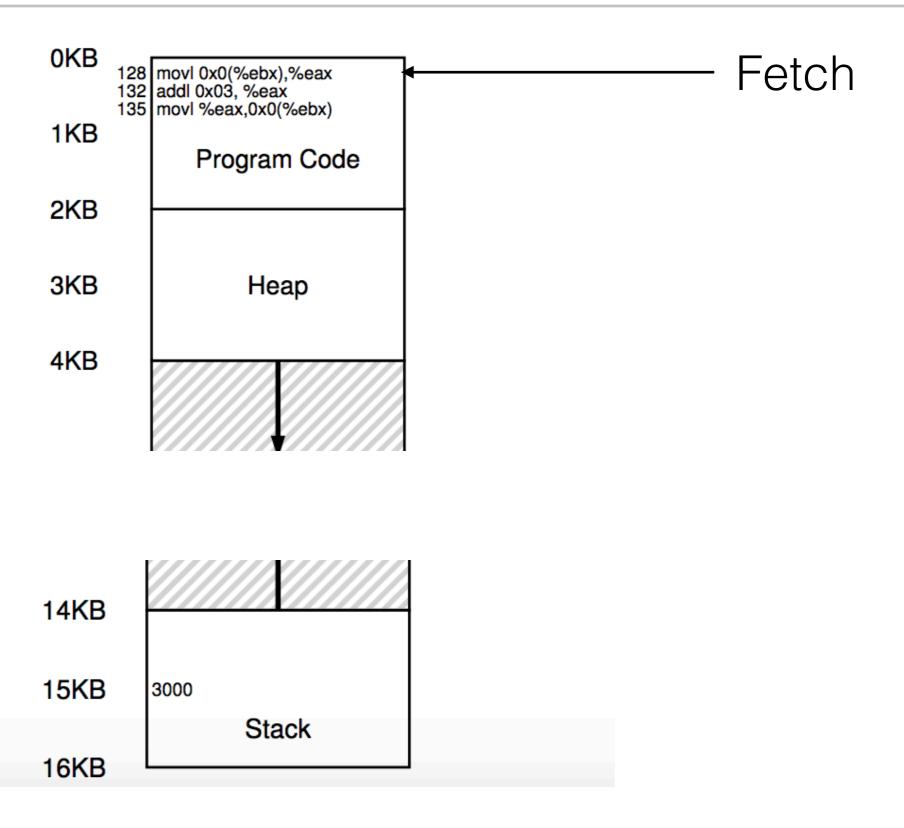


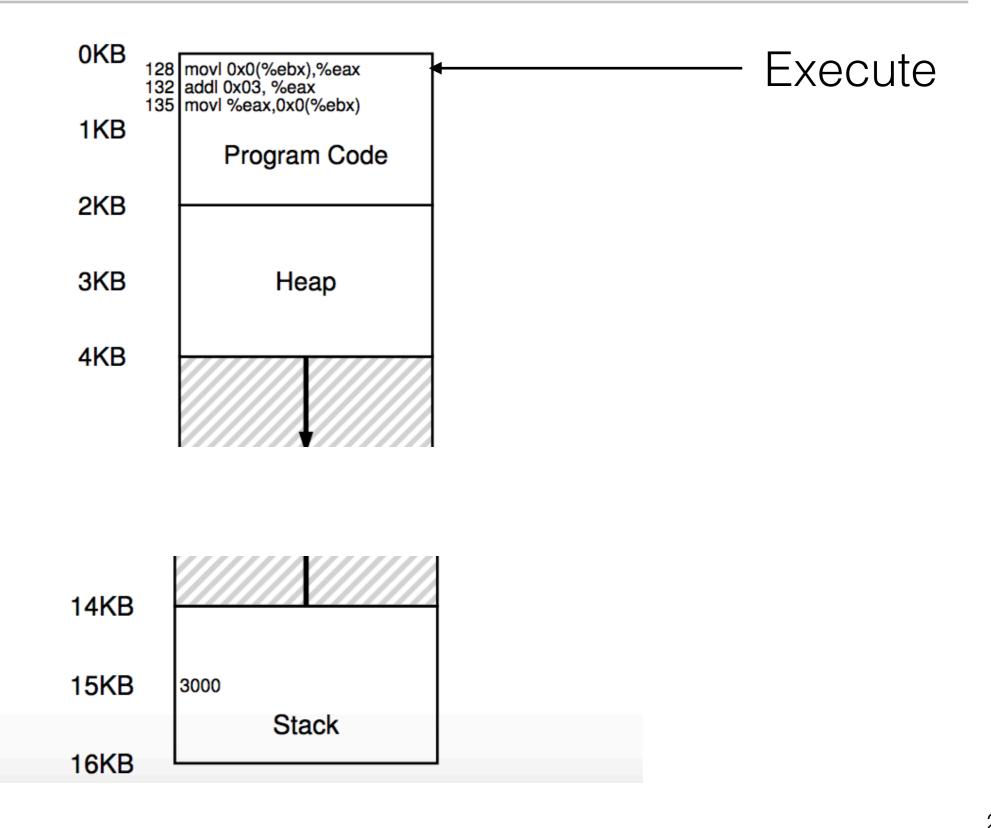


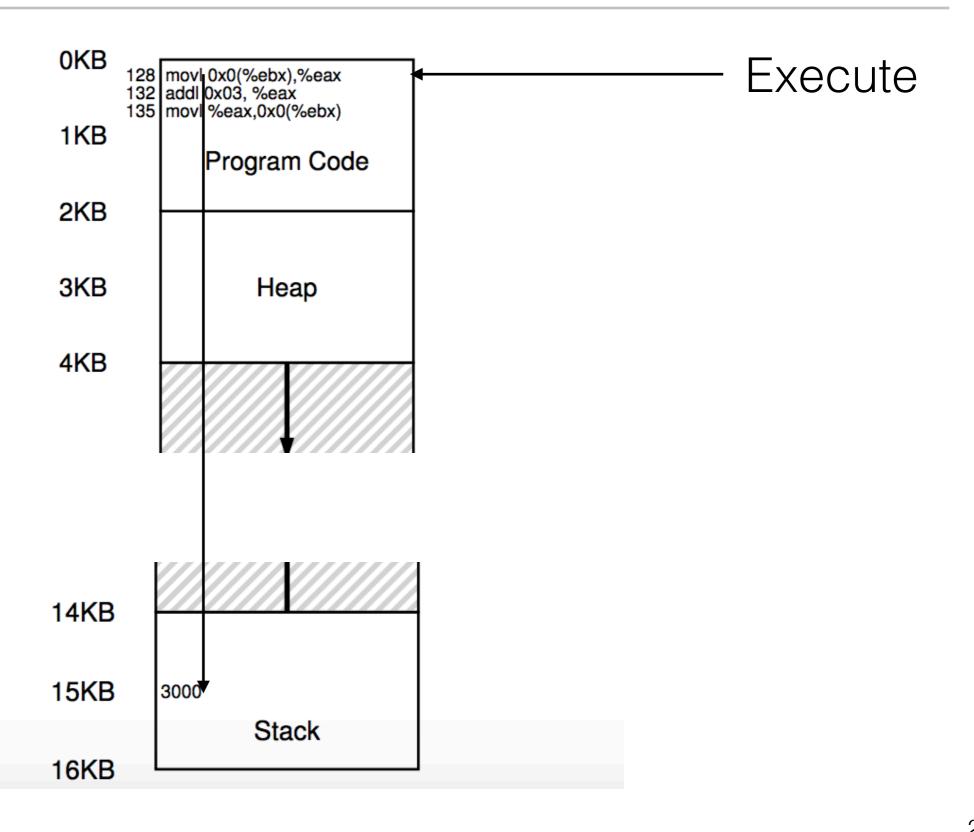
Fetch

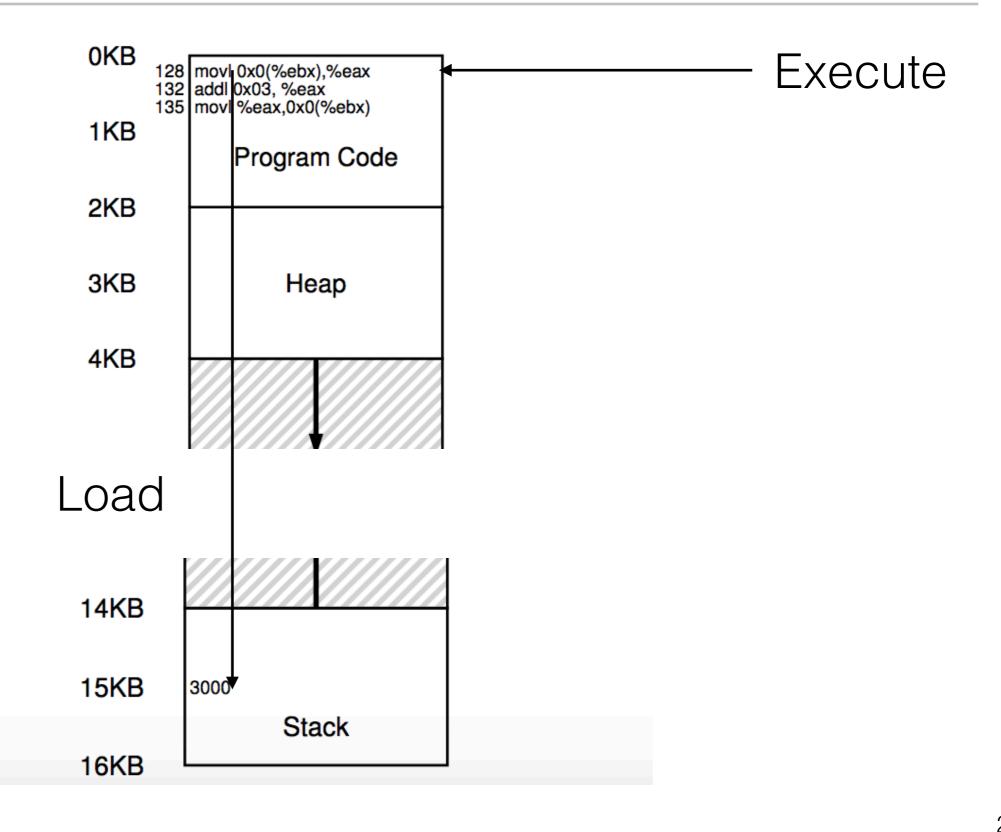


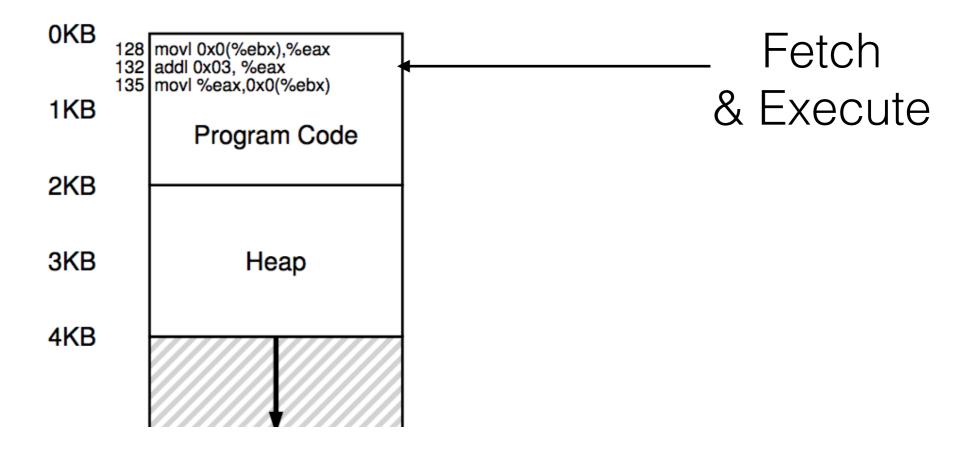
Fetch

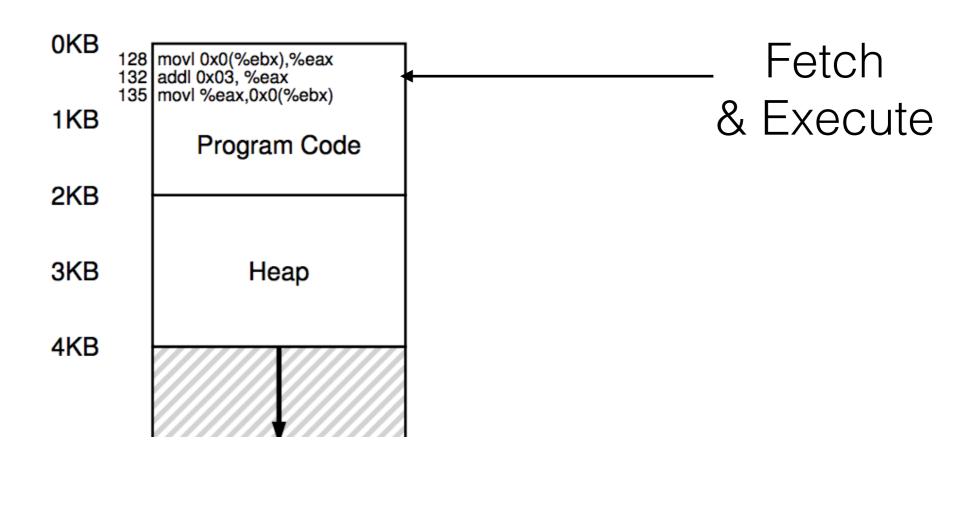


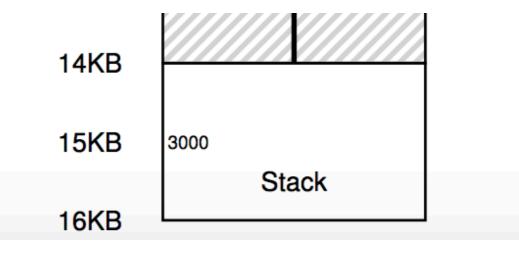


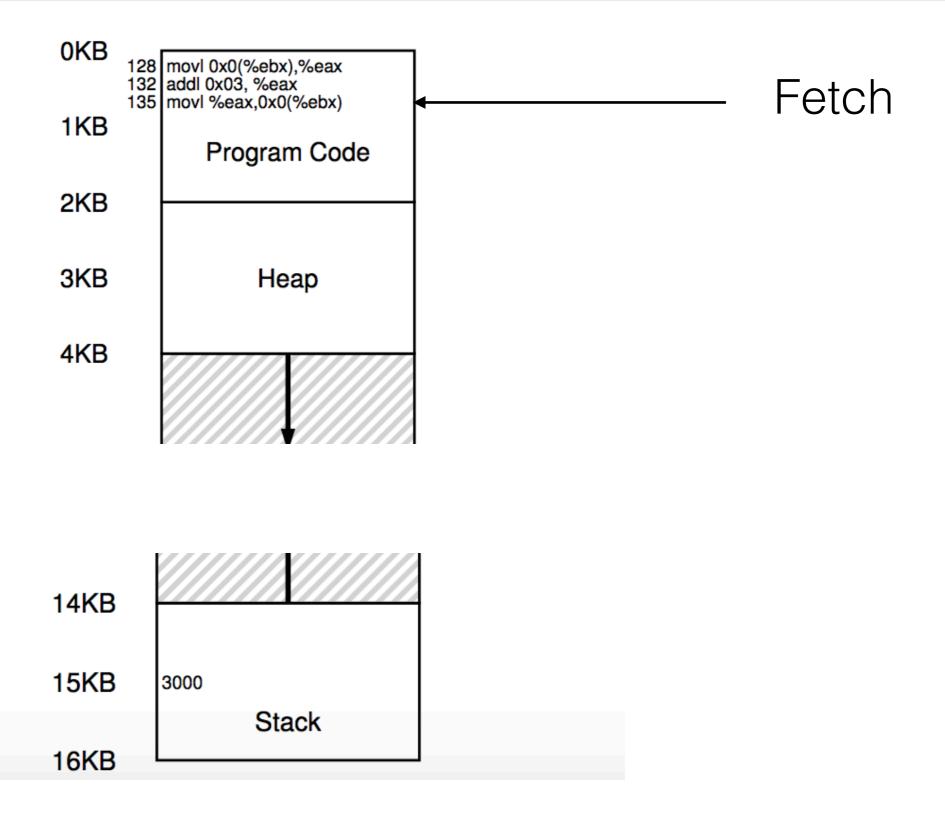


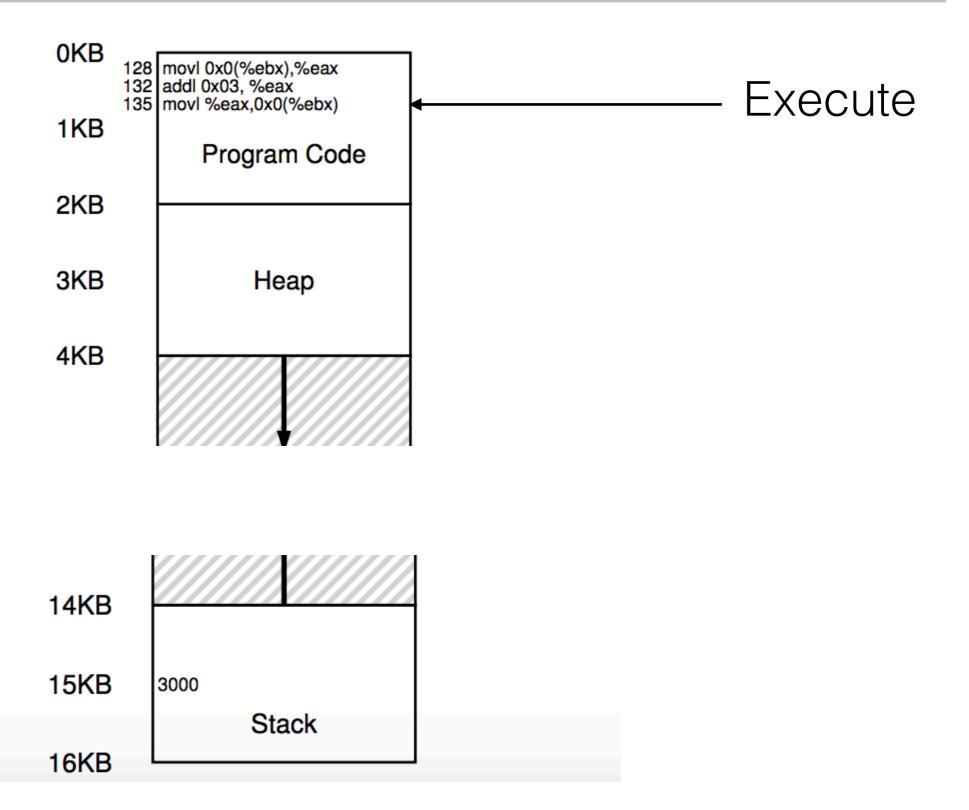


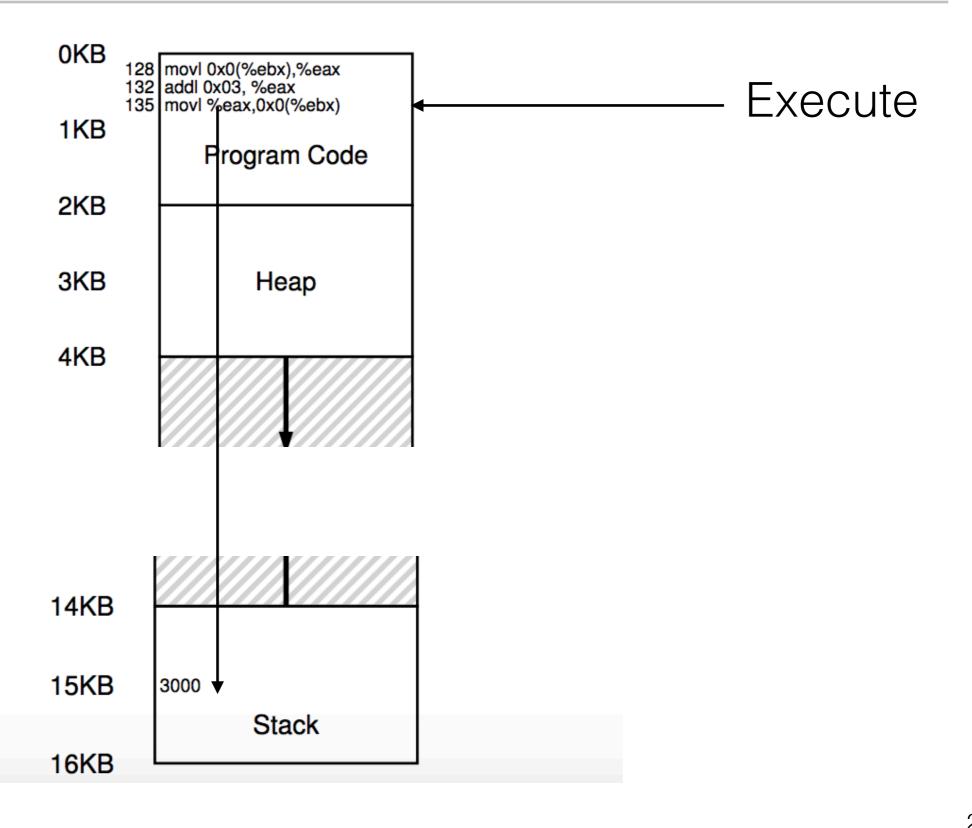


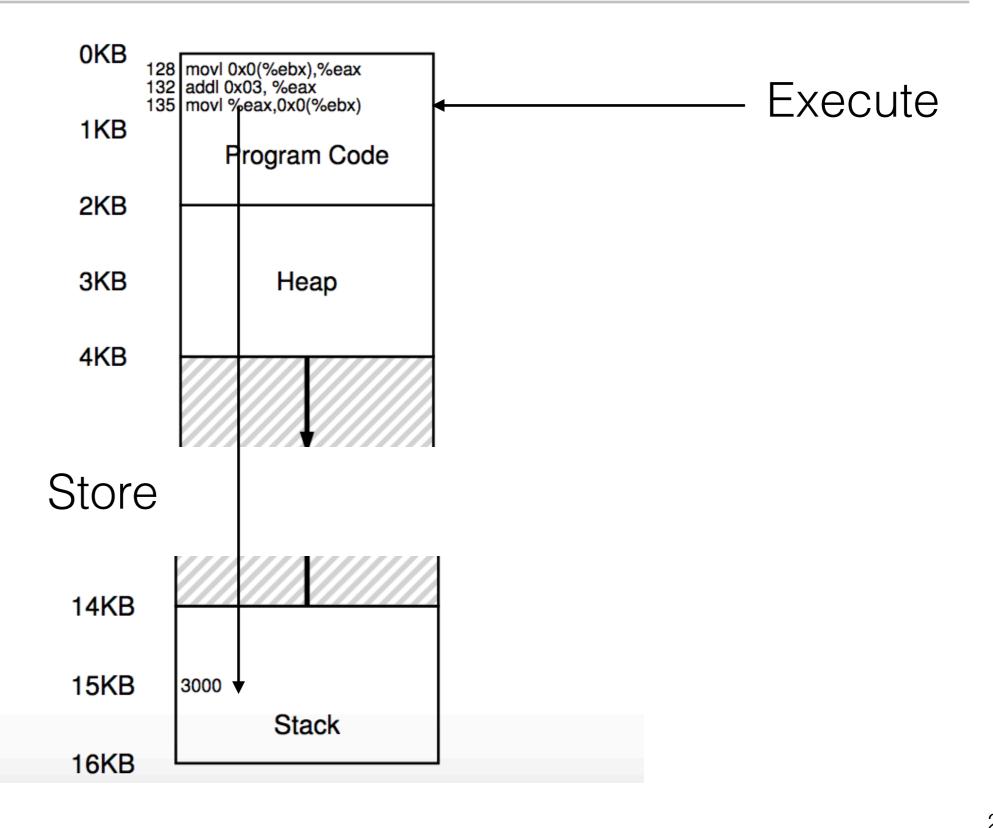




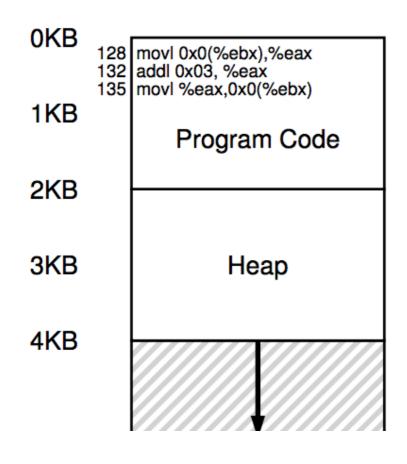




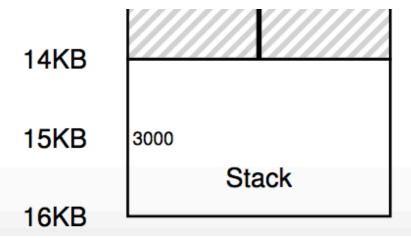




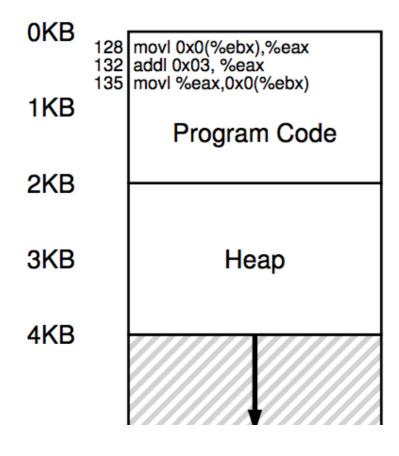
## Pop Quiz

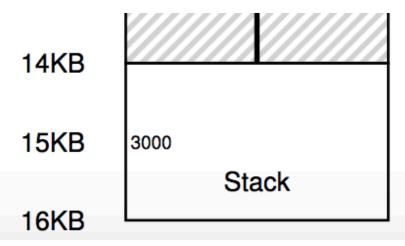


Do all process start and end from 0 KB and 16 KB?

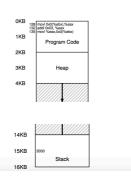


## Relocation

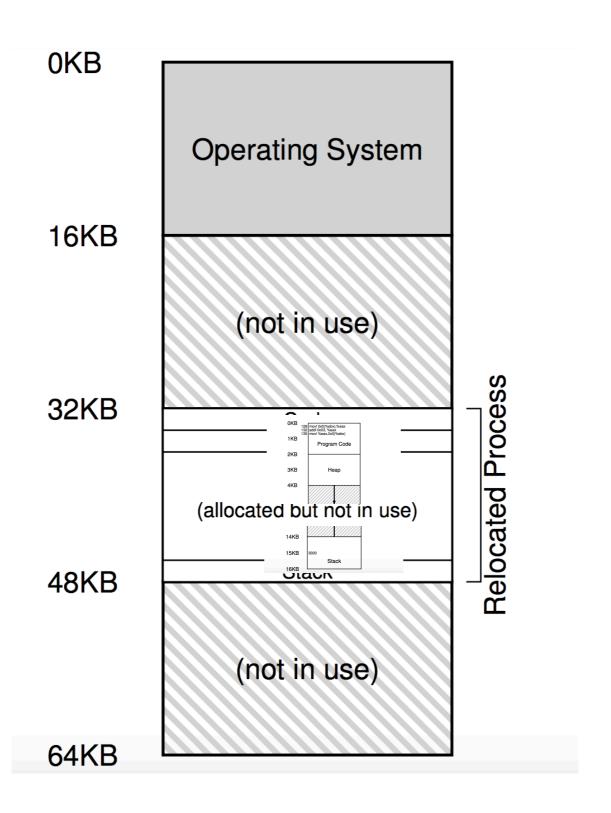




## Relocation



## Relocation



Kernel

CPU

MMU



Kernel

CPU

MMU



Kernel

CPU

MMU

Physical Memory

Virtual Address



Kernel

MMU

CPU

Physical Memory

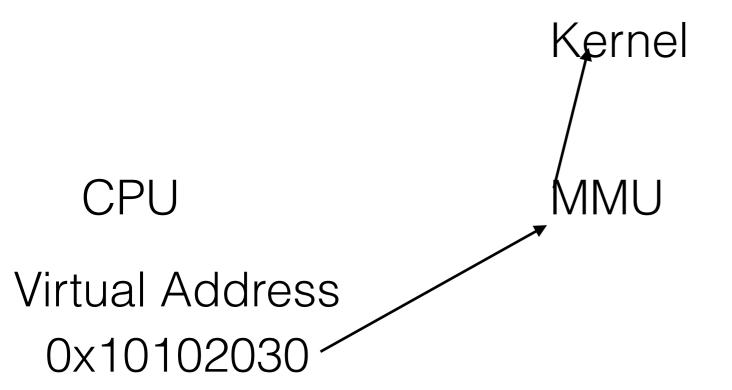
Virtual Address 0x10102030



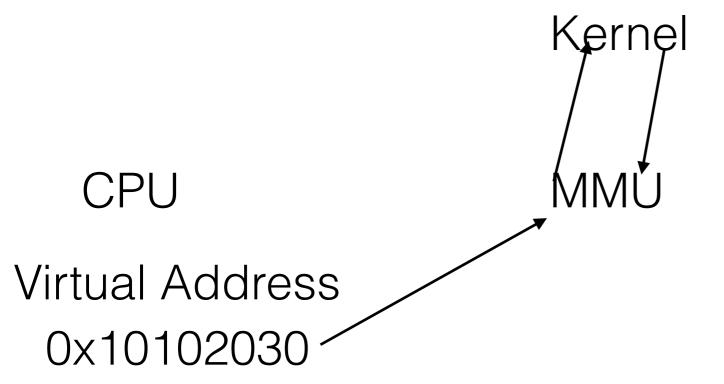
Kernel

CPU MMU
Virtual Address
0x10102030









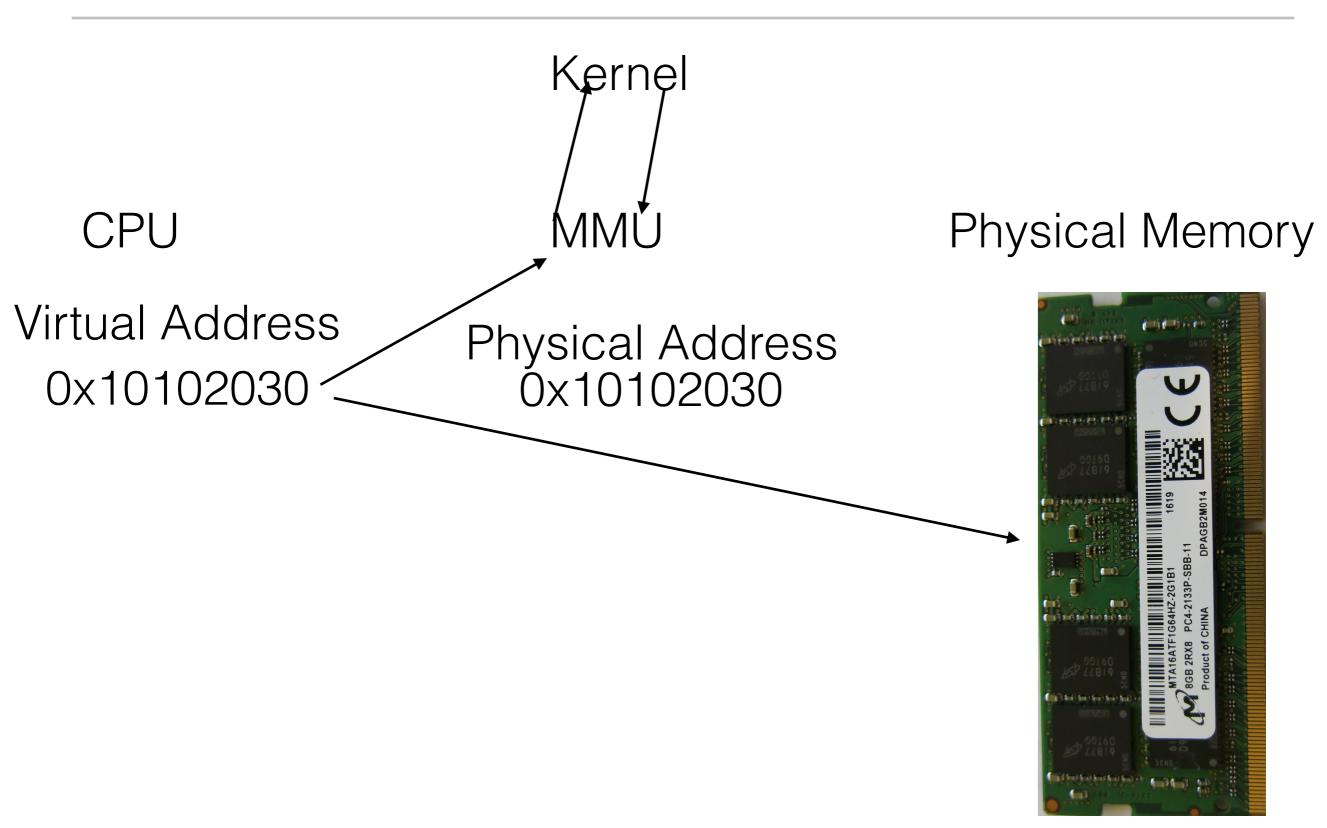


CPU MMU
Virtual Address
0x10102030 Physical Address



CPU MMU
Virtual Address
0x10102030 Physical Address
0x10102030





Kernel **CPU** Virtual Address Physical Address 0x10102030 0x10102030 What if you want to translate same virtual address again?



Kernel **CPU** Physical Memory Virtual Address Physical Address 0x10102030 0x10102030 What if you want to translate same virtual address again?

Kernel **CPU** Virtual Address Physical Address 0x10102030 0x10102030 What if you want to translate same

virtual address again?

Cache!!



Kernel **CPU** Virtual Address Physical Address 0x10102030 0x10102030 What do you do with cache if there

is a context switch?

