Operating Systems Lecture 2: The Process

Nipun Batra

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- 4. Energy efficiency (esp. for mobile systems!)

Process

- Process = Running program
- Review example from previous lecture
 - Output of top
 - Activity monitor

Memory in C

– static: global variable storage, permanent for the entire run of the program.

– stack: local variable storage (automatic, continuous memory).

– heap: dynamic storage (large pool of memory, not allocated in contiguous order).

https://craftofcoding.wordpress.com/2015/12/07/memory-in-c-thestack-the-heap-and-static/

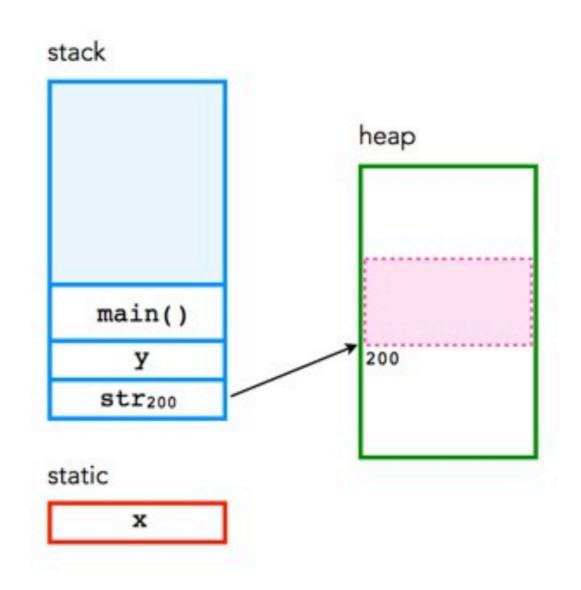
Memory in C

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#include <stdio.h>
#include <stdlib.h>
int x;
int main(void)
{
    int y;
    char *str;
    y = 4;
    printf("stack memory: %d\n", y);
    str = malloc(100*sizeof(char));
    str[0] = 'm';
    printf("heap memory: %c\n", str[0]);
    free(str);
    return 0;
}
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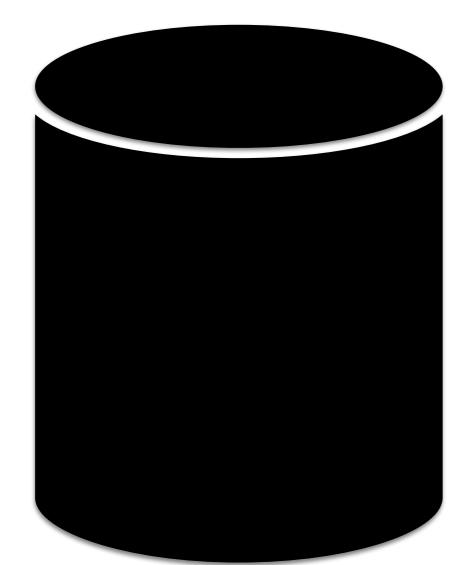
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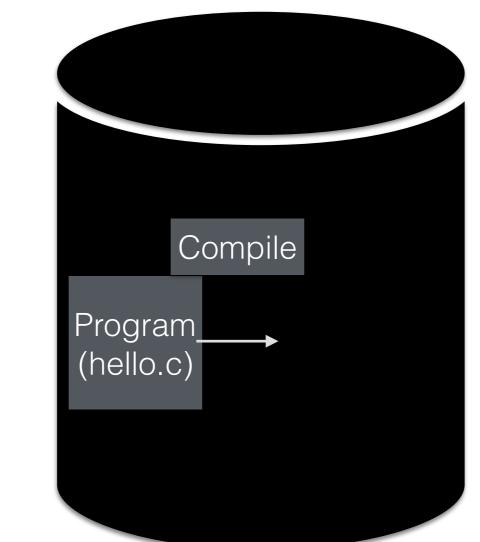
Process Execution



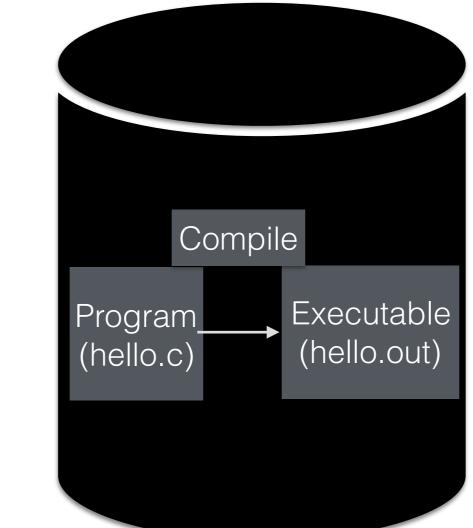
Disk



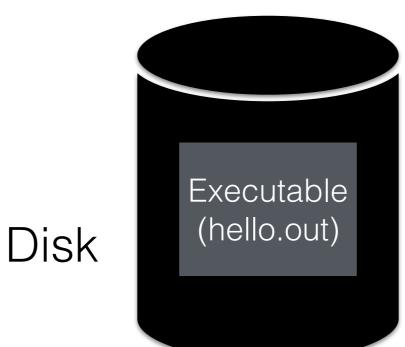
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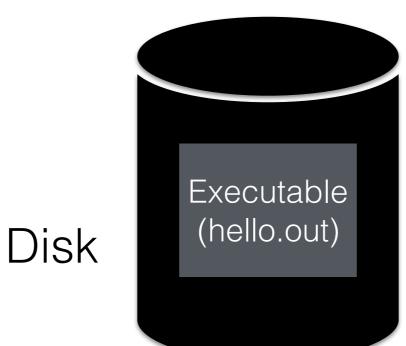


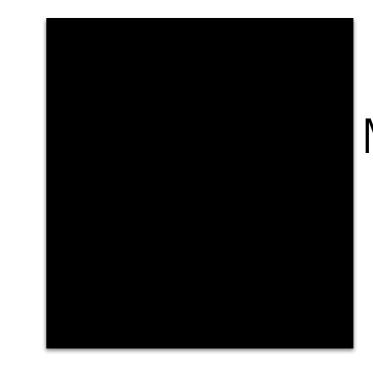
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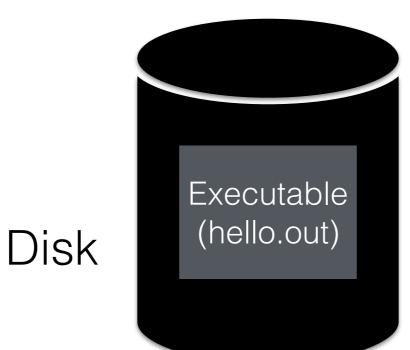
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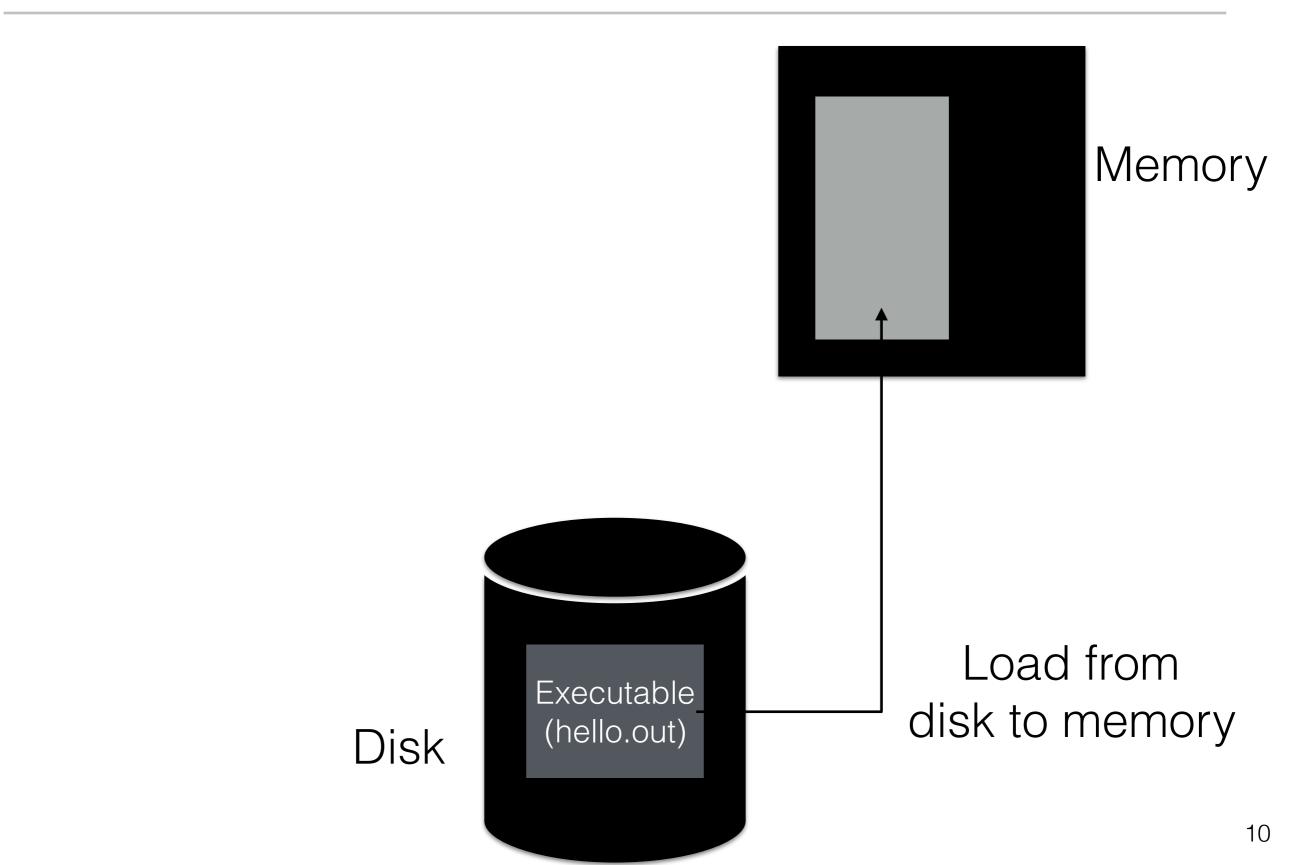


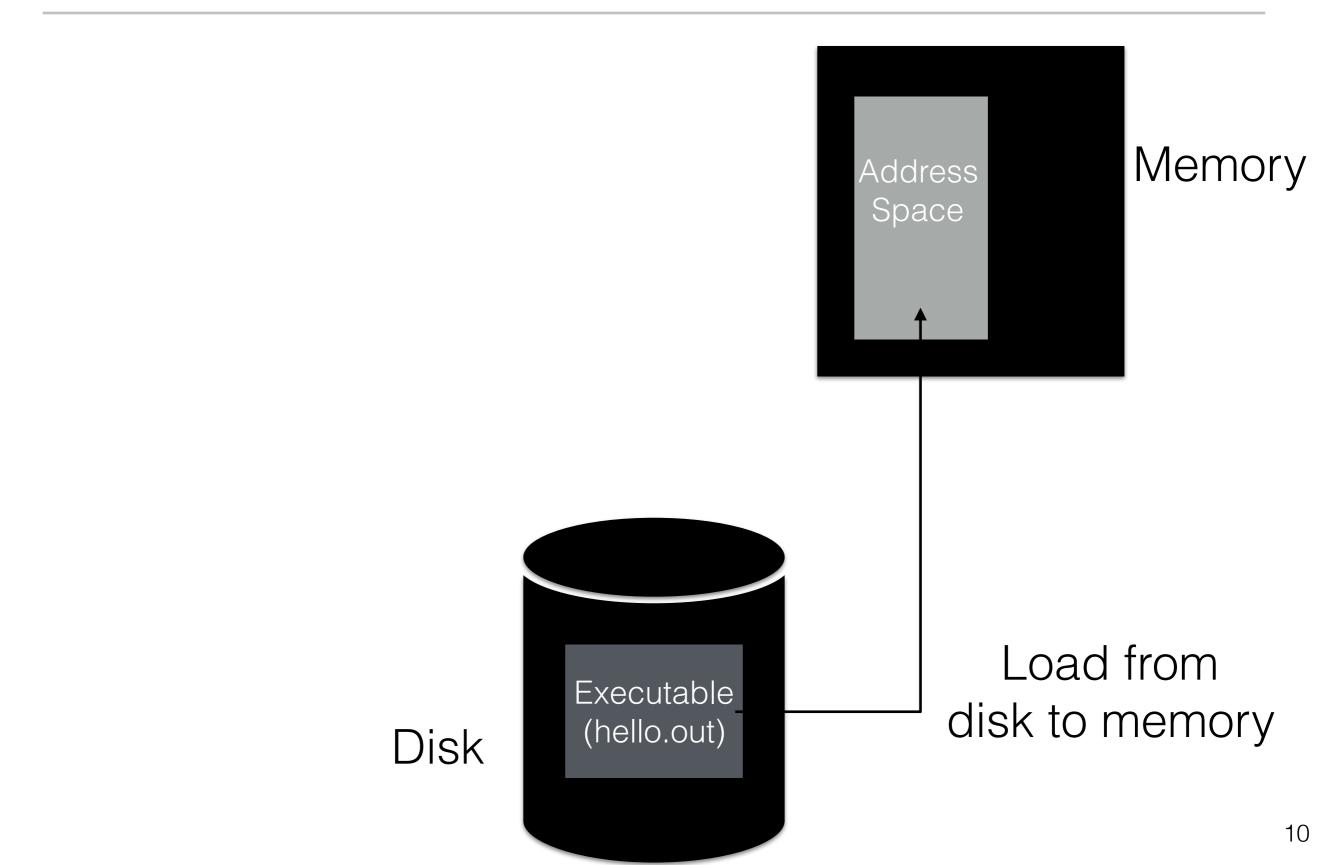


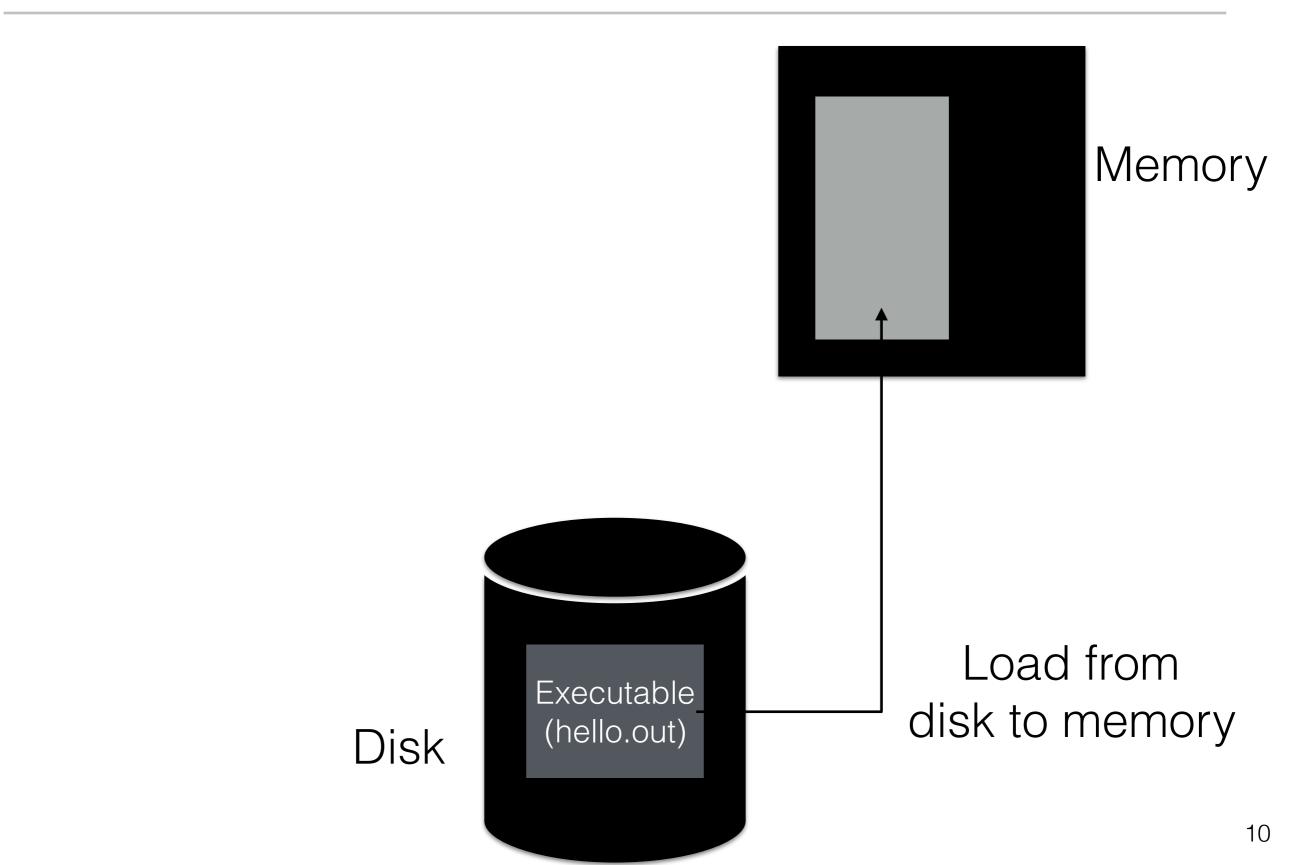


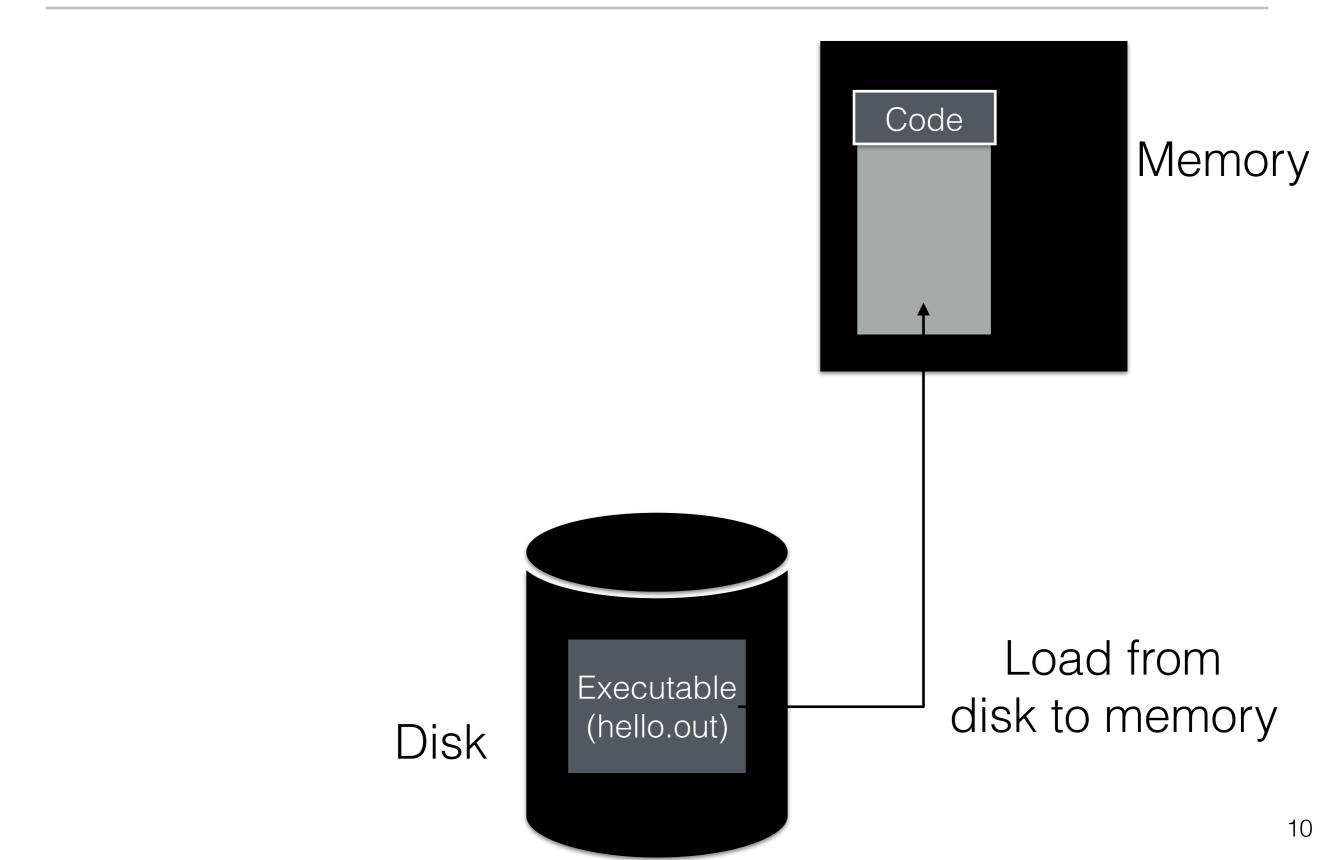
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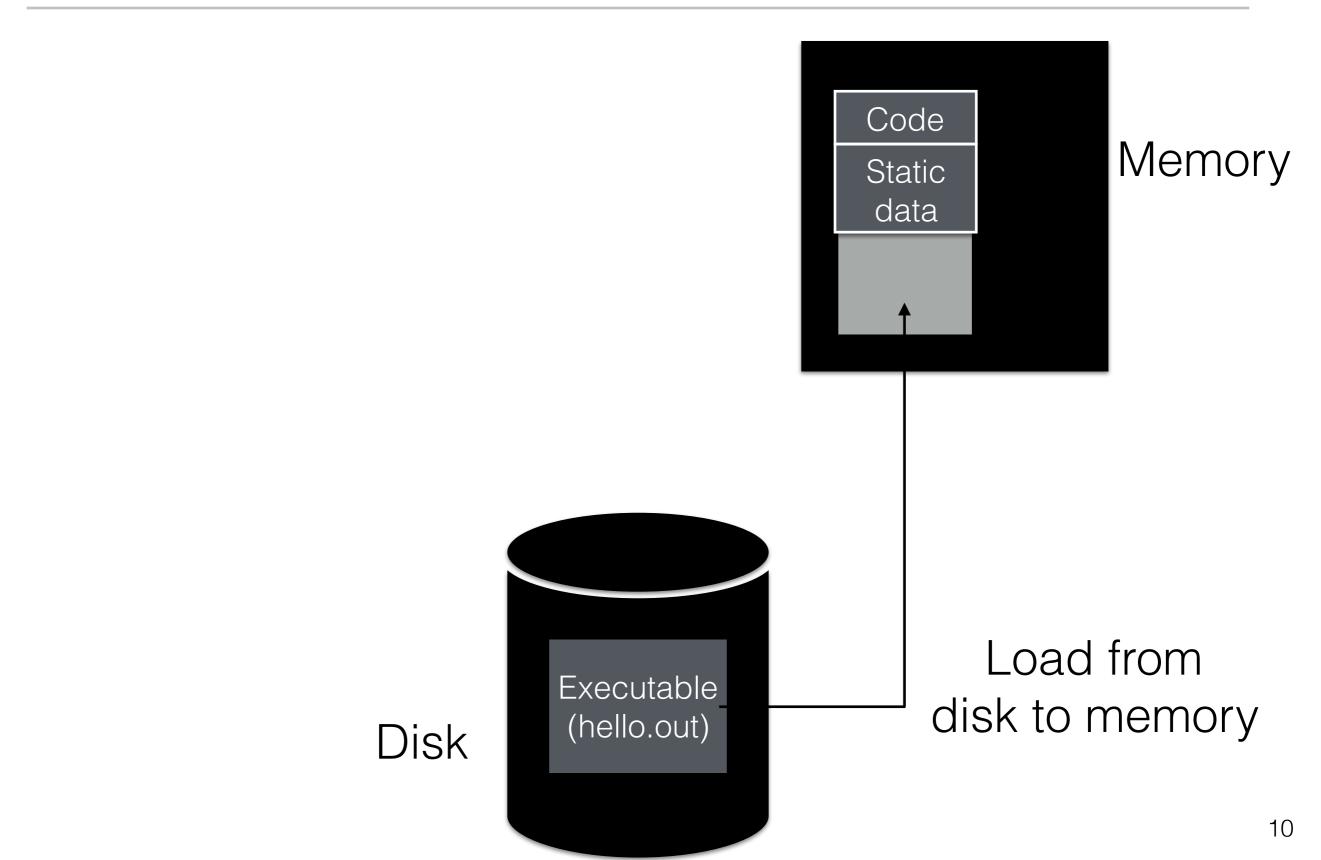


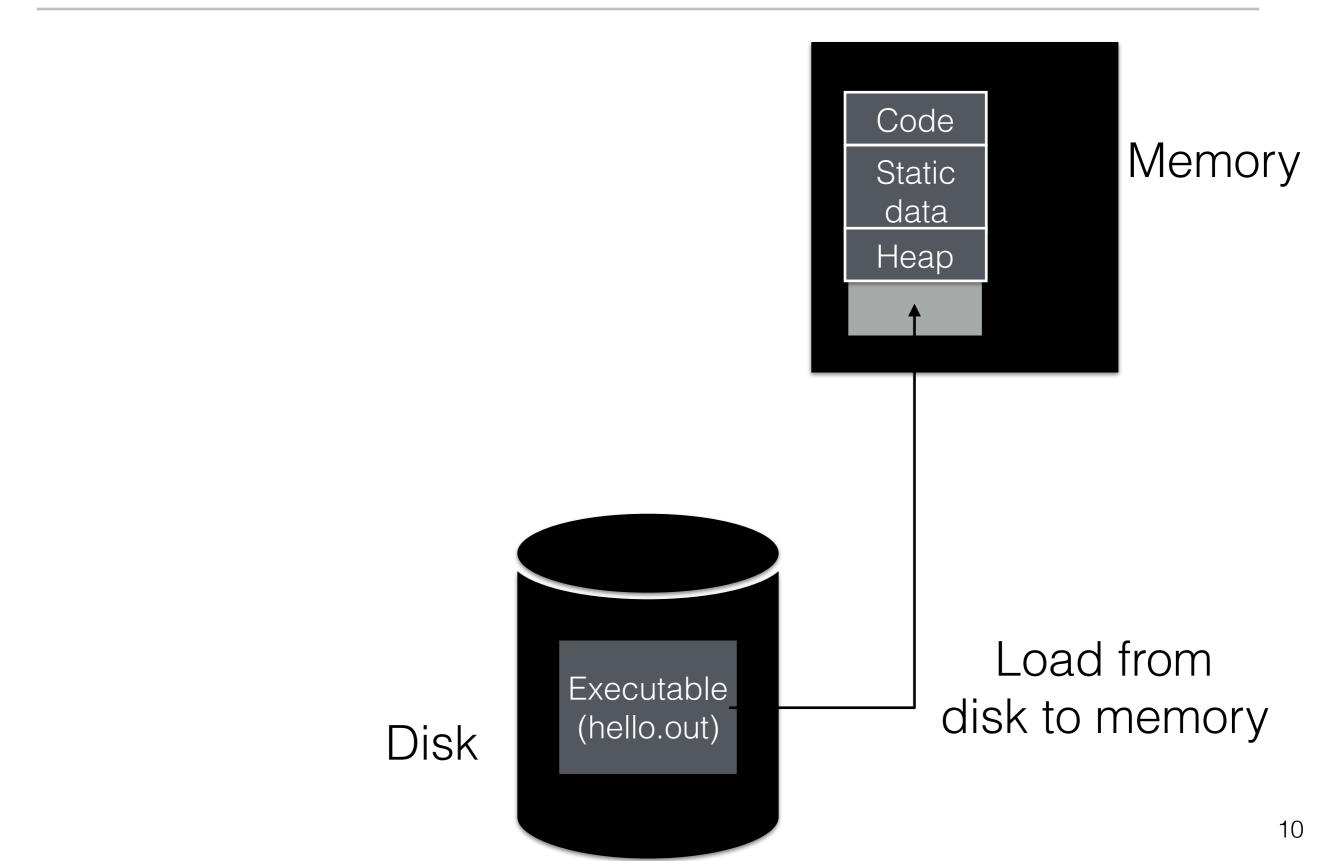


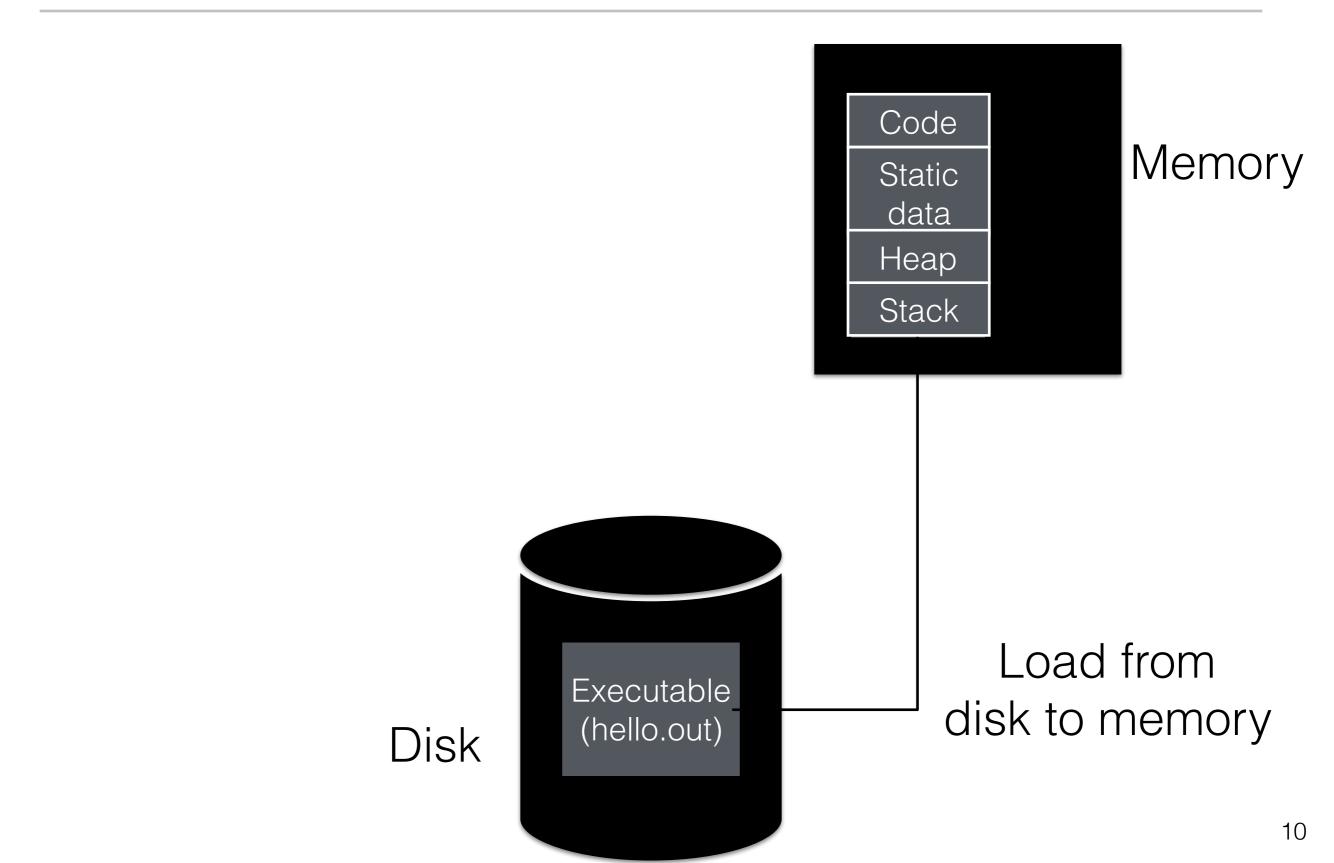


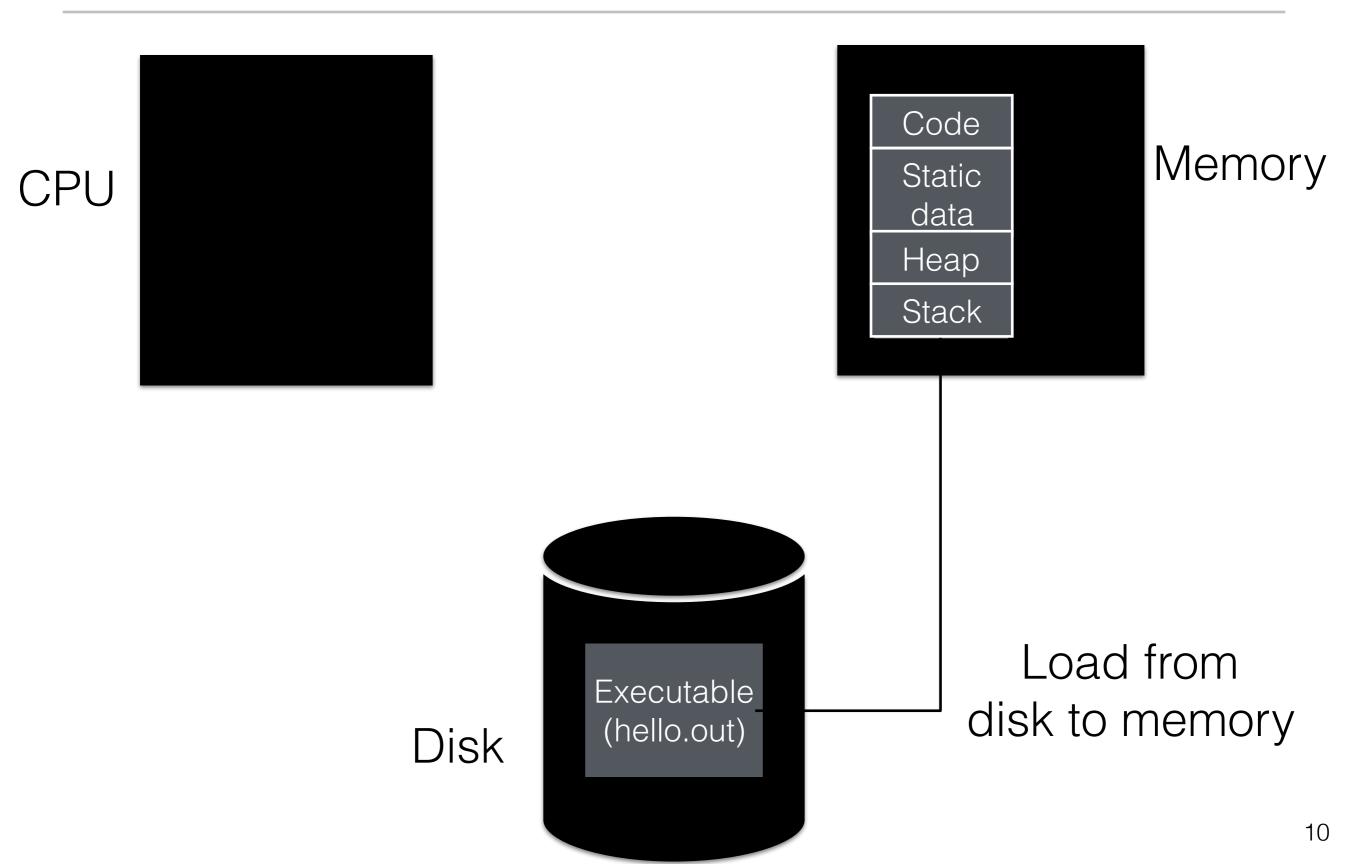


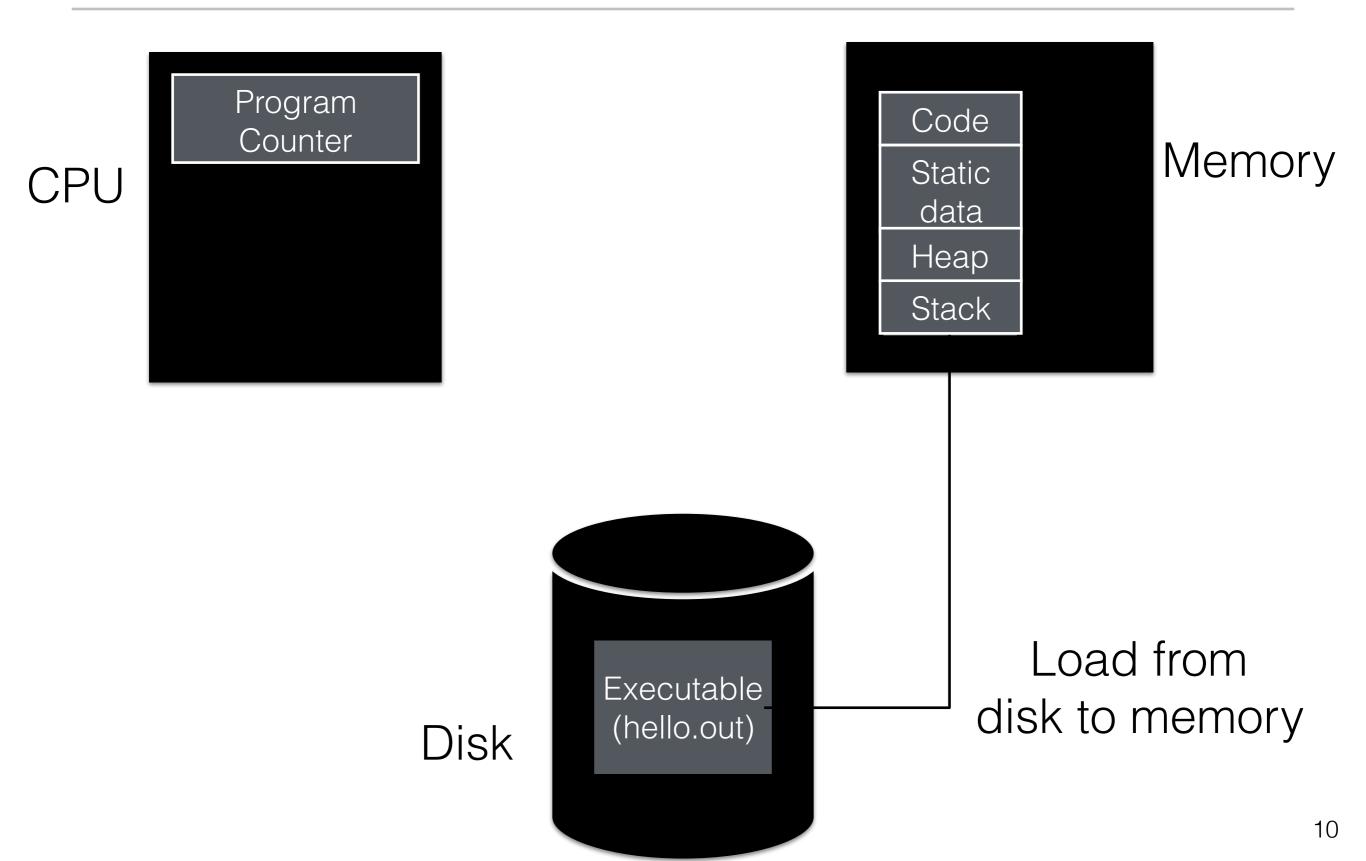


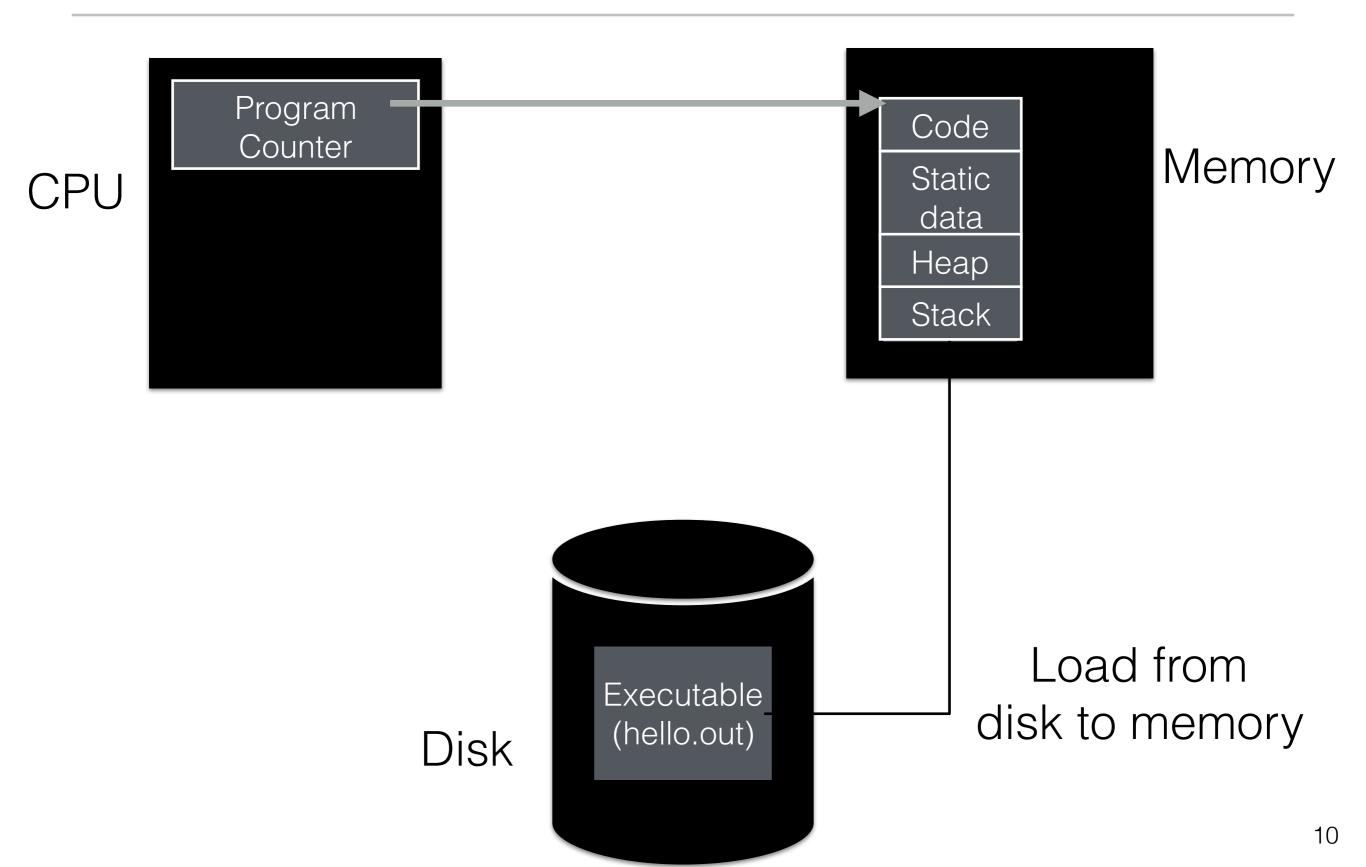


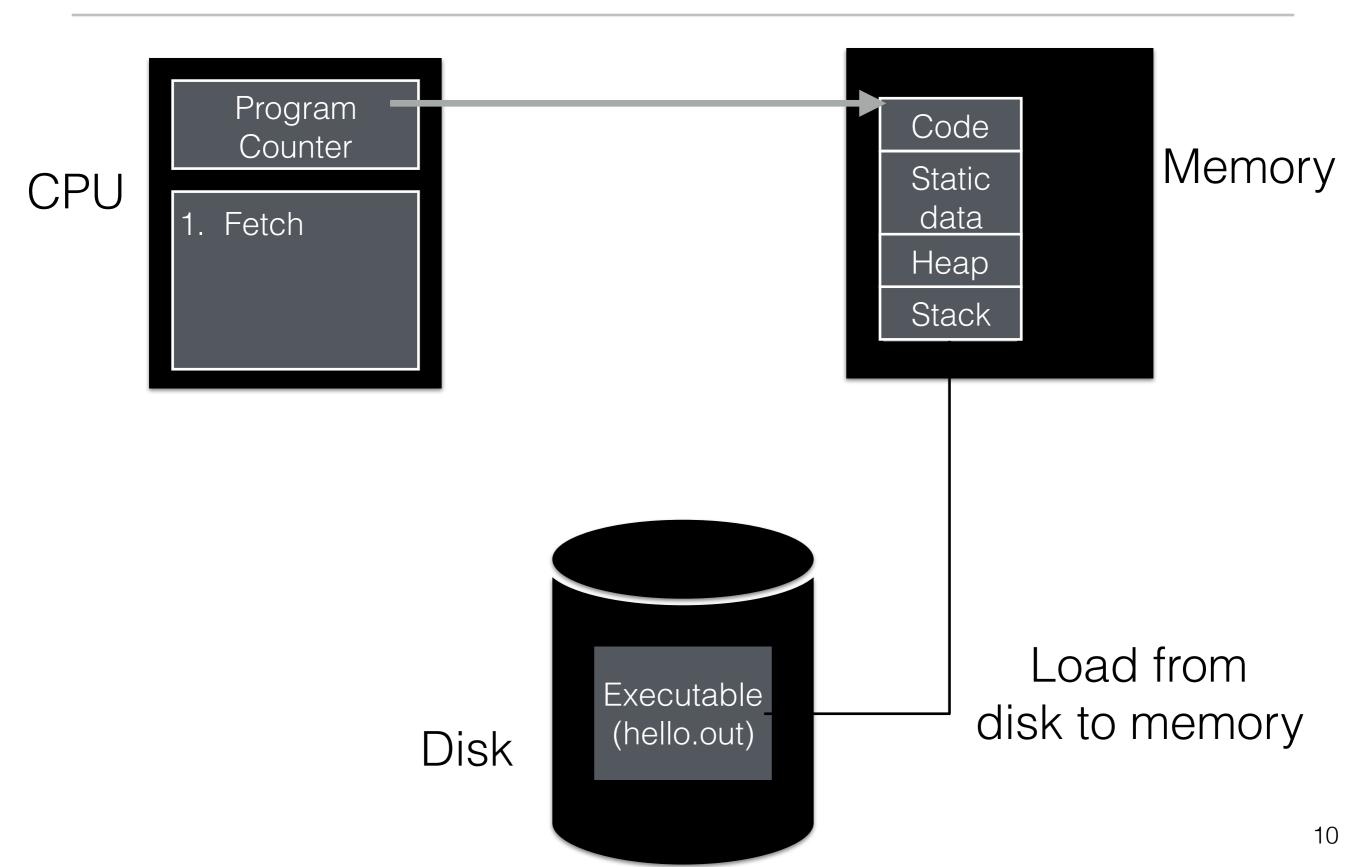


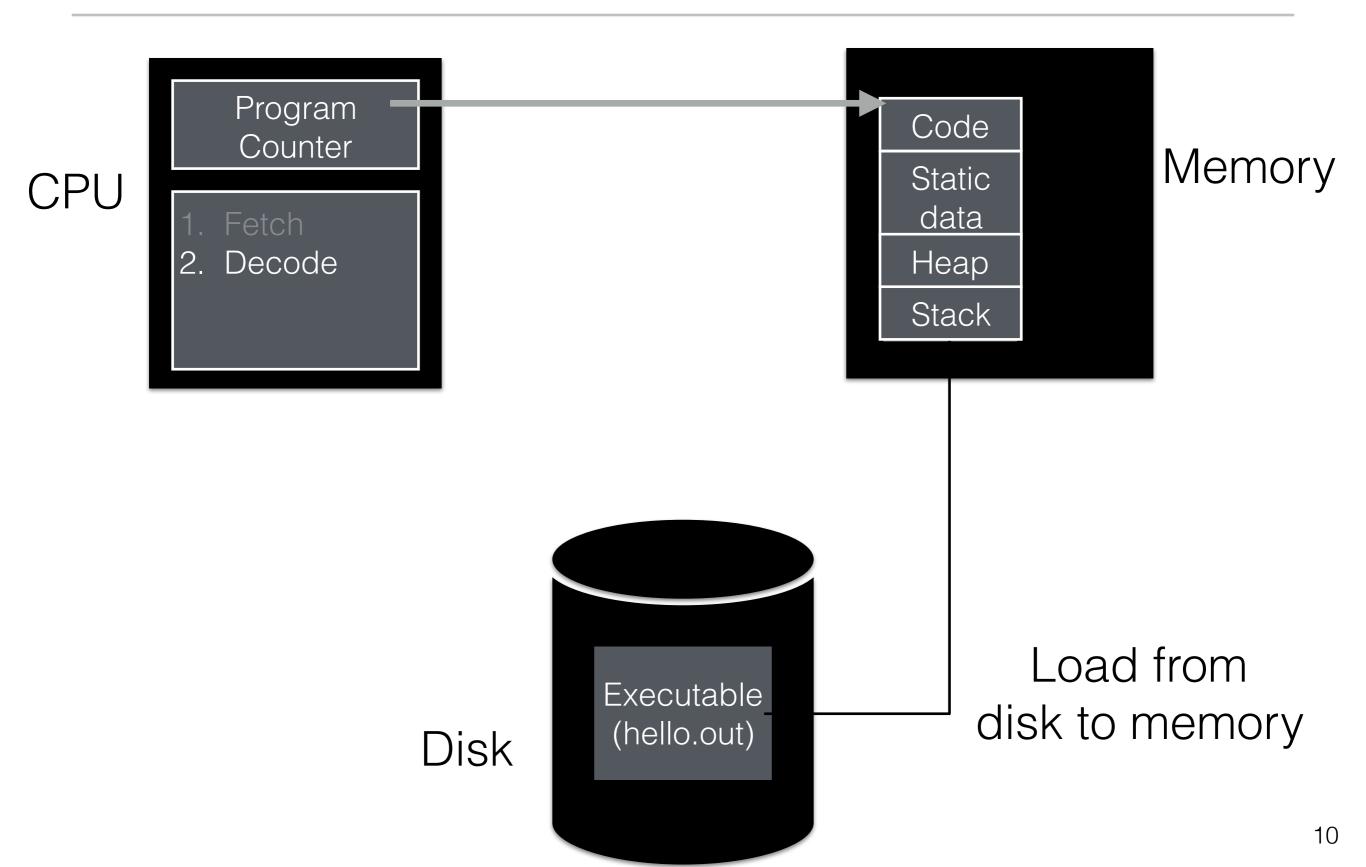


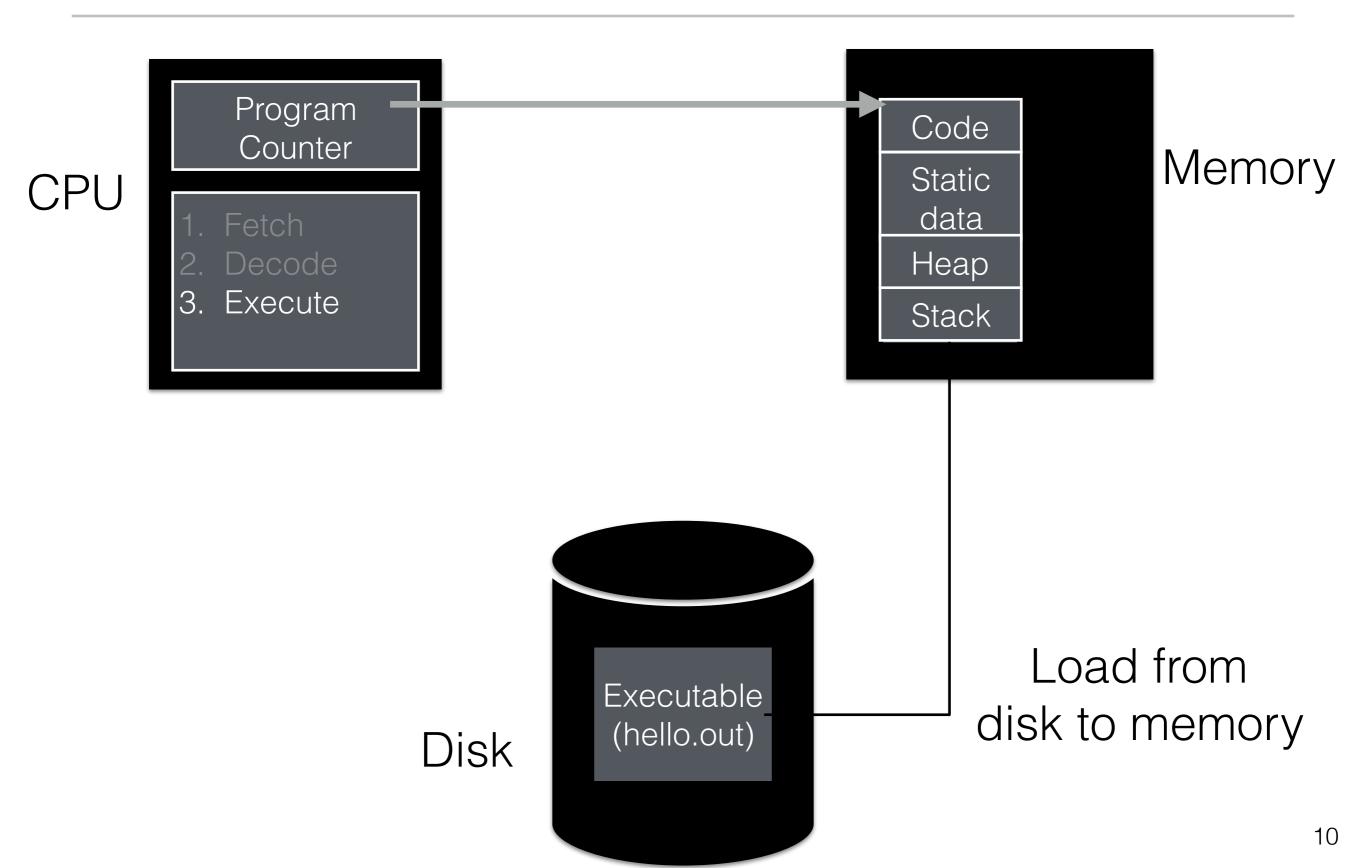


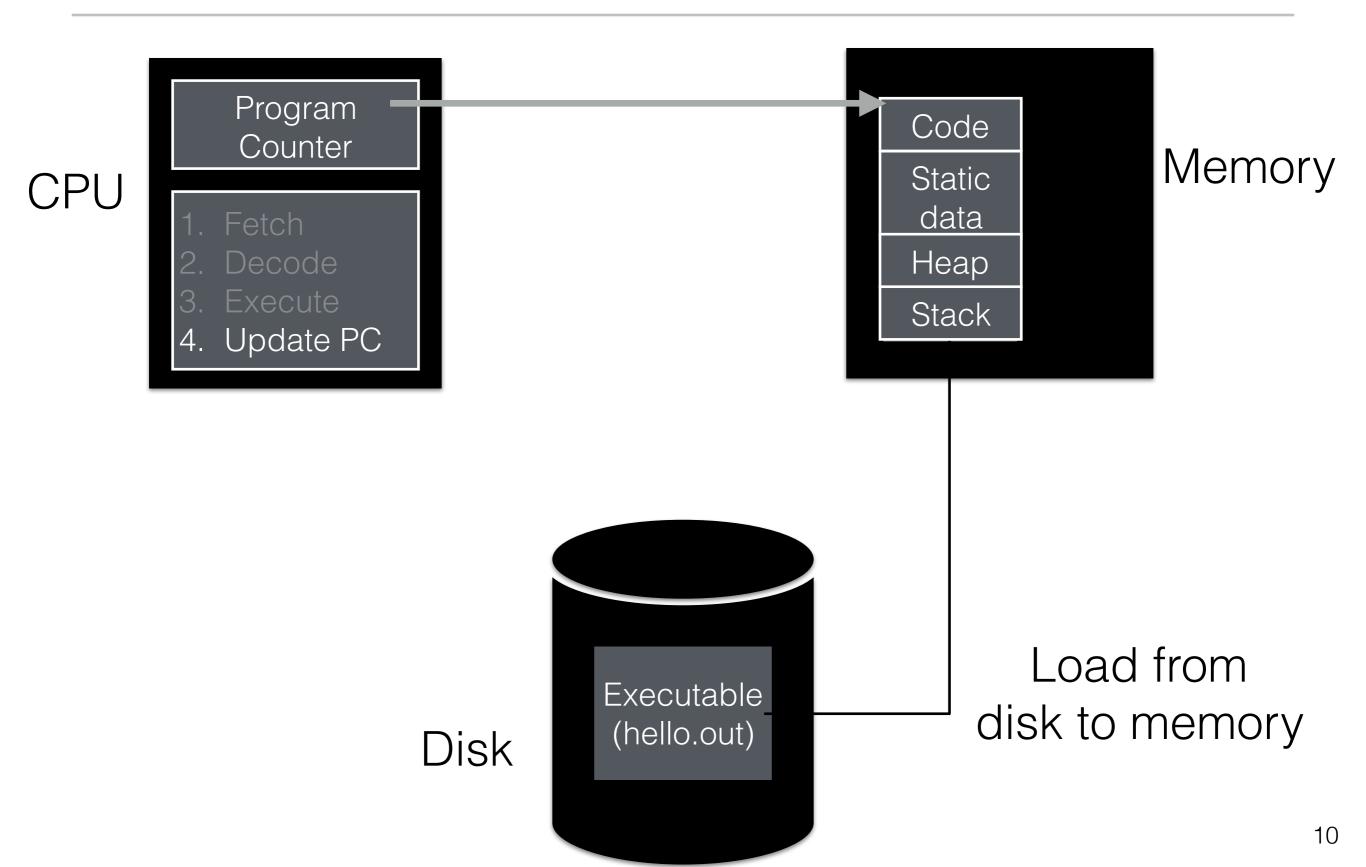


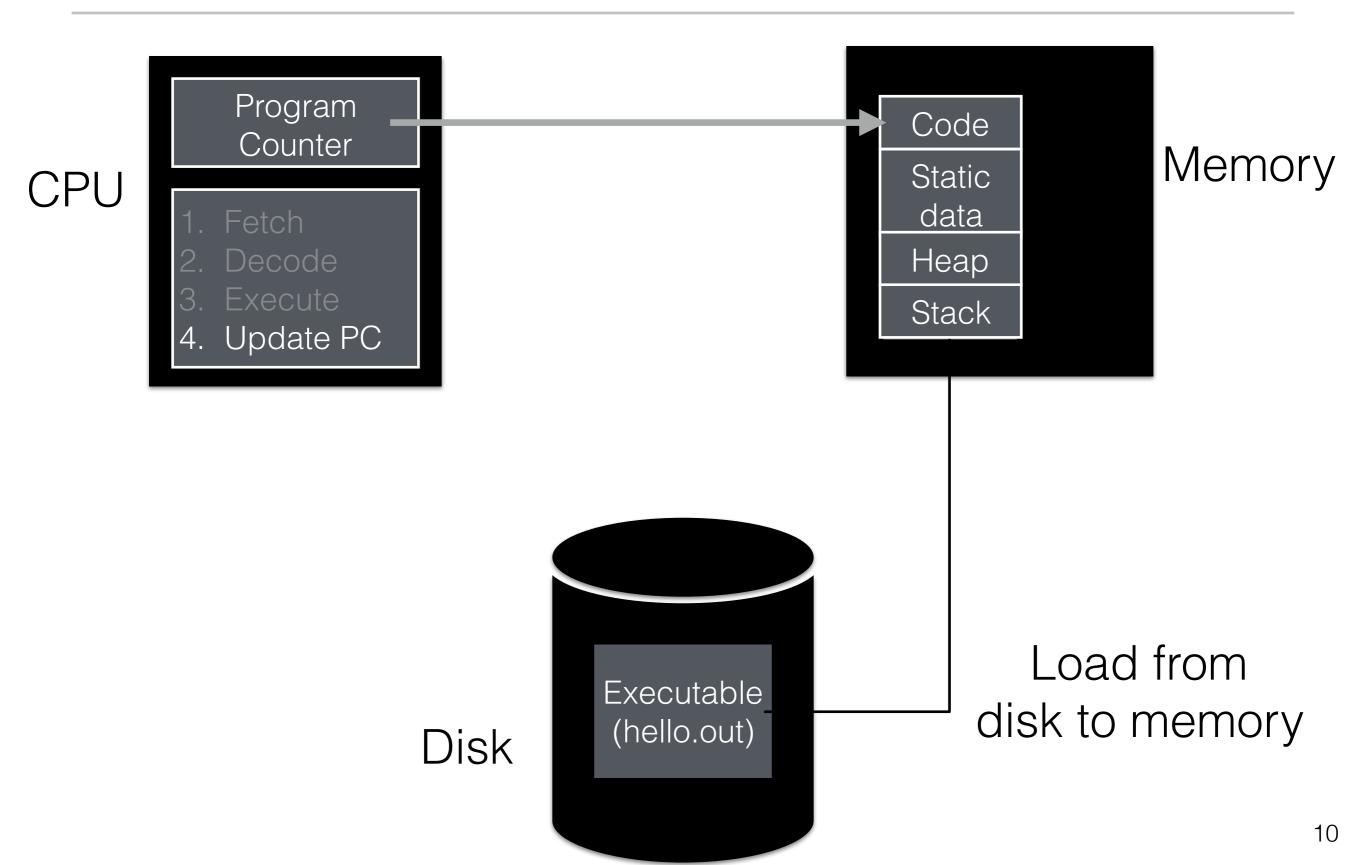












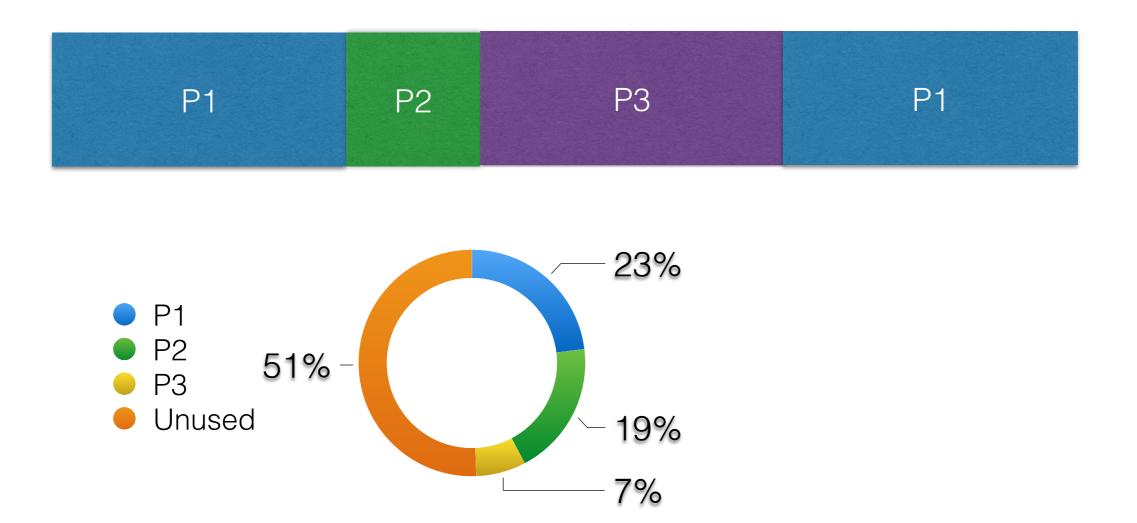
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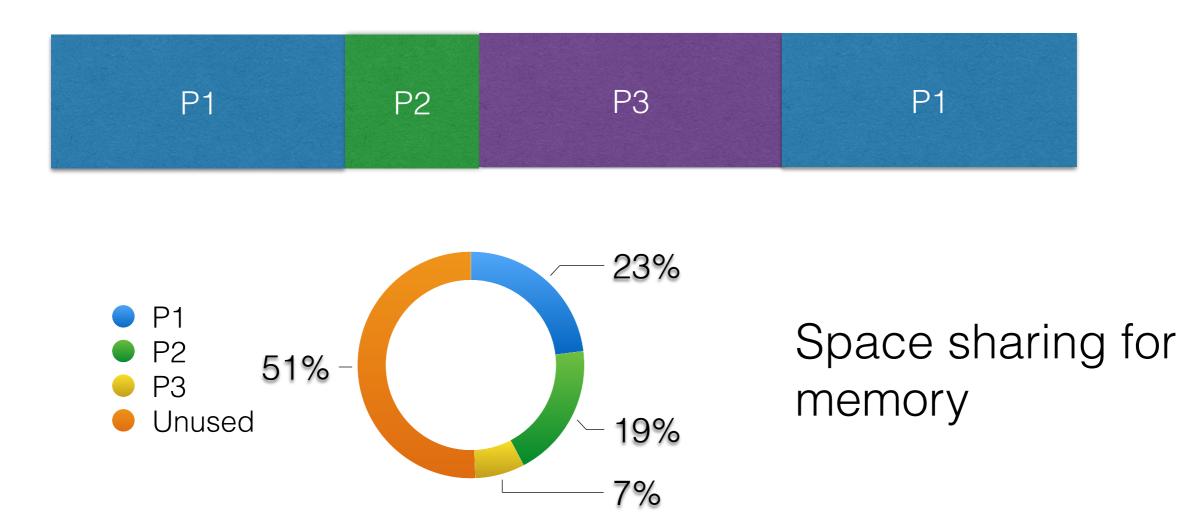
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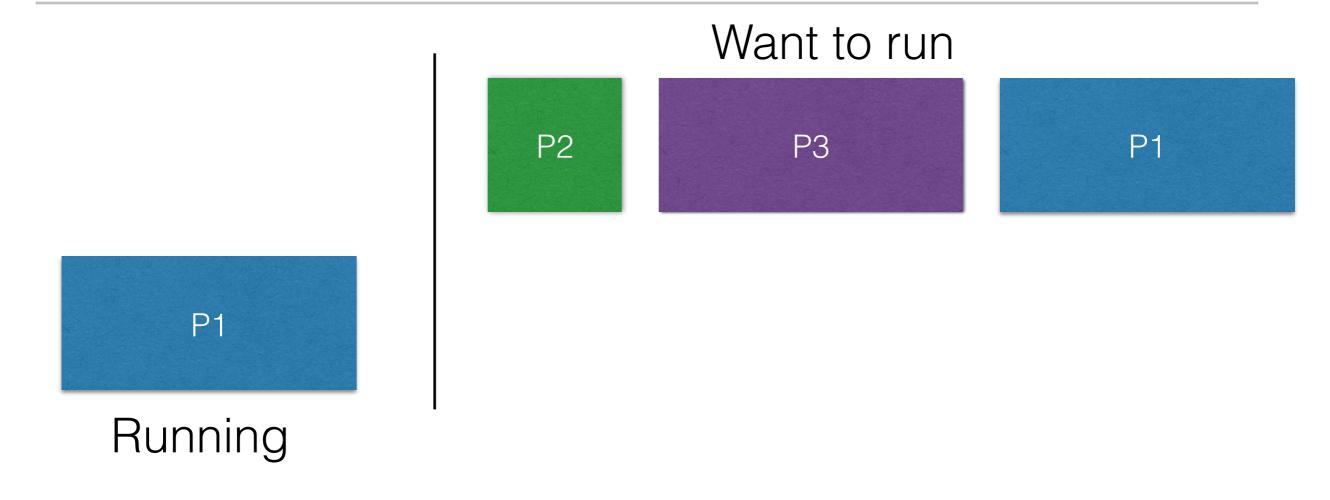


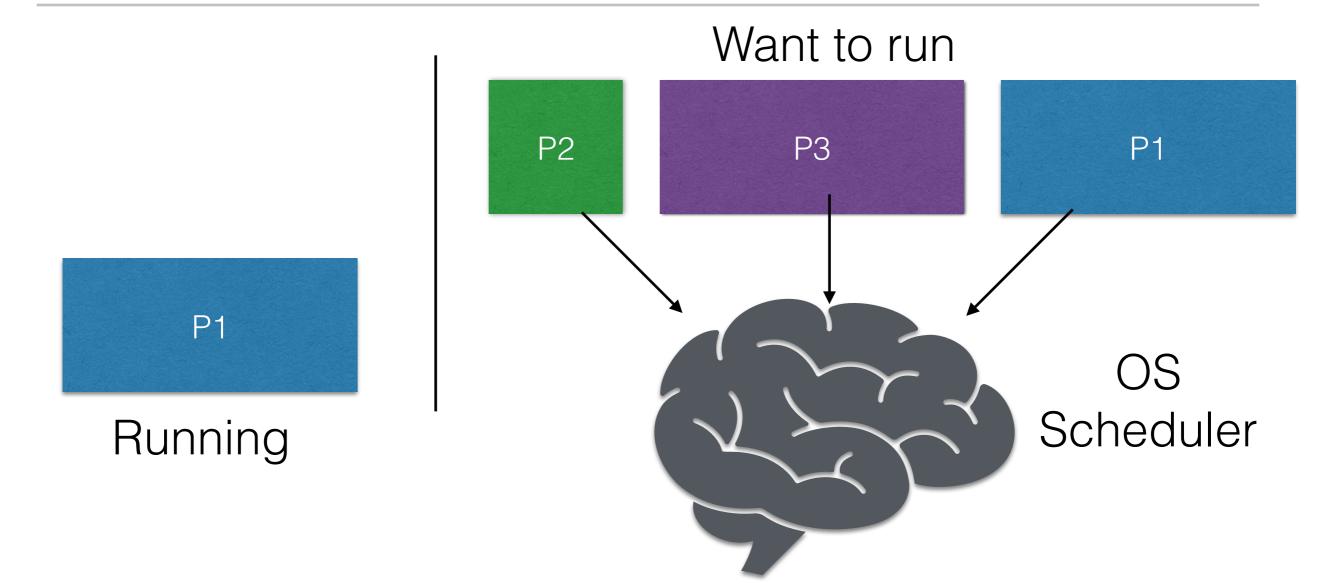
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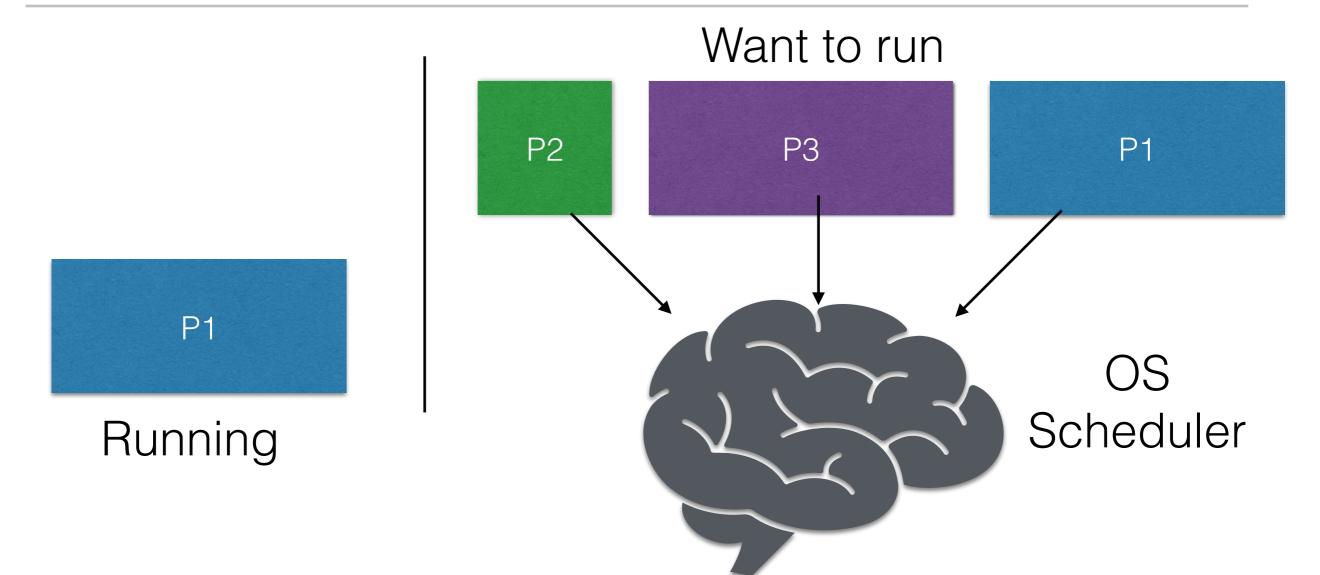




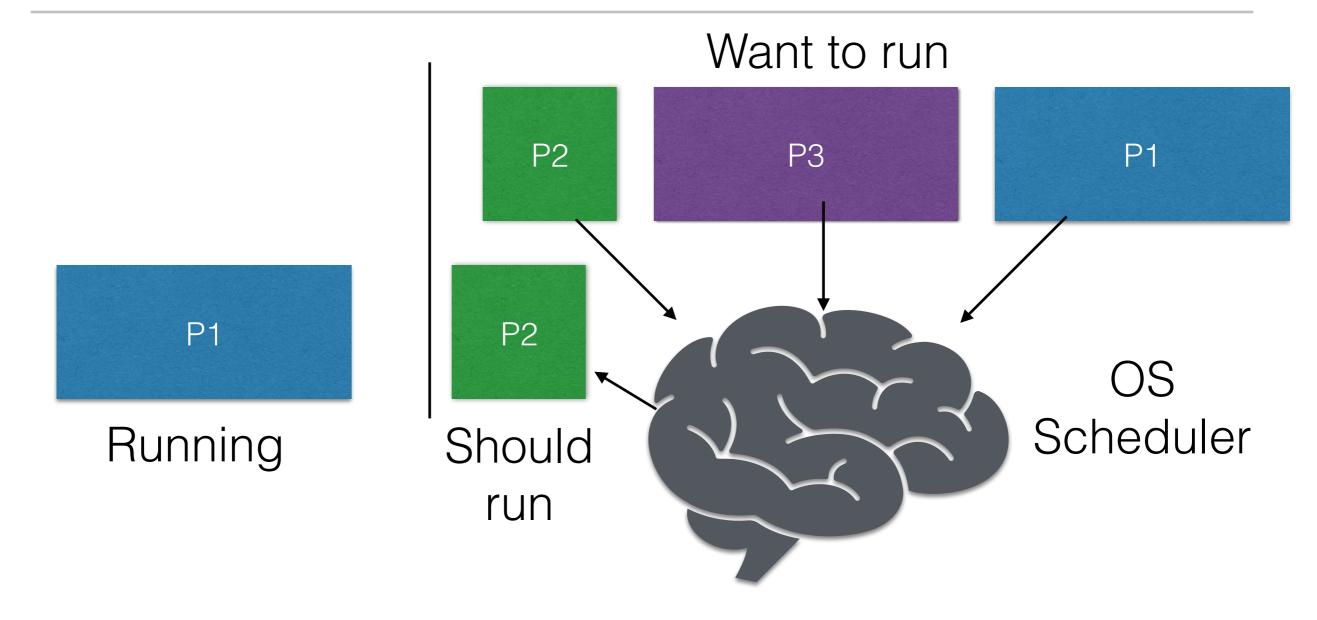
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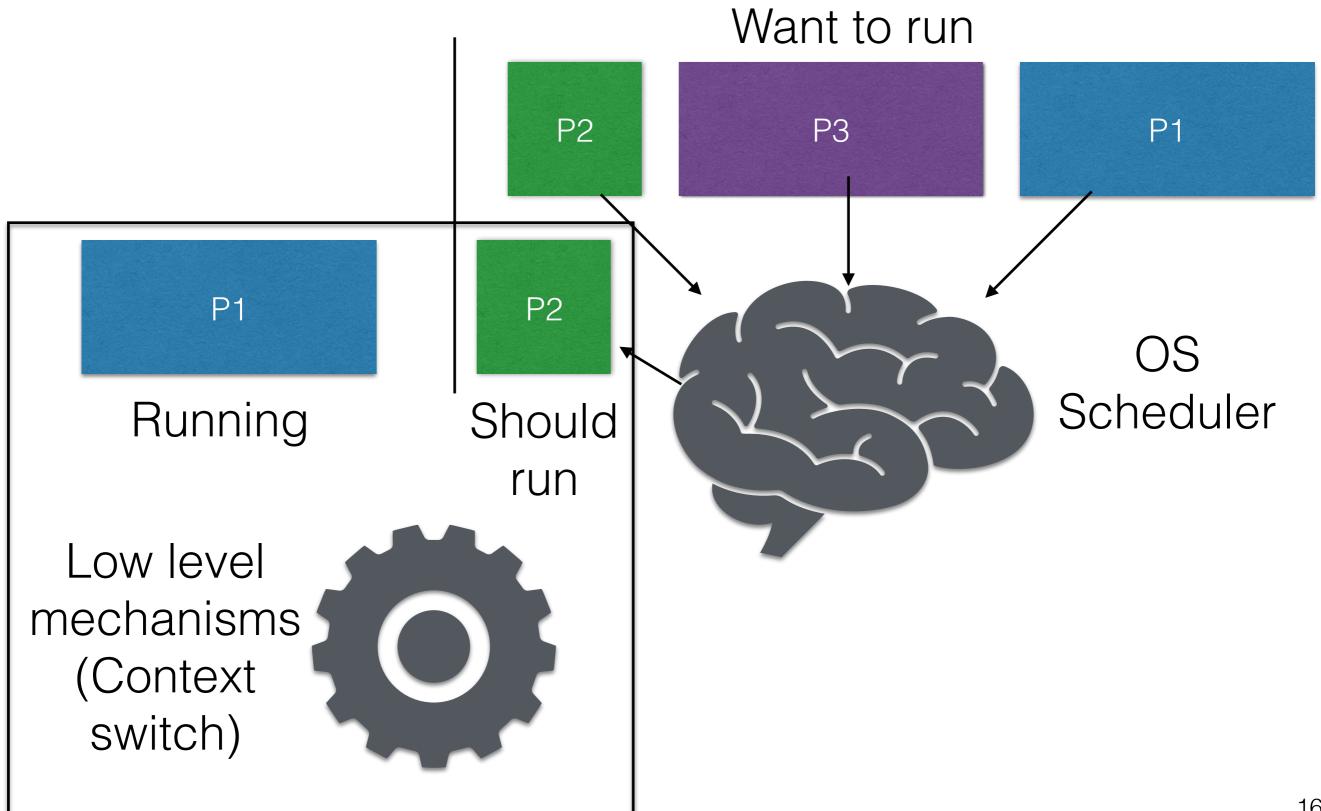


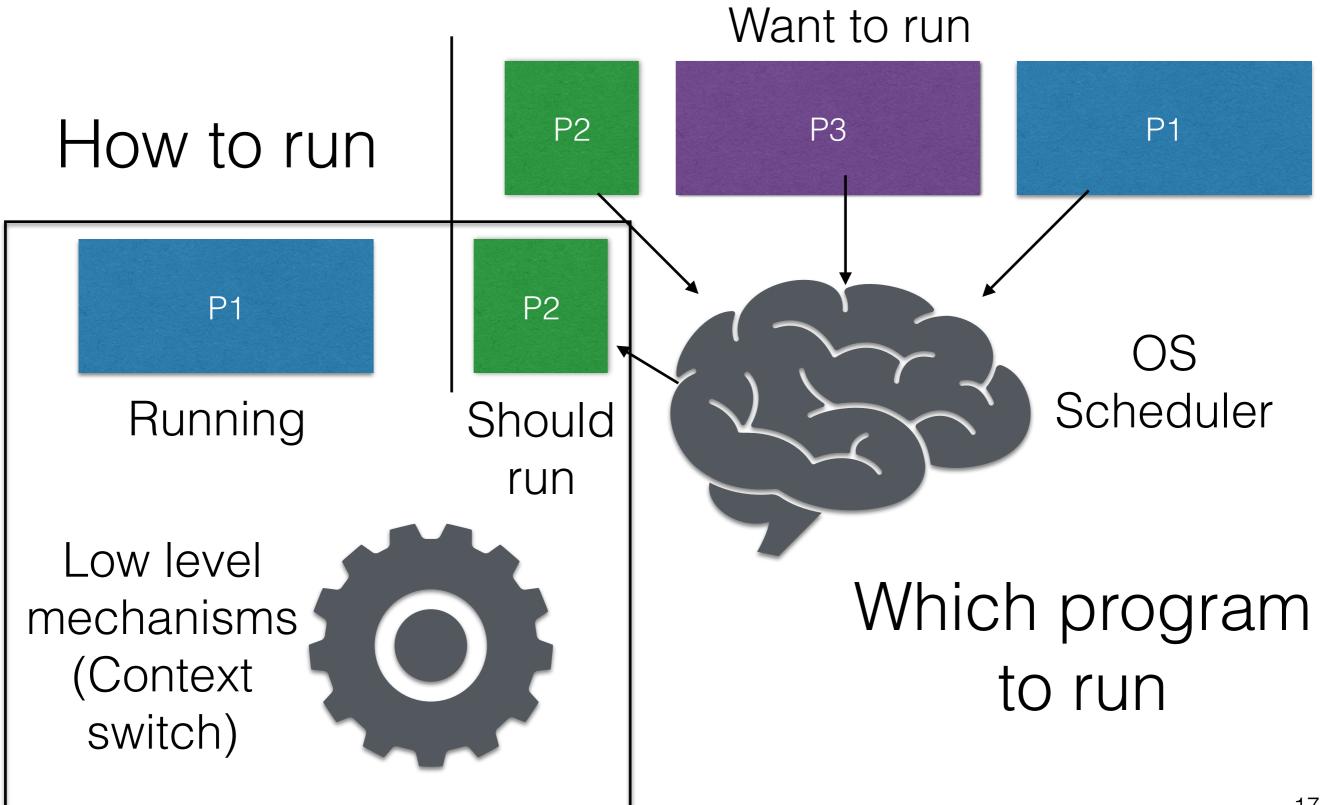




Next P = f(run time, metric, type of process, ...)







Process API

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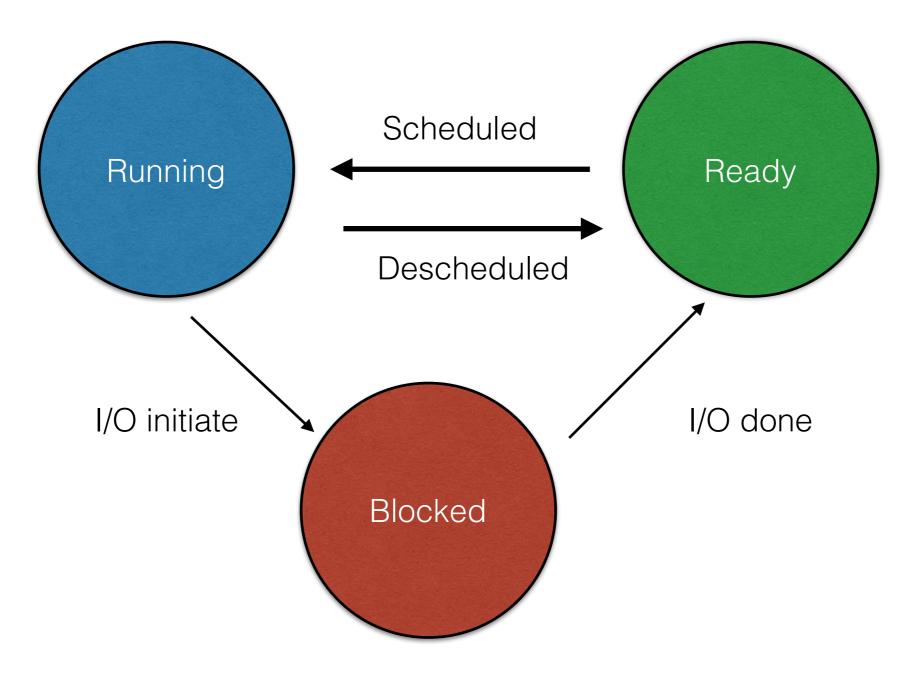
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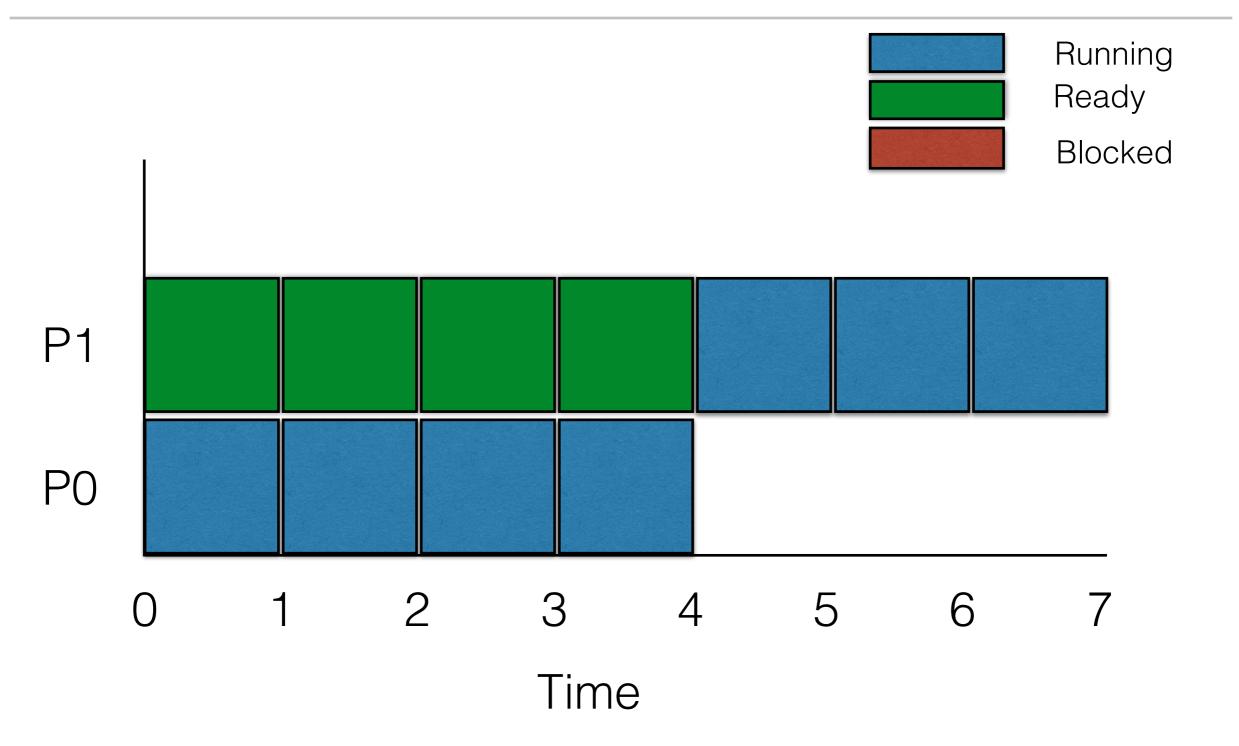
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