

# Operating Systems

## Lecture 12: Paging + TLB

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# CS stories

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<https://www.youtube.com/watch?v=kTn56jJW4zY>

# Revision

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## 1. Segmentation

### 1. Registers containing:

1. Start VA
2. Bounds
3. .... (think Stack)
4. ... (save memory using identical code segment)

2.  $\text{Segment} = (\text{VA} \ \& \ \text{SEG\_MASK}) \gg \text{SEG\_SHIFT}$

3.  $\text{Offset} = \text{VA} \ \& \ \text{OFF\_Mask}$

### 4. Segmentation cons:

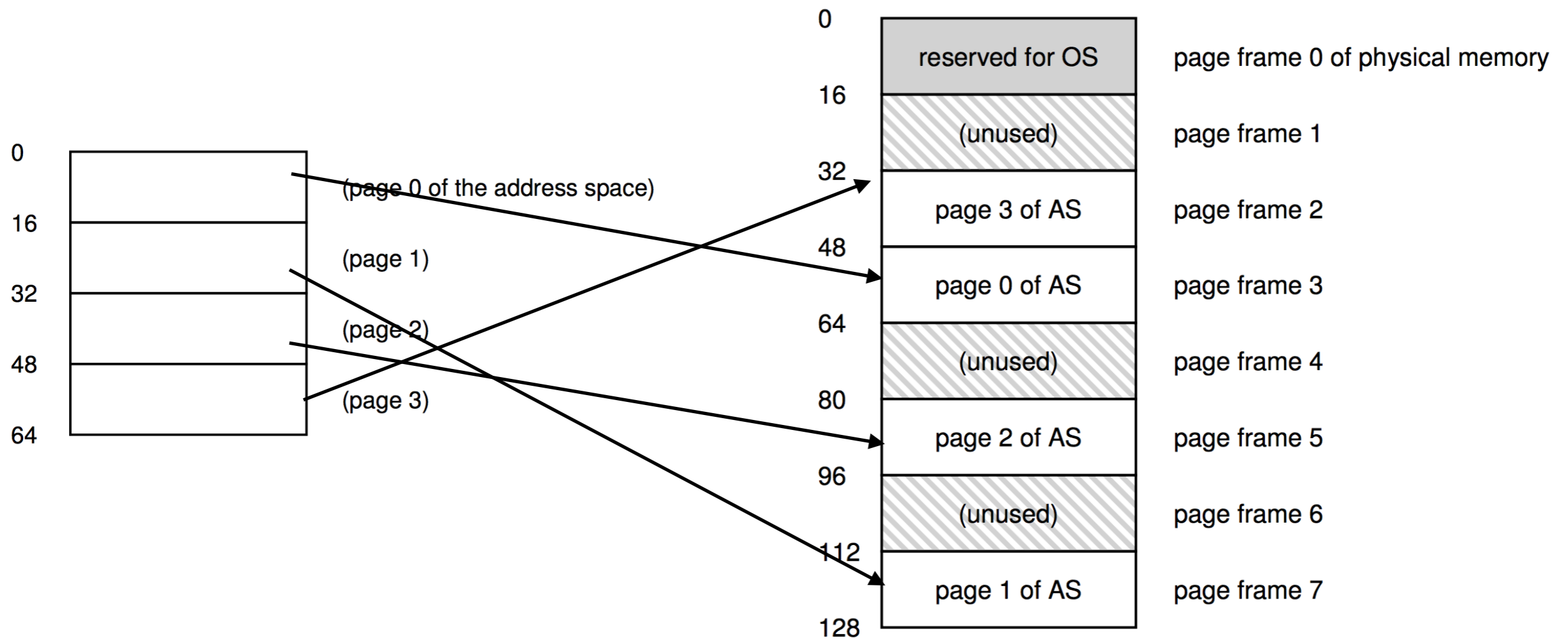
1. Requires \_\_\_\_\_ block of memory for each segment
  1. Can lead to \_\_\_\_\_ and \_\_\_\_\_

# Revision

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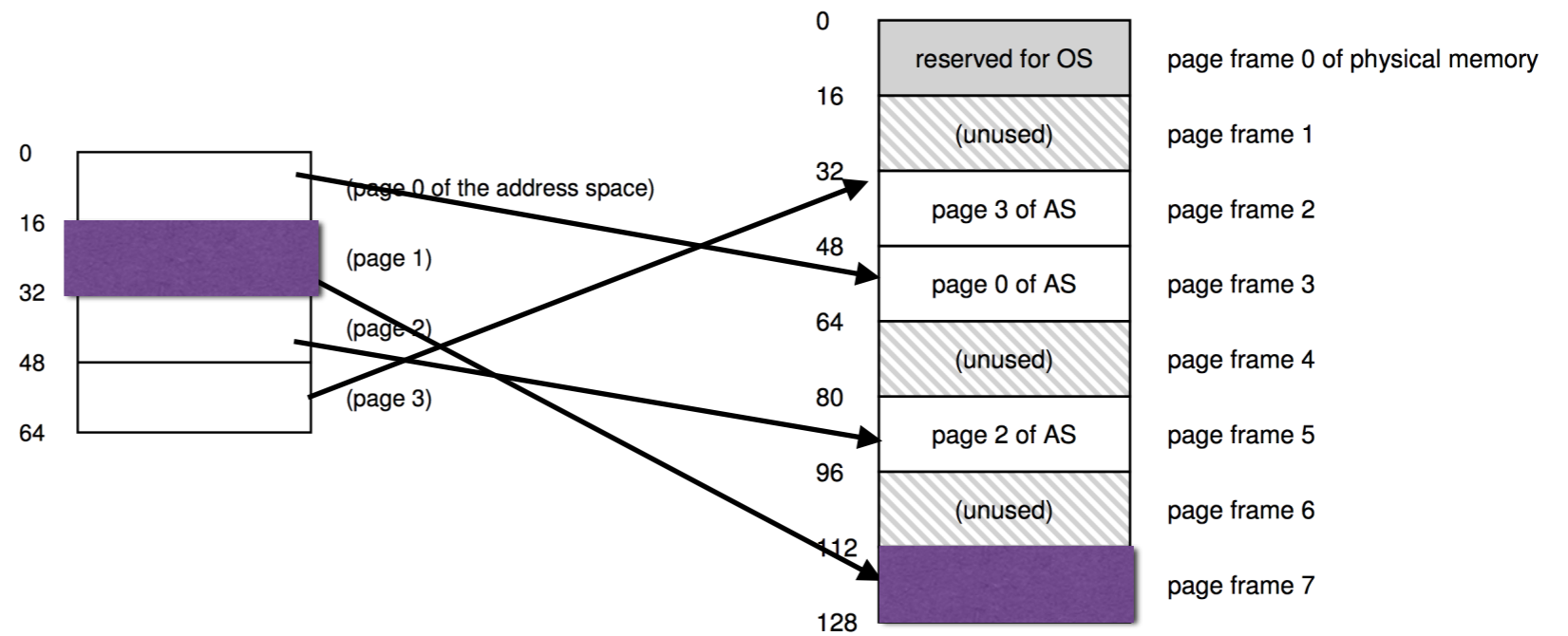
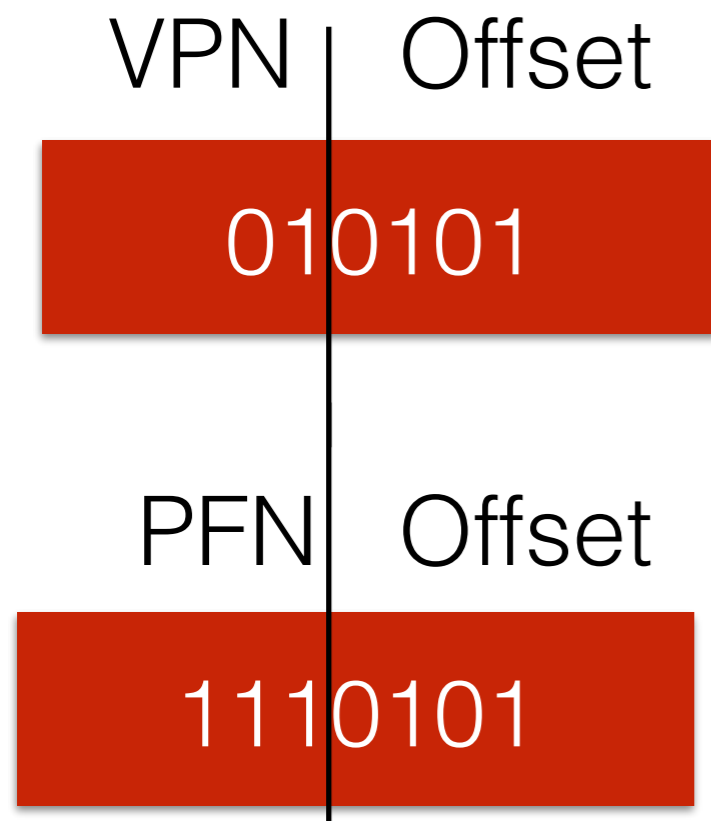
1. Large contiguous memory causes problems
  1. What happens if we map every byte of VA to a byte of PA?
    1. Reduces fragmentation?
      1. External?
      2. Internal?
    2. How much space needed per-process to store mapping?
  2. Middle ground?

# Revision : Paging



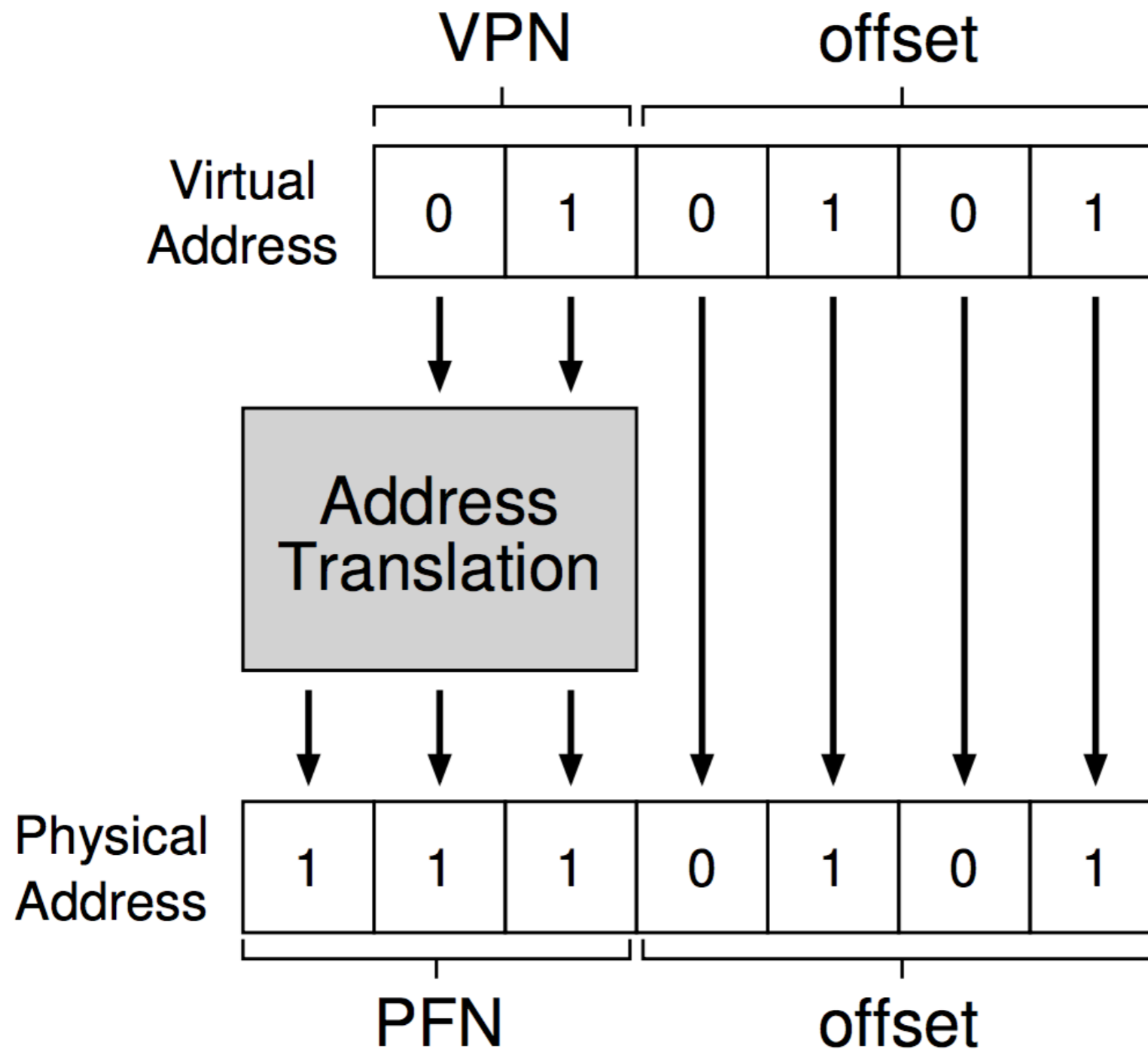
# Example

movl 21, %eax



# Address Translation Summary

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# Page Table Storage

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- Let's consider 32 bit address space
- 32 bit address space with 4 KB pages
- 4 KB pages -> \_\_\_\_\_ bits?
  - 12 bits Offset
- Remaining bits =  $32 - 12 = 20$ 
  - 20 bit VPN
  - # pages =  $2^{20}$
  - # translations required = \_\_\_\_\_
    - $2^{20}$
- 4 bytes per translation ->  $4 * 2^{20} \text{ MB} = 4 \text{ MB/}$   
process



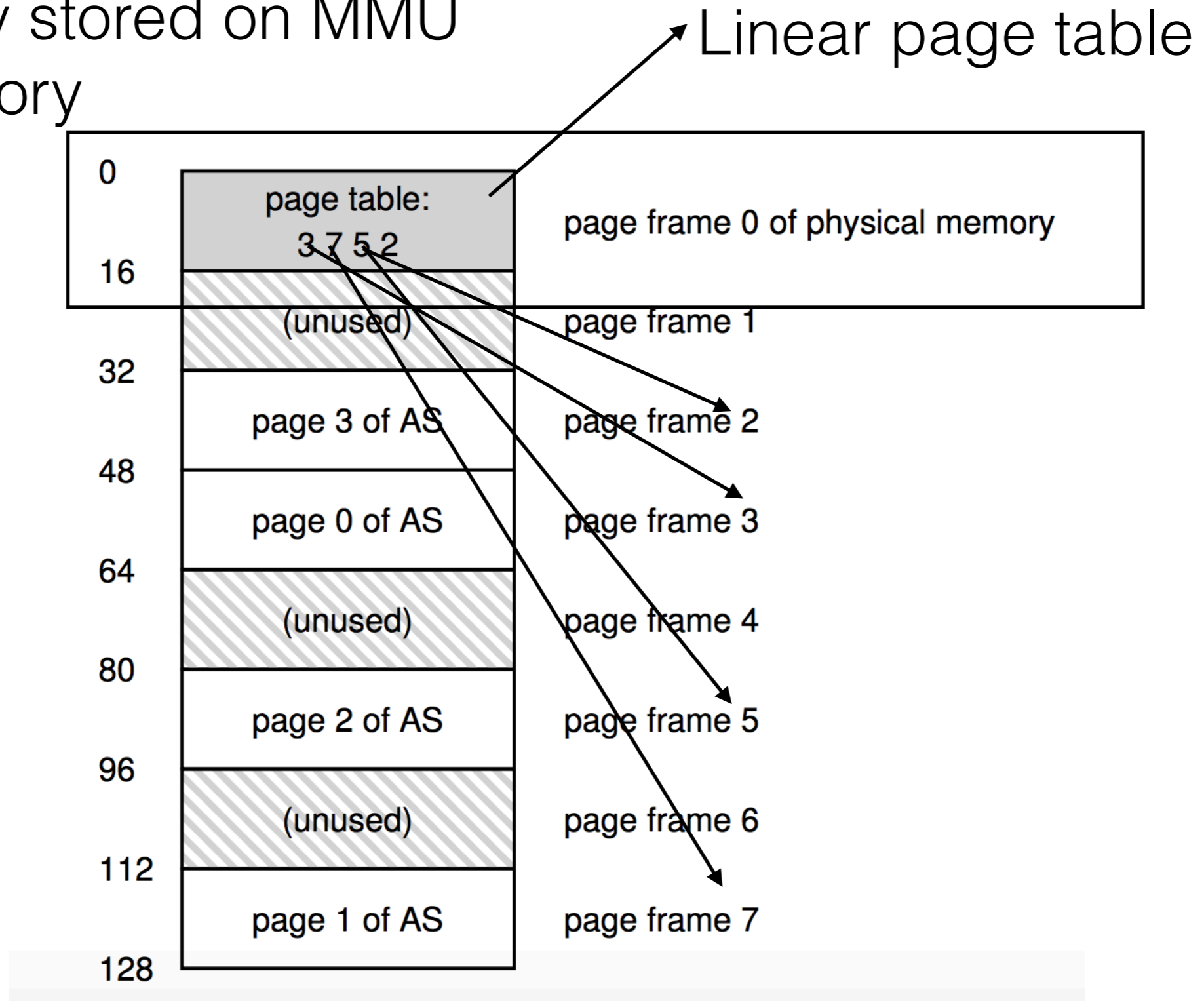
# Page Size Tradeoffs?

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- Small size
  - More # of translations
    - More memory overhead/process
    - Less chances of fragmentation
- Large size
  - Less # of translations
    - Less memory overhead/process
    - More chances of fragmentation

# Page Table Storage

Not really stored on MMU  
- In memory



# What else is in the Page Table?

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- Protection bit : Read/Write/Execute?
- Present bit: On Memory or HDD/SSD?
- Reference bit: Is the page popular/being referenced?
  - Else?
- Valid bit: Is translation valid?
- Dirty bit: Modified since brought to memory?

# Worked Out Example

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```
int array[1000];  
...  
for (i = 0; i < 1000; i++)  
    array[i] = 0;
```

```
1024 movl $0x0, (%edi,%eax,4)  
1028 incl %eax  
1032 cmpl $0x03e8,%eax  
1036 jne 0x1024
```

←  $*(EDI + 4 * EAX) = 0$

# Worked Out Example

---

```
int array[1000];  
...  
for (i = 0; i < 1000; i++)  
    array[i] = 0;
```

Address of array[0]

```
1024 movl $0x0, (%edi,%eax,4)  
1028 incl %eax  
1032 cmpl $0x03e8,%eax  
1036 jne 0x1024
```

$*(EDI + 4 * EAX) = 0$

# Worked Out Example

---

```
int array[1000];  
...  
for (i = 0; i < 1000; i++)  
    array[i] = 0;
```

Index into array (i)

```
1024 movl $0x0, (%edi,%eax,4)  
1028 incl %eax  
1032 cmpl $0x03e8,%eax  
1036 jne 0x1024
```

$*(EDI + 4 * EAX) = 0$

# Worked Out Example

---

```
int array[1000];  
...  
for (i = 0; i < 1000; i++)  
    array[i] = 0;
```

Index into array (i)

```
1024 movl $0x0, (%edi,%eax,4)  
1028 incl %eax  
1032 cmpl $0x03e8,%eax  
1036 jne 0x1024
```

$*(EDI + 4 * EAX) = 0$   
 $i = i + 1$

# Worked Out Example

---

```
int array[1000];  
...  
for (i = 0; i < 1000; i++)  
    array[i] = 0;
```

Index into array (i)

```
1024 movl $0x0, (%edi,%eax,4)  
1028 incl %eax  
1032 cmpl $0x03e8,%eax  
1036 jne 0x1024
```

↓  
 $*(EDI + 4 * EAX) = 0$   
 $i = i + 1$   
 $!s i == 1000$



# Worked Out Example

---

```
int array[1000];  
...  
for (i = 0; i < 1000; i++)  
    array[i] = 0;
```

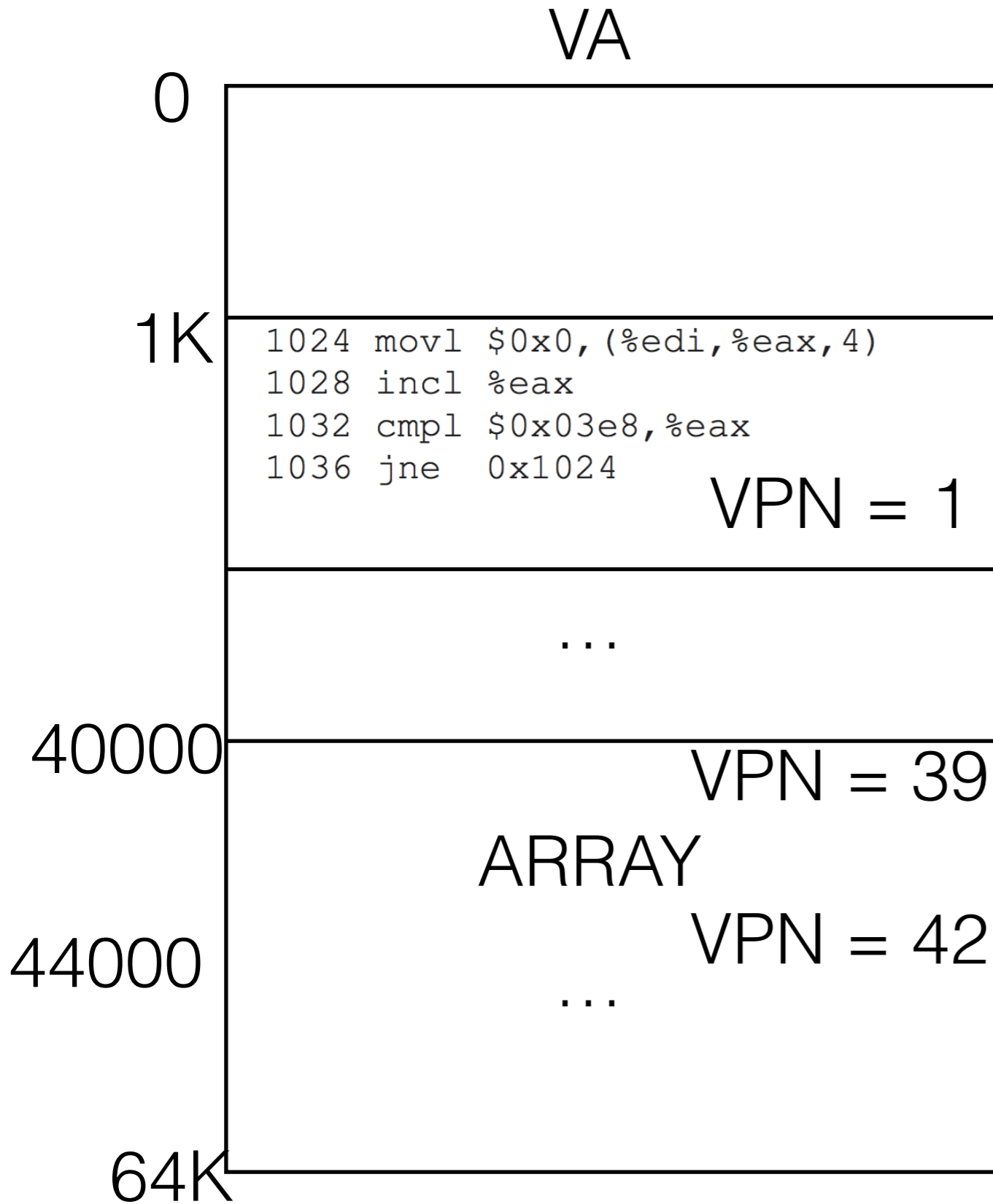
Index into array (i)

```
1024 movl $0x0, (%edi,%eax,4)  
1028 incl %eax  
1032 cmpl $0x03e8,%eax  
1036 jne 0x1024
```

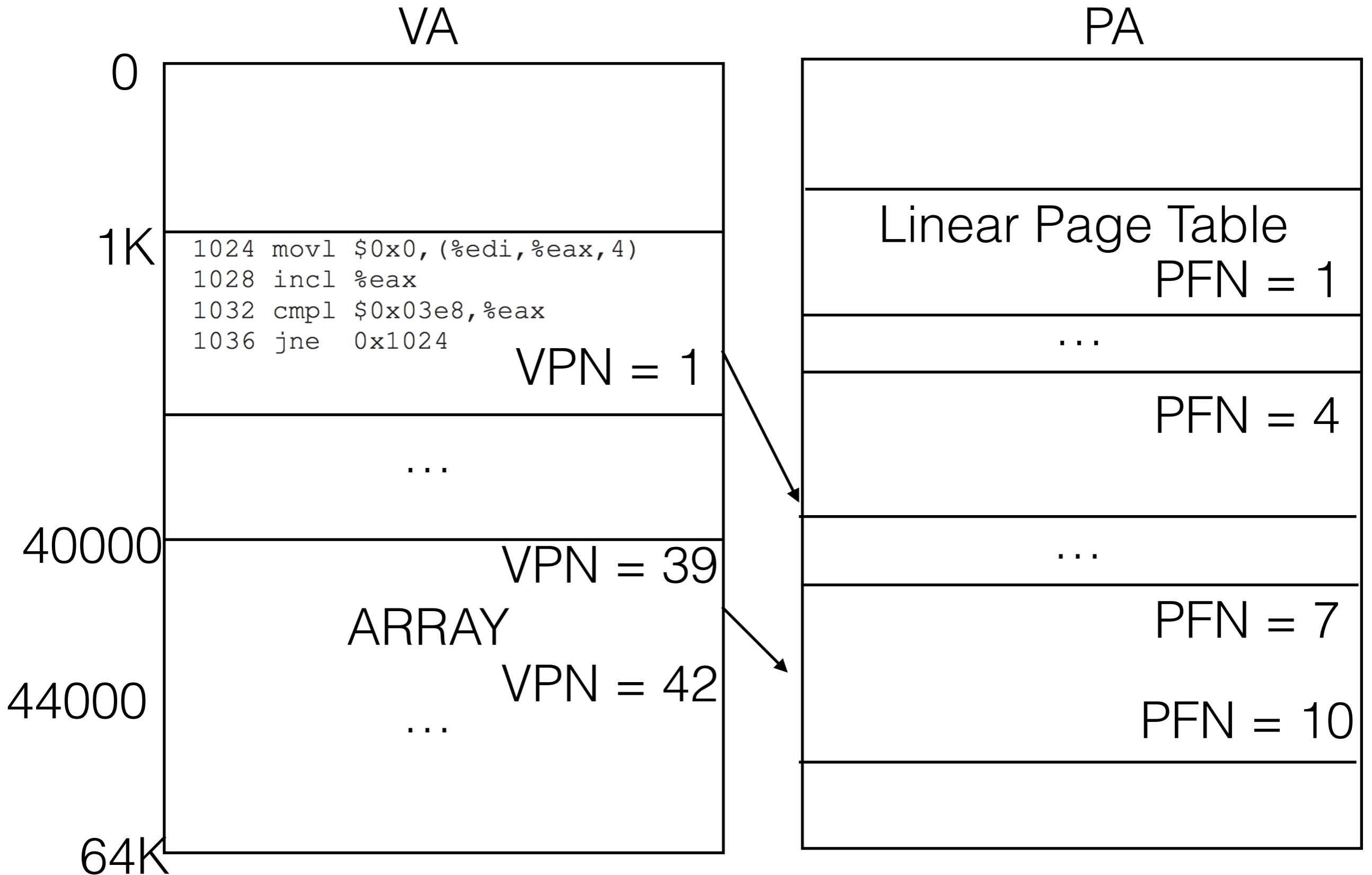
↓  
 $*(EDI + 4 * EAX) = 0$   
 $I = I + 1$   
Is  $I == 1000$   
If Above is False

# Worked Out Example

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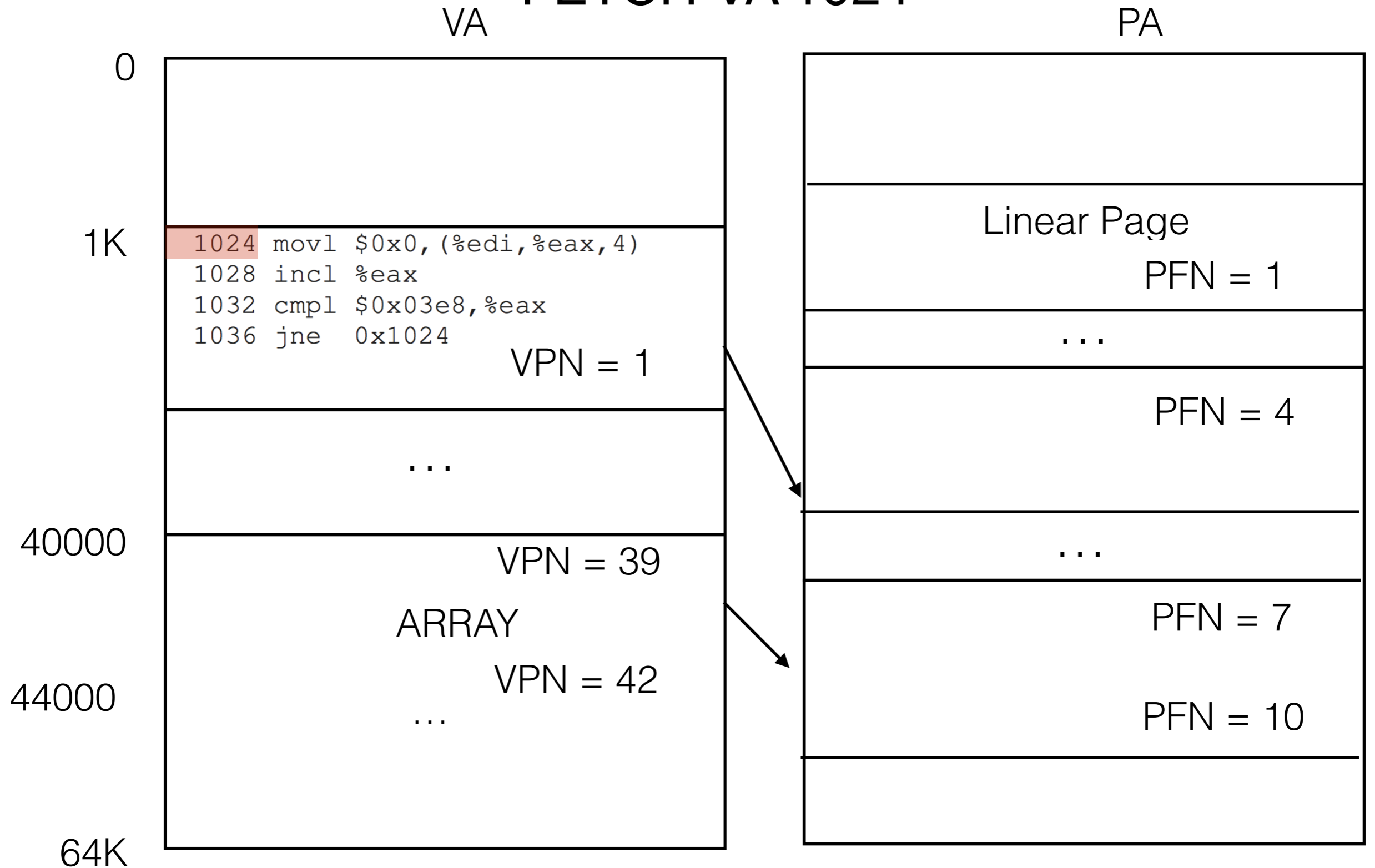


# Worked Out Example



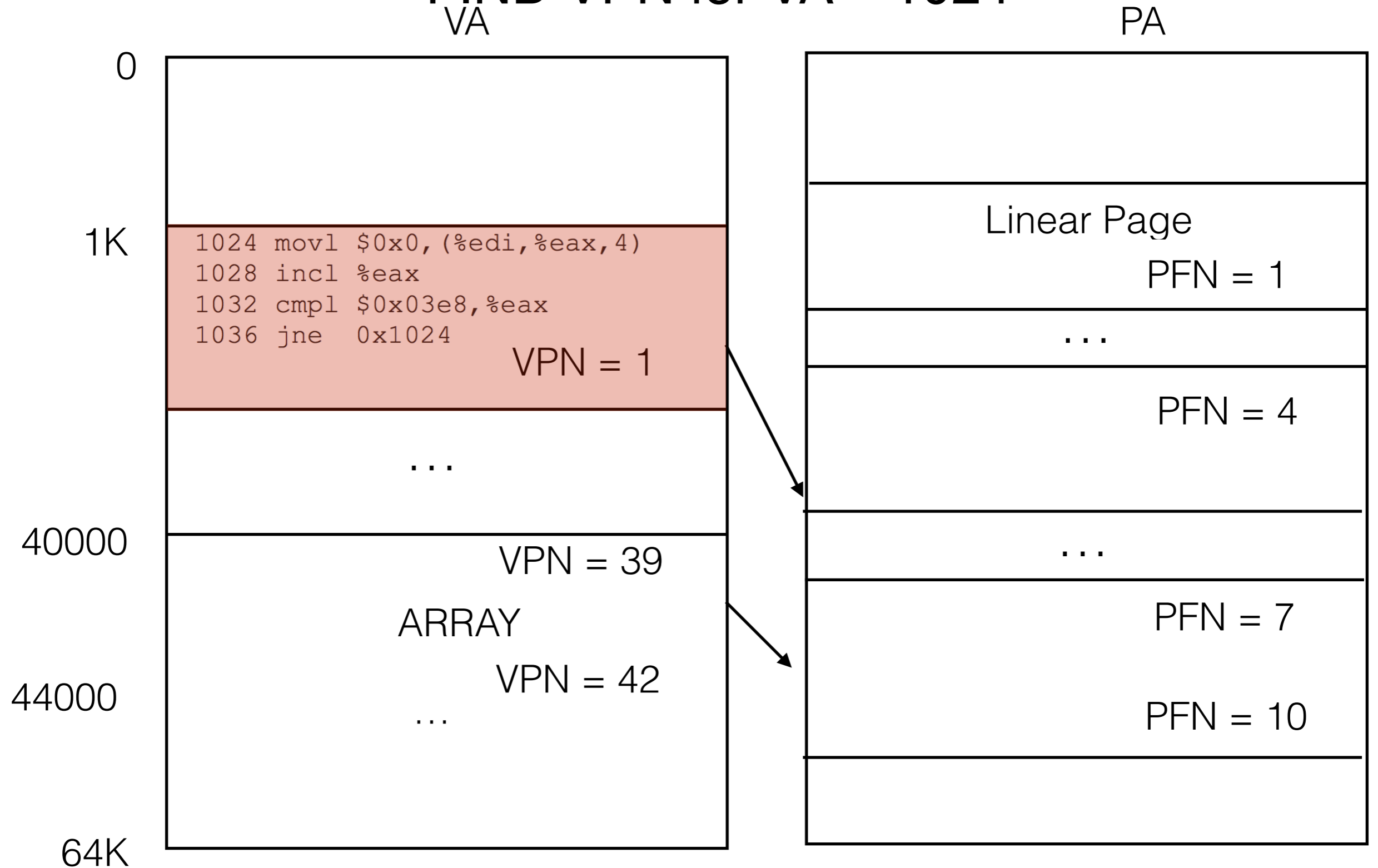
# Worked Out Example

## FETCH VA 1024



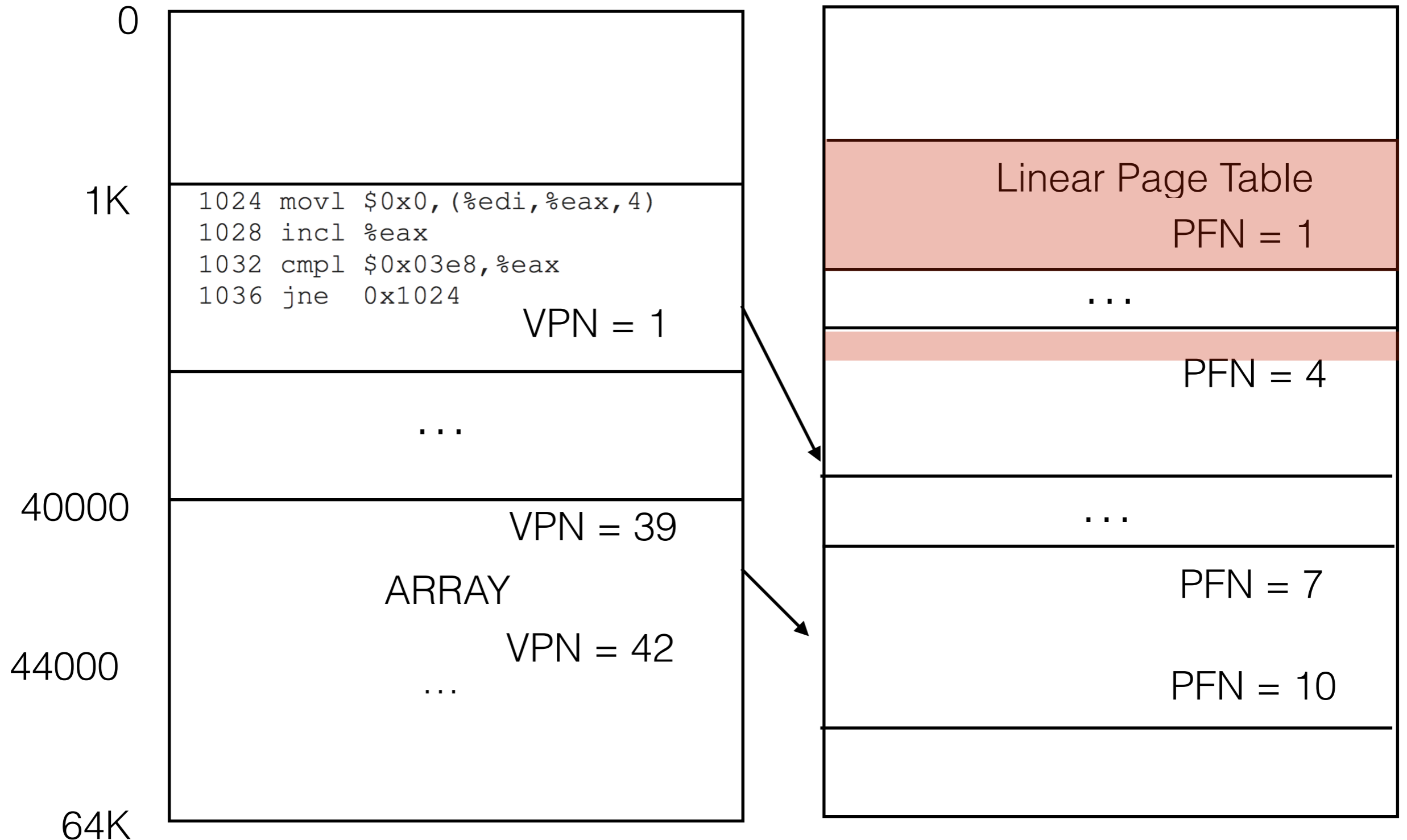
# Worked Out Example

## FIND VPN for VA = 1024



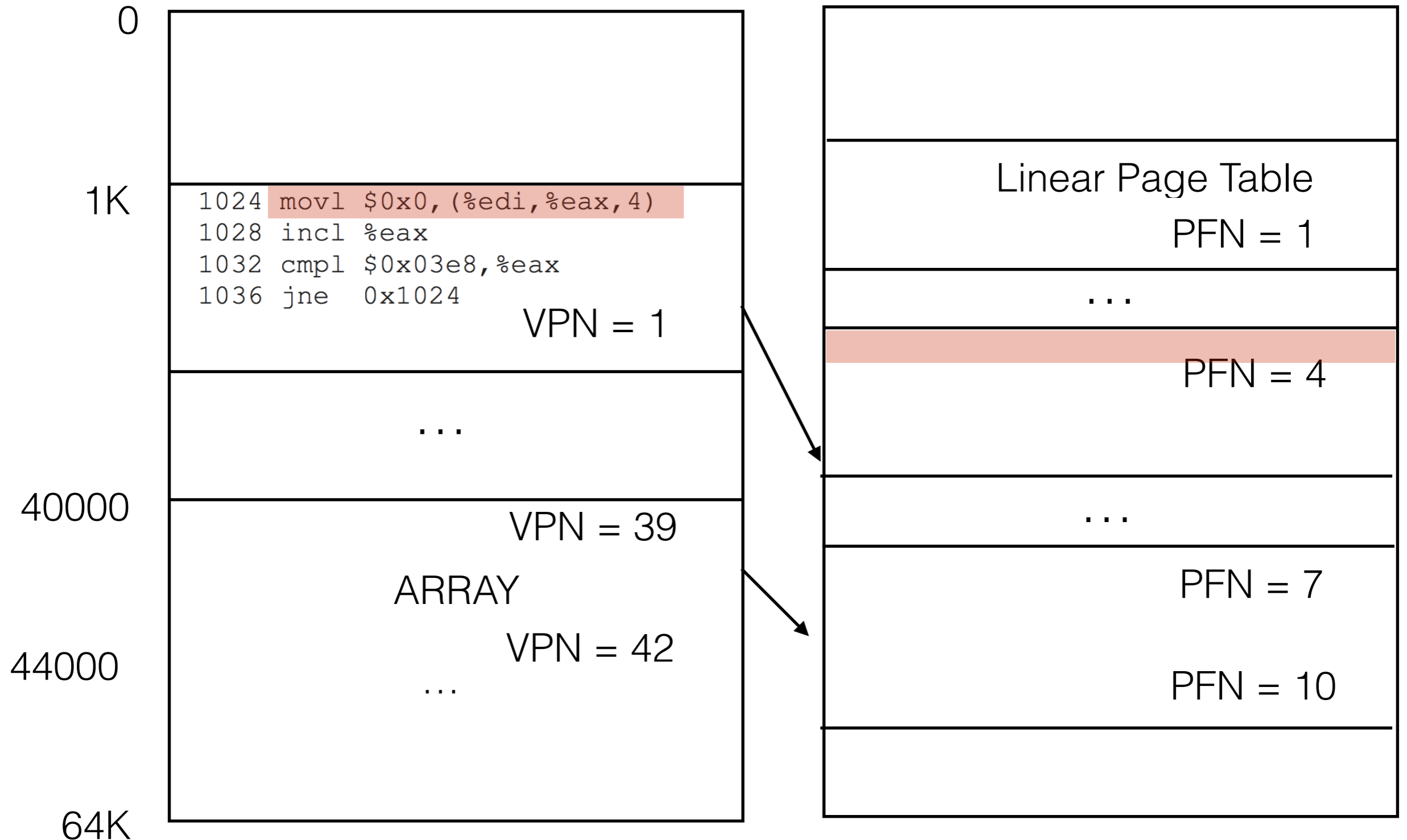
# Worked Out Example

## FIND PA FOR VA 1024 (VPN = 1)<sub>VA</sub> <sub>PA</sub>



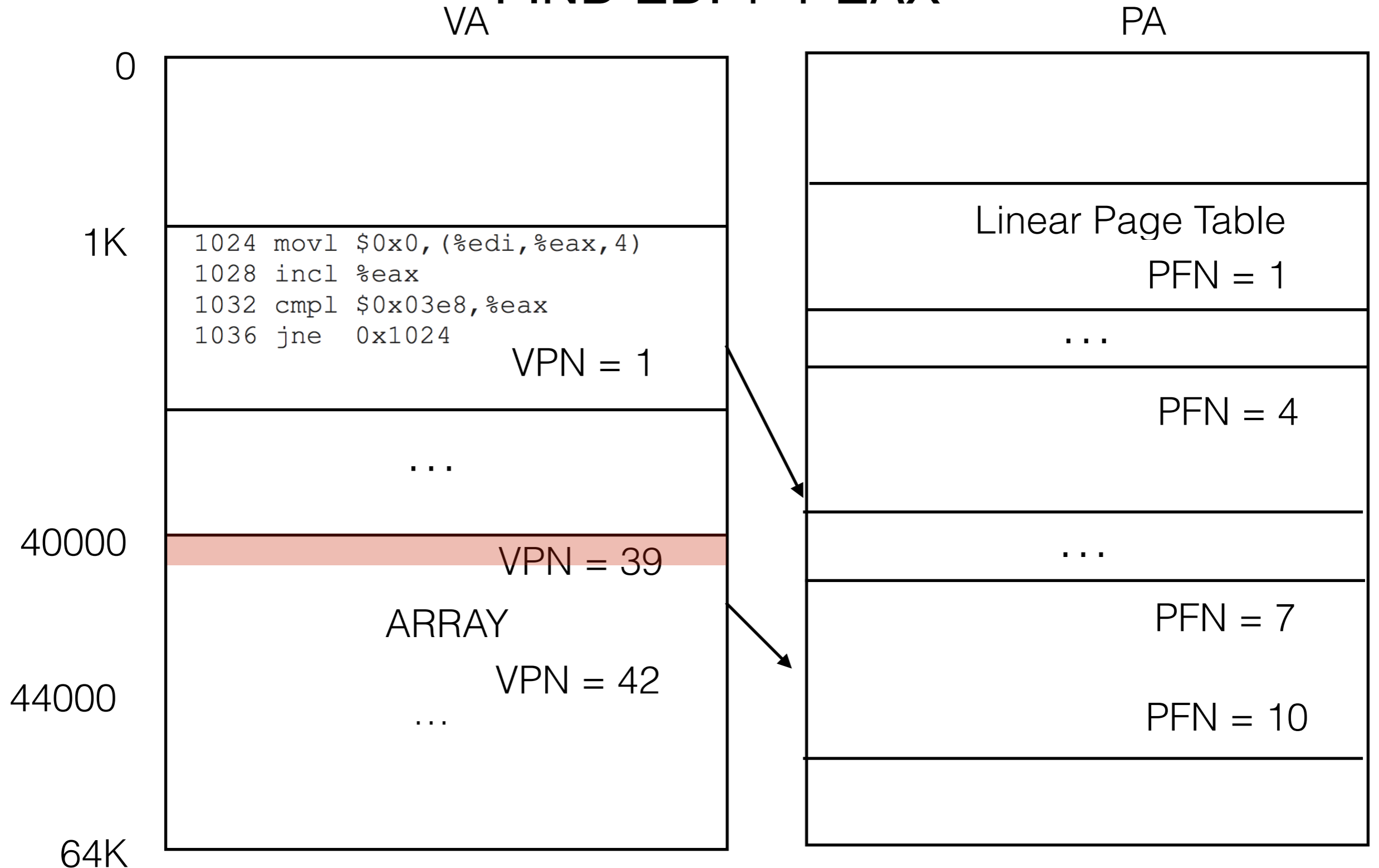
# Worked Out Example

## READ INSTRUCTION at PA(1024)<sub>PA</sub>



# Worked Out Example

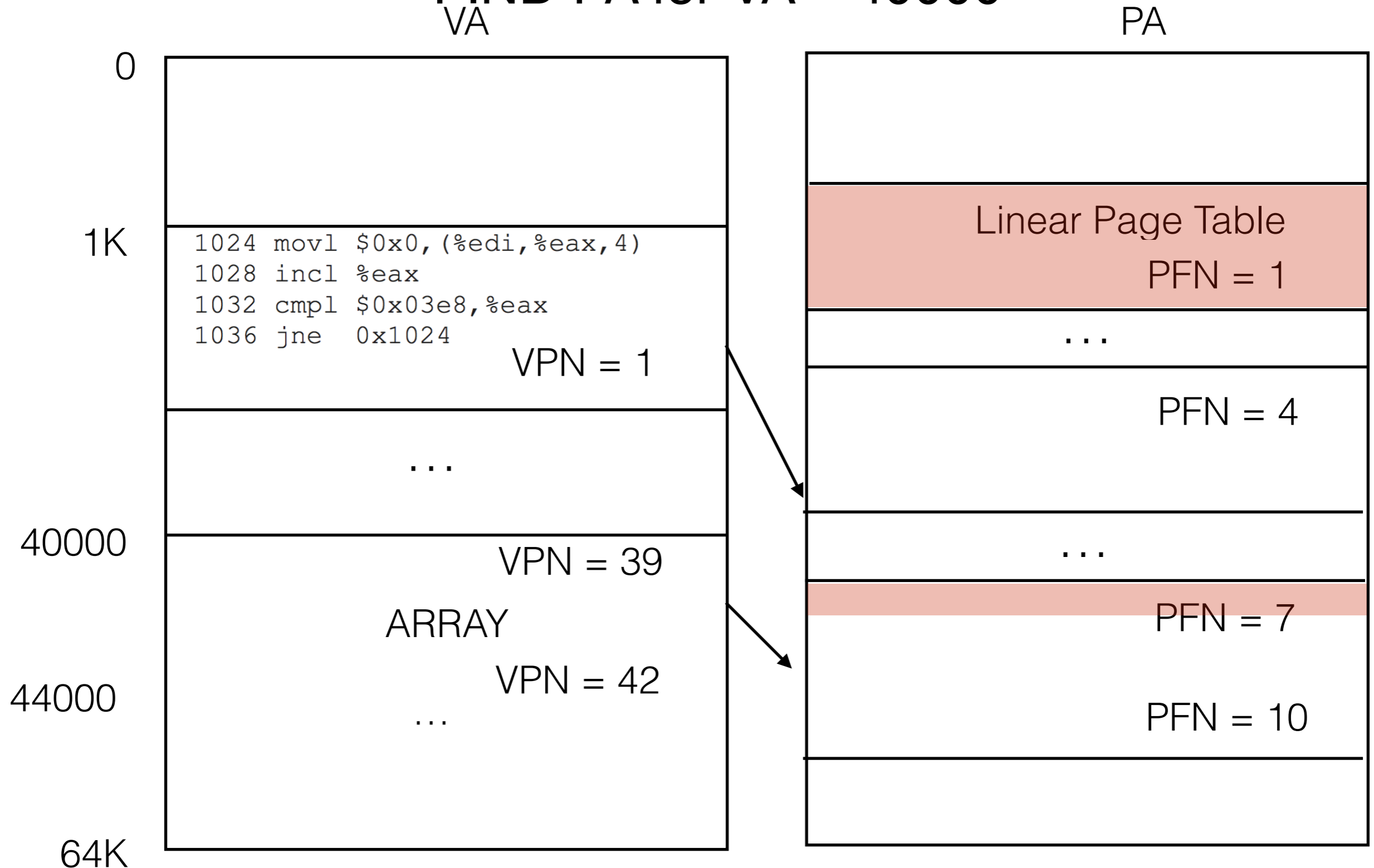
## VA **FIND EDI + 4\*EAX**





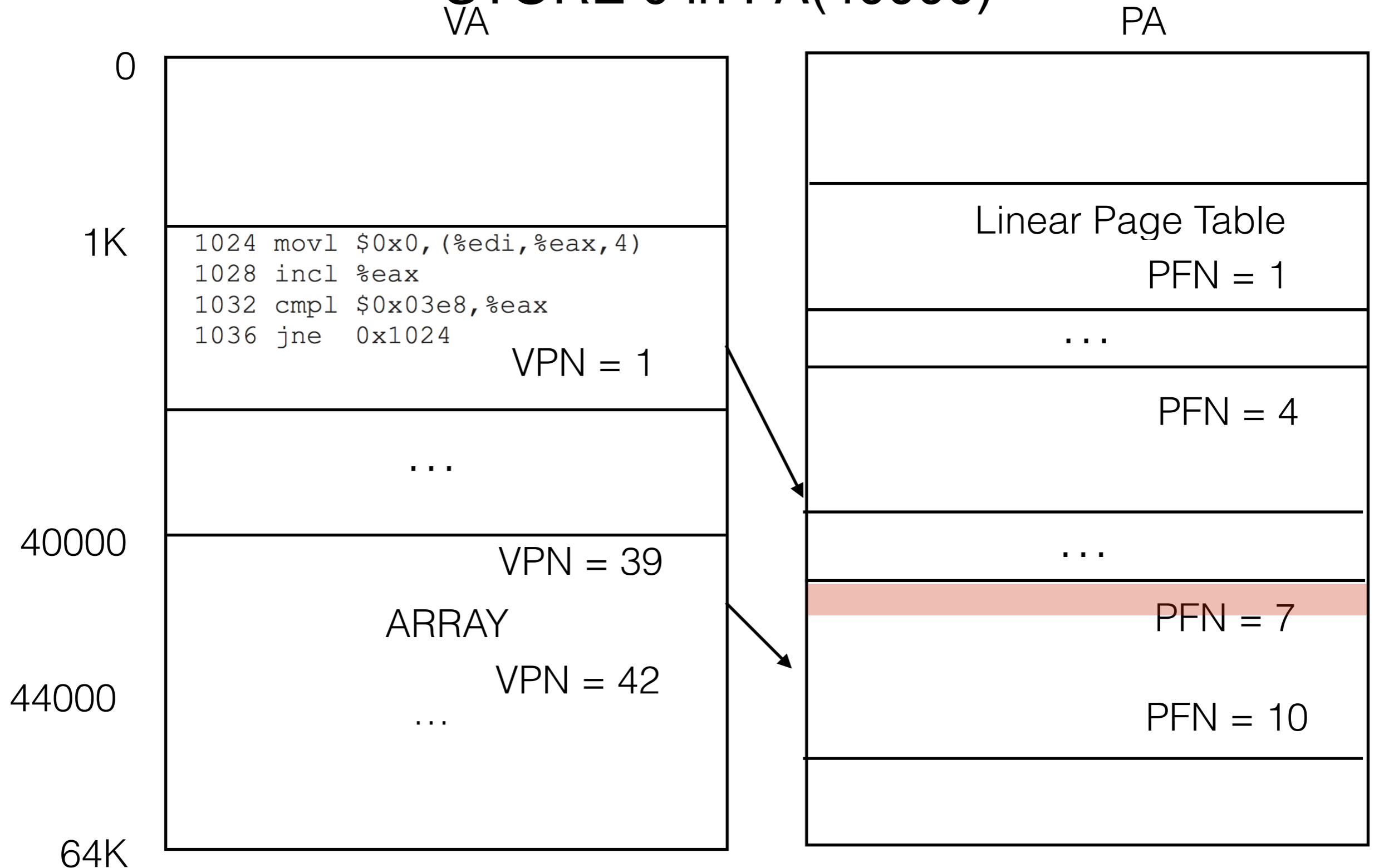
# Worked Out Example

## FIND PA for VA = 40000



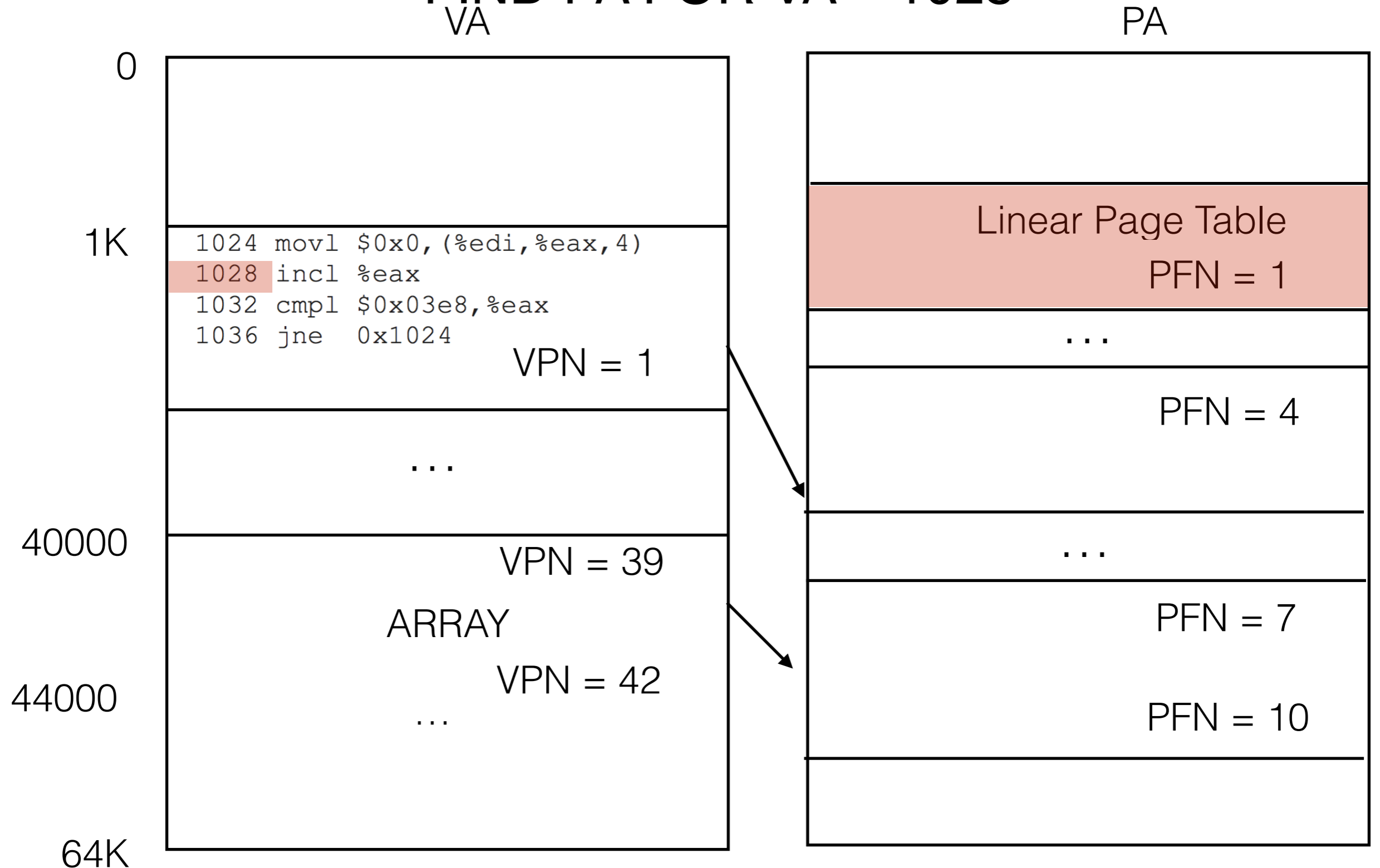
# Worked Out Example

## STORE 0 in PA(40000)



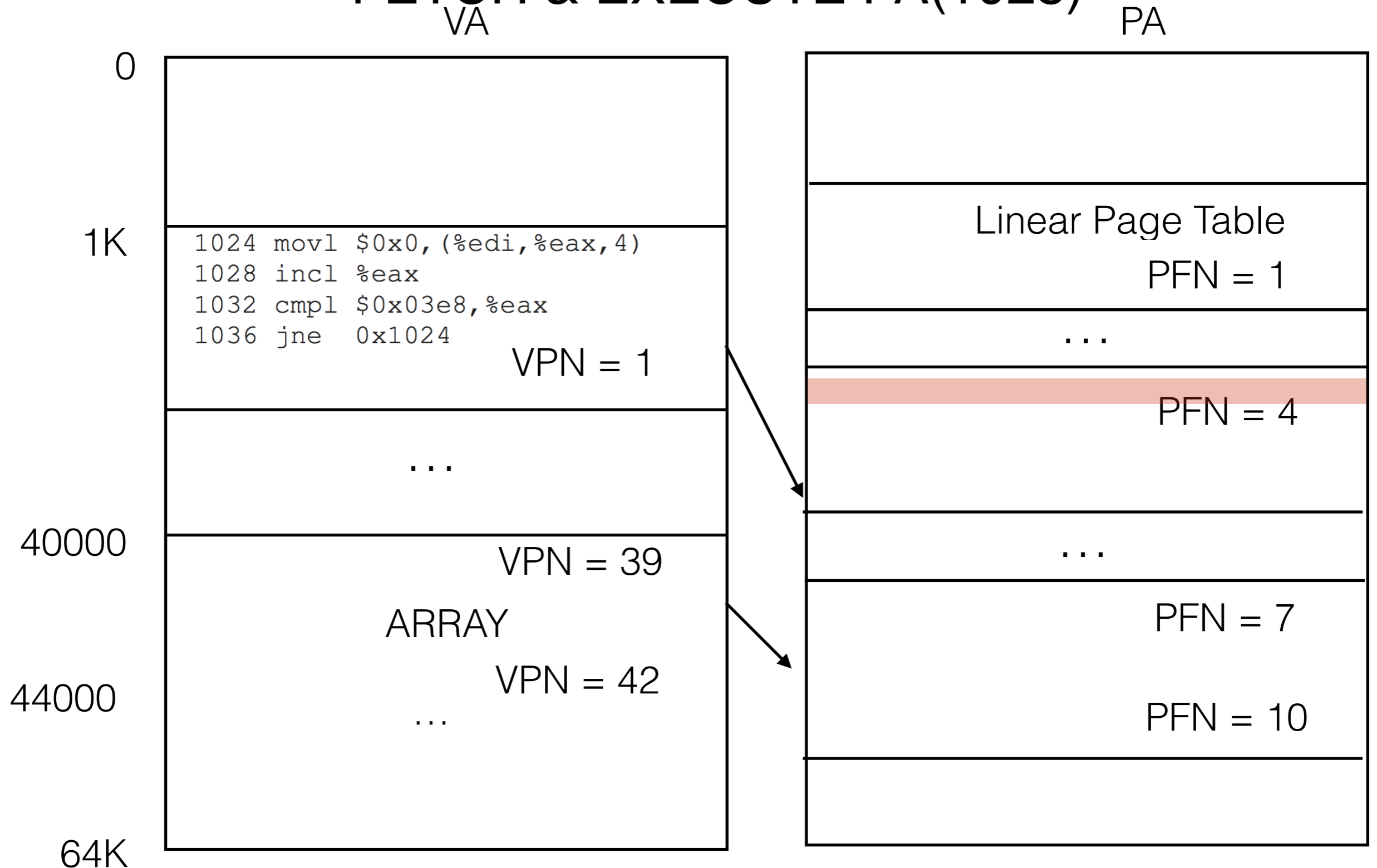
# Worked Out Example

## FIND PA FOR VA = 1028



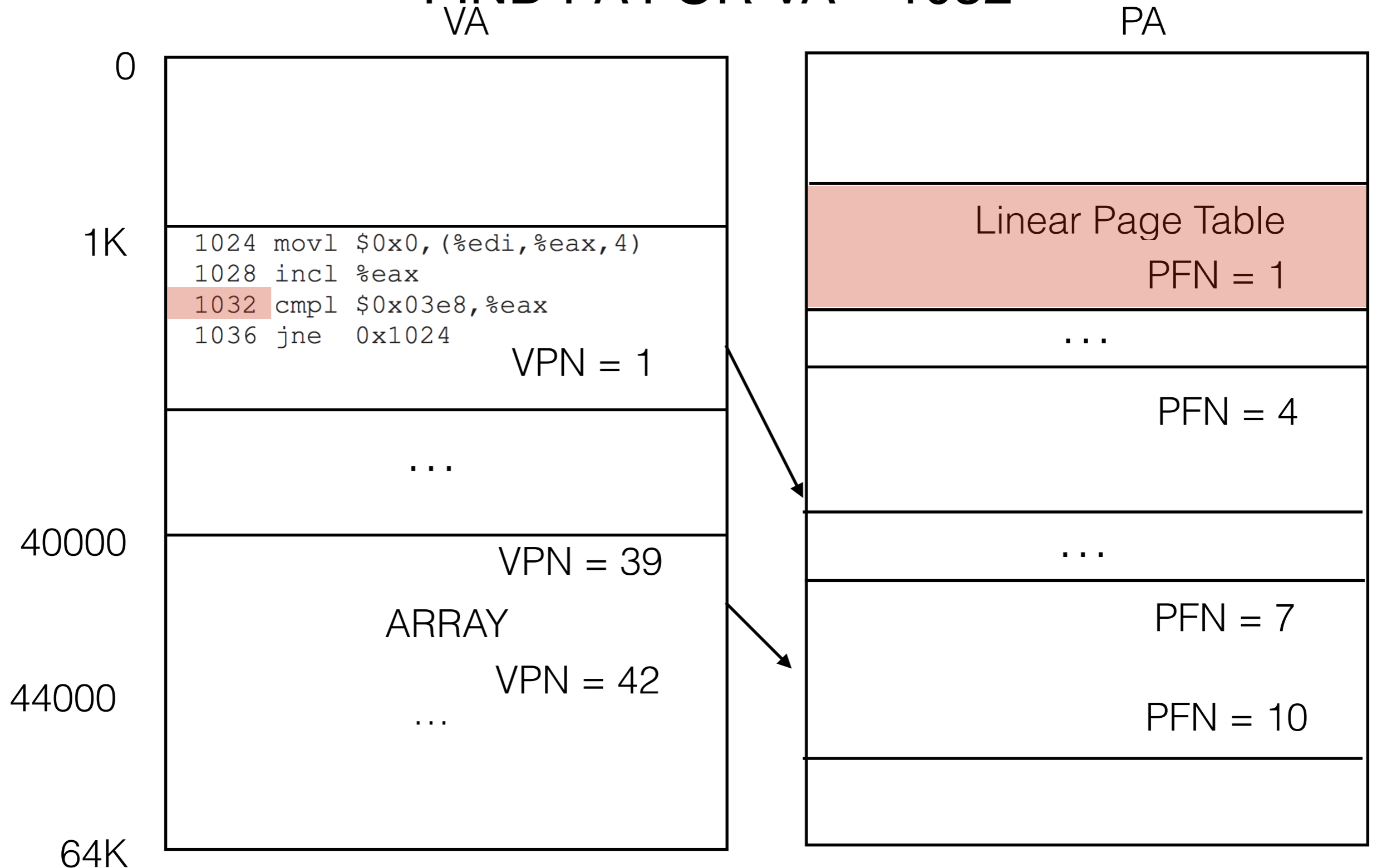
# Worked Out Example

## FETCH & EXECUTE PA(1028)



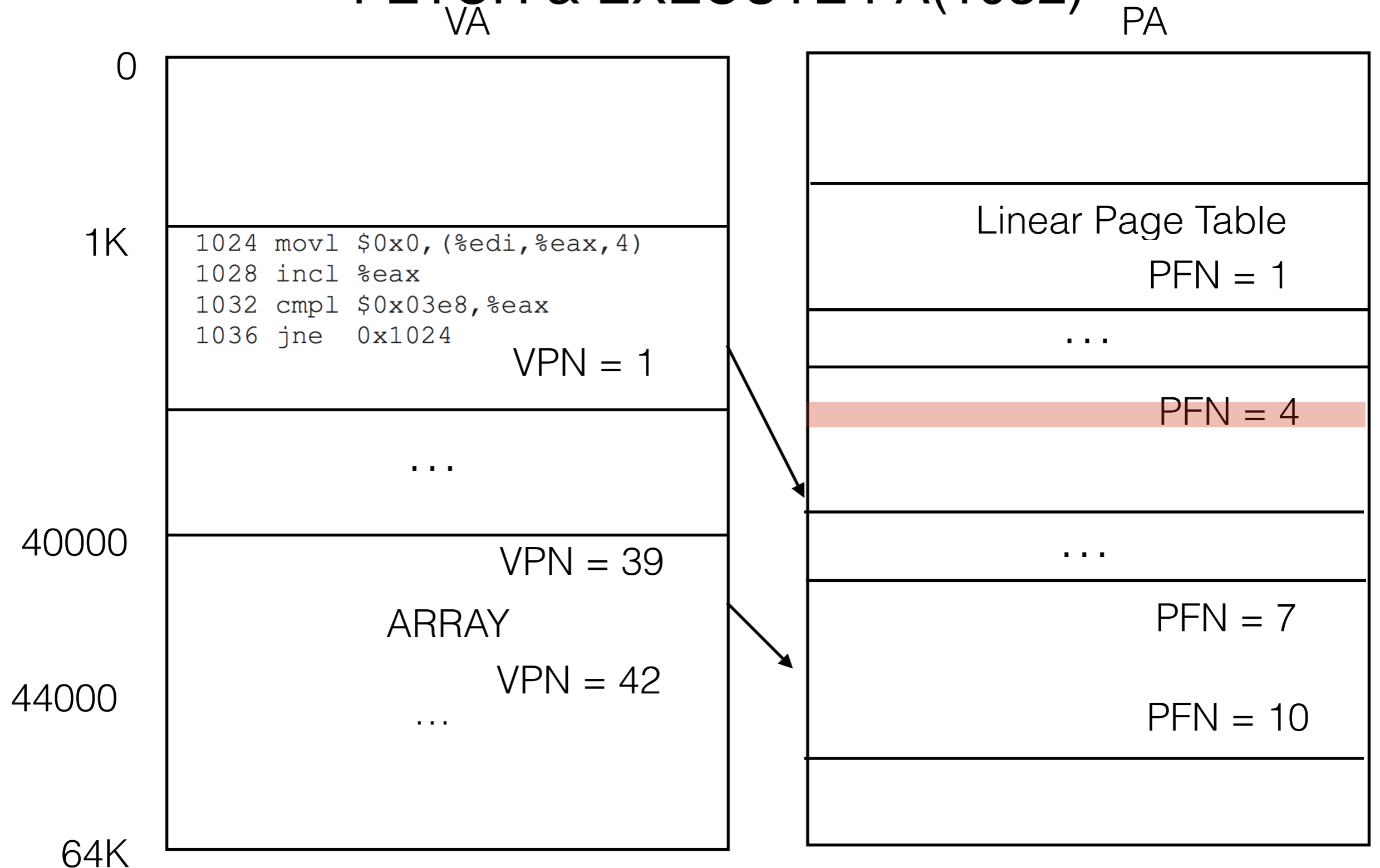
# Worked Out Example

## FIND PA FOR VA = 1032



# Worked Out Example

## FETCH & EXECUTE PA(1032)



# Example Summary

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1. Extract VPN (virt page num) from VA (virt addr)
2. Calculate addr of PTE (page table entry)
3. Read PTE from memory
4. Extract PFN (page frame num) **SLOW!**
5. Build PA (phys addr)
6. Read contents of PA from memory into register

# Caching Makes Sense!

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Factorial with and without memoization



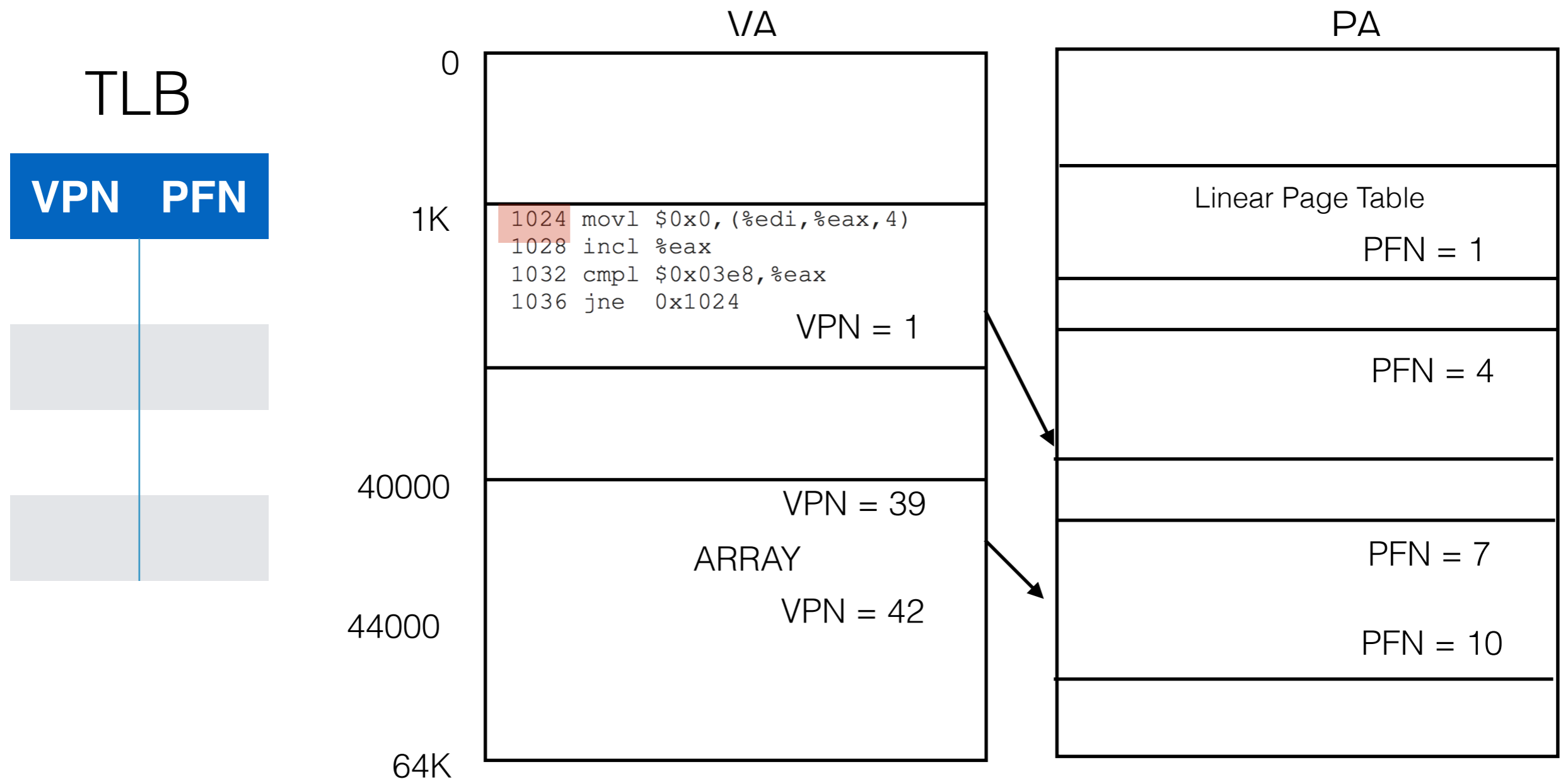
# Caching - Translation Lookaside Buffer (TLB)

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1. Get the VPN from VA
2. Check if TLB has VA
3. If found, it is a TLB Hit. Yay!
  - Extract the PFN from TLB ( $\text{PFN} = \text{TLB}[\text{VPN}]$ )
  - Generate PA from PFN (Add offset)
  - Access memory assuming protection checks work
4. If not found, it is a TLB Miss. :(
  - Access Page table to find the translation
  - Add translation to TLB ( $\text{TLB}[\text{VPN}] = \text{PFN}$ )
  - Goto Step 2

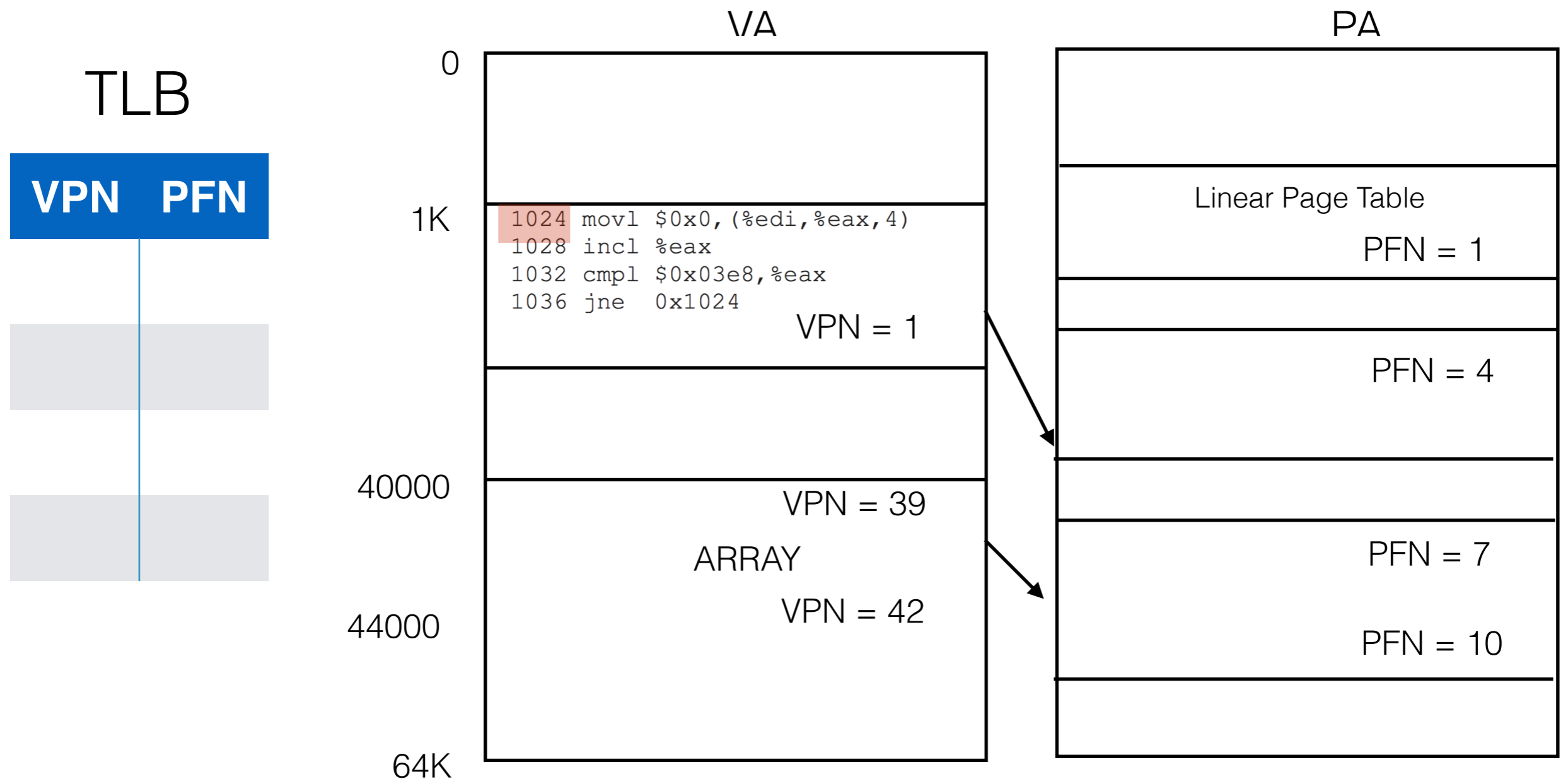
# Worked Out Example

## FETCH VA 1024



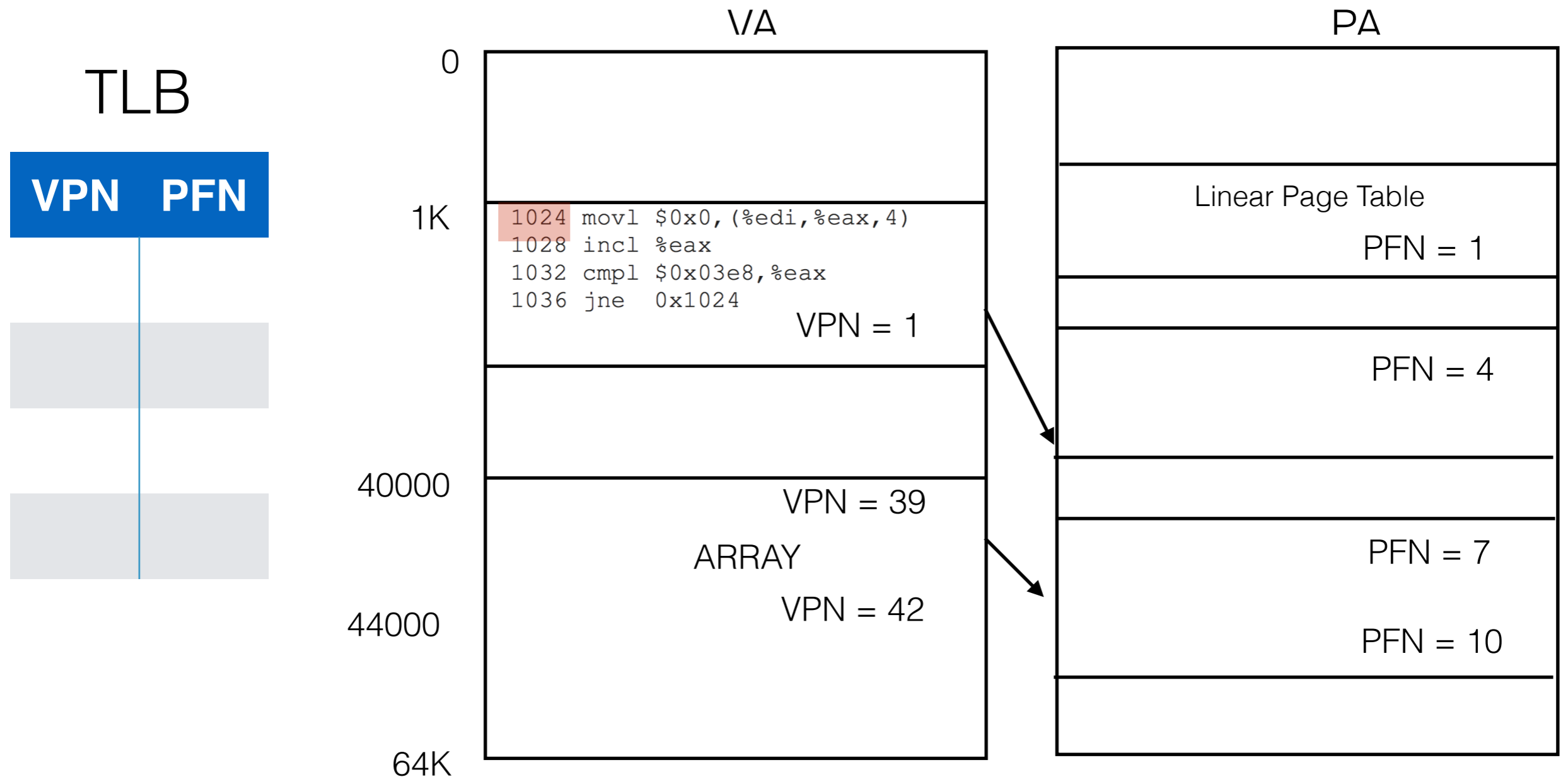
# Worked Out Example

Get VPN for VA 1024. VPN = 1



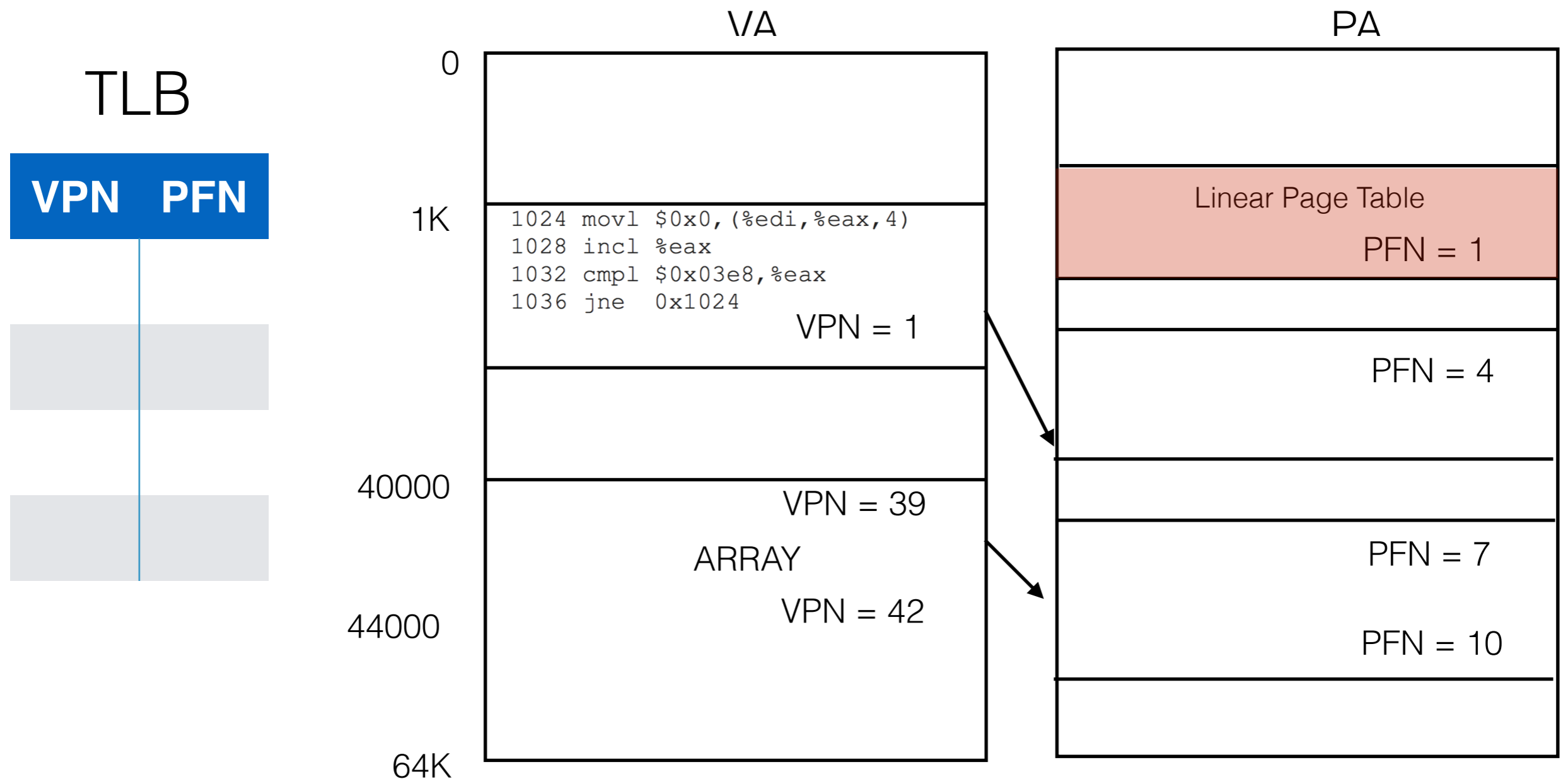
# Worked Out Example

**LOOK IN TLB for VPN = 1. Not found. TLB Miss!**



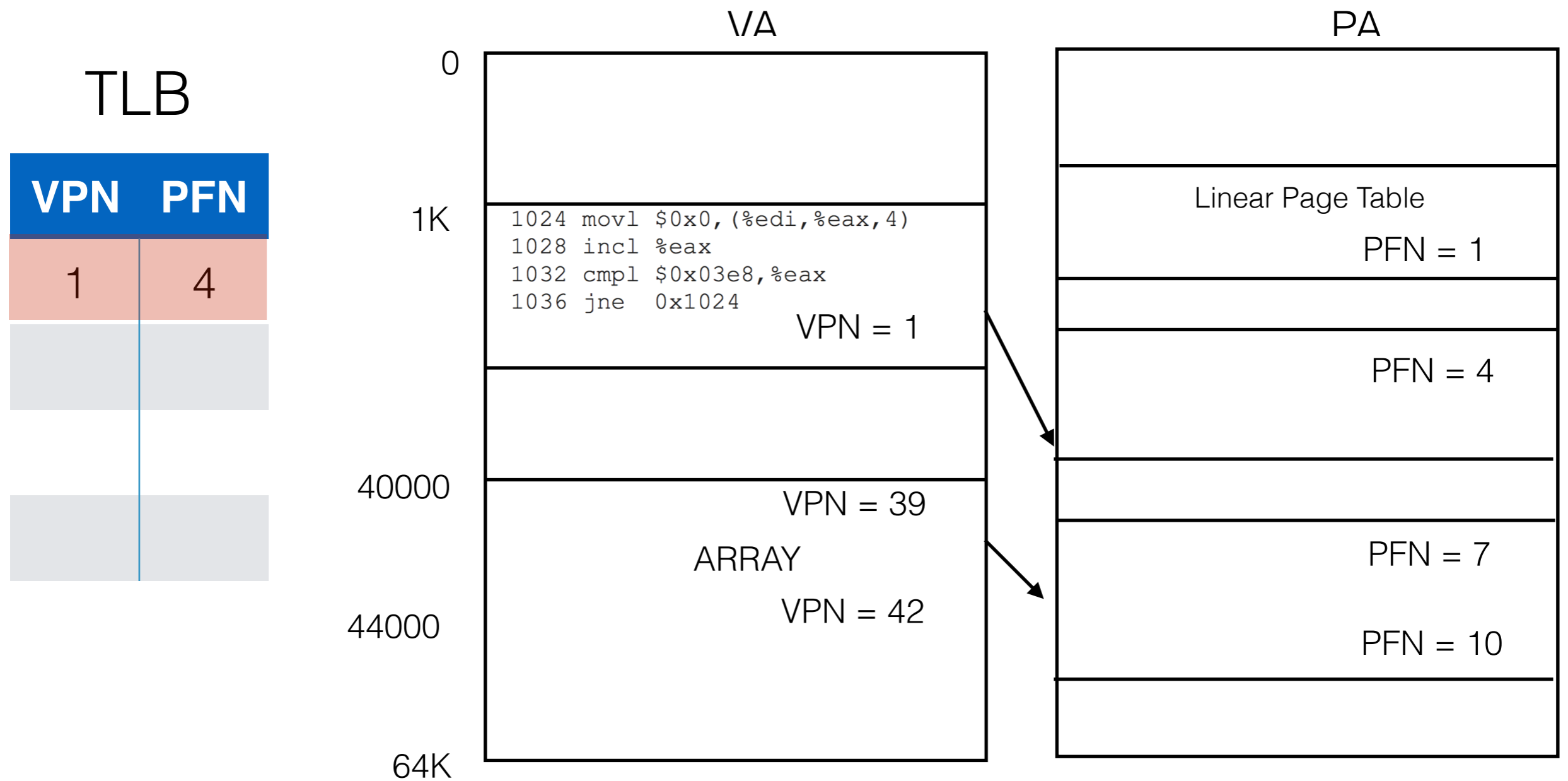
# Worked Out Example

Find PFN for VPN = 1 by accessing Page Table



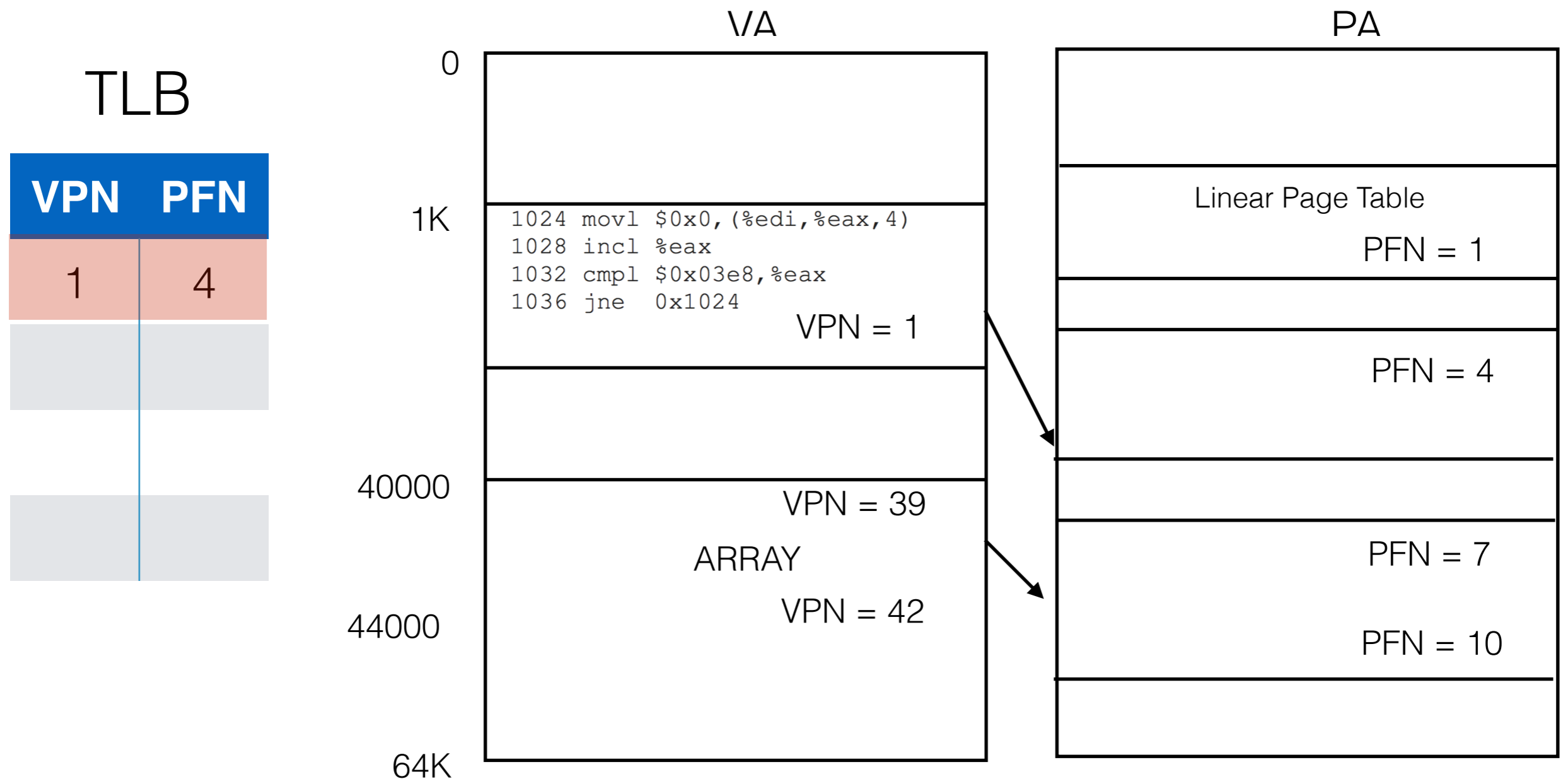
# Worked Out Example

## Add entry to TLB



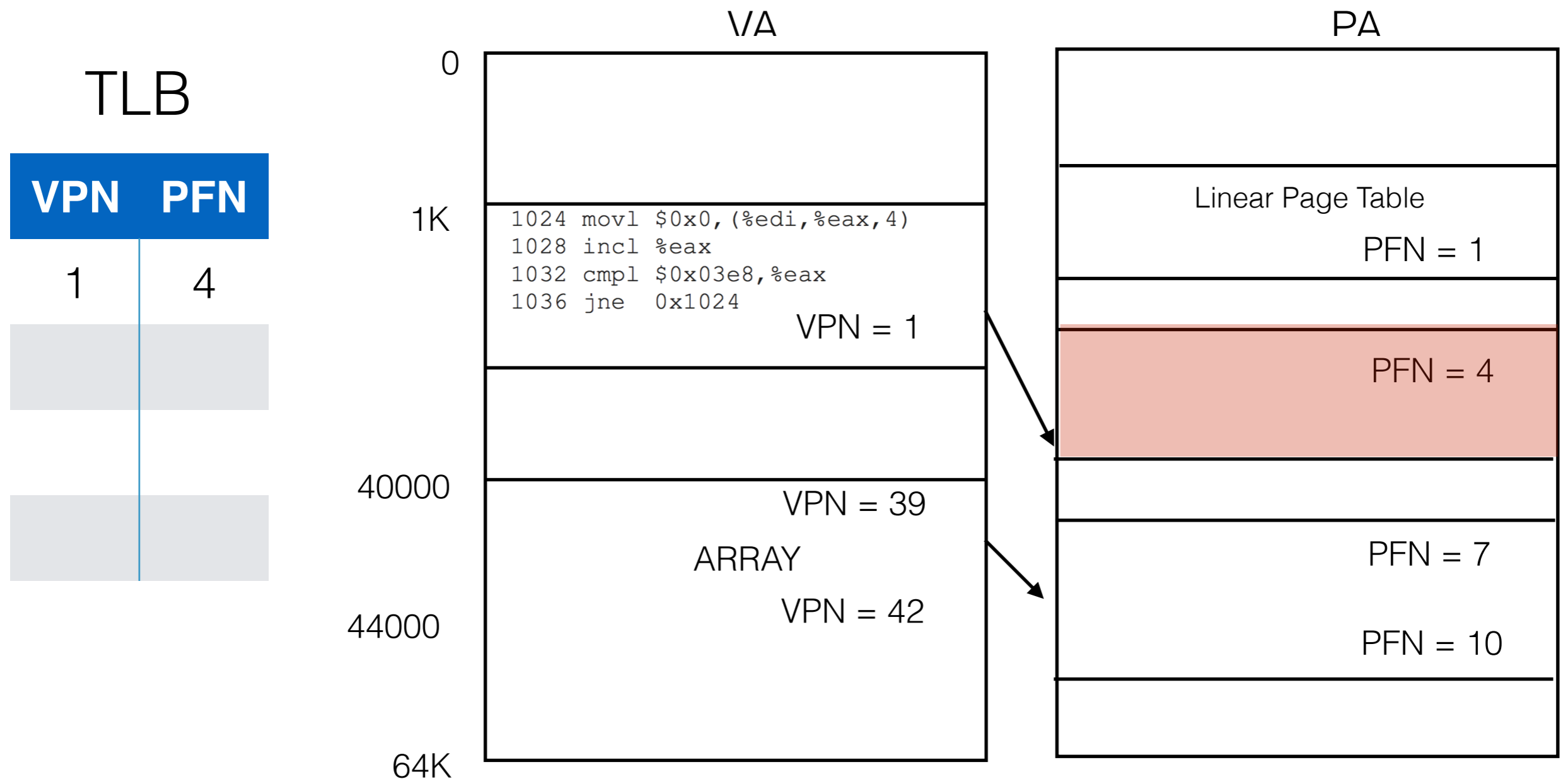
# Worked Out Example

## Search for translation of VPN = 1 on TLB



# Worked Out Example

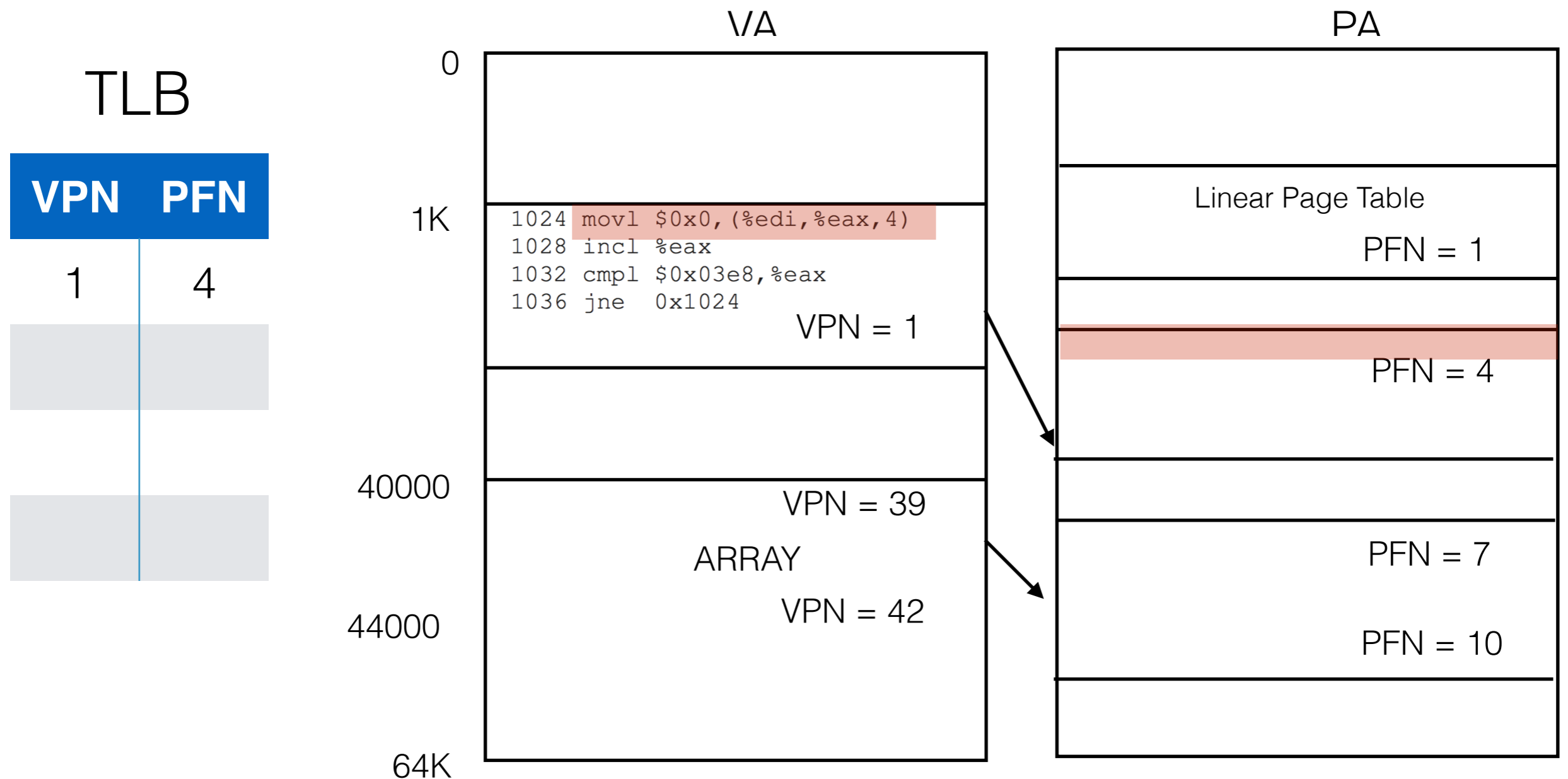
Goto PFN 4 and create PA by adding offset





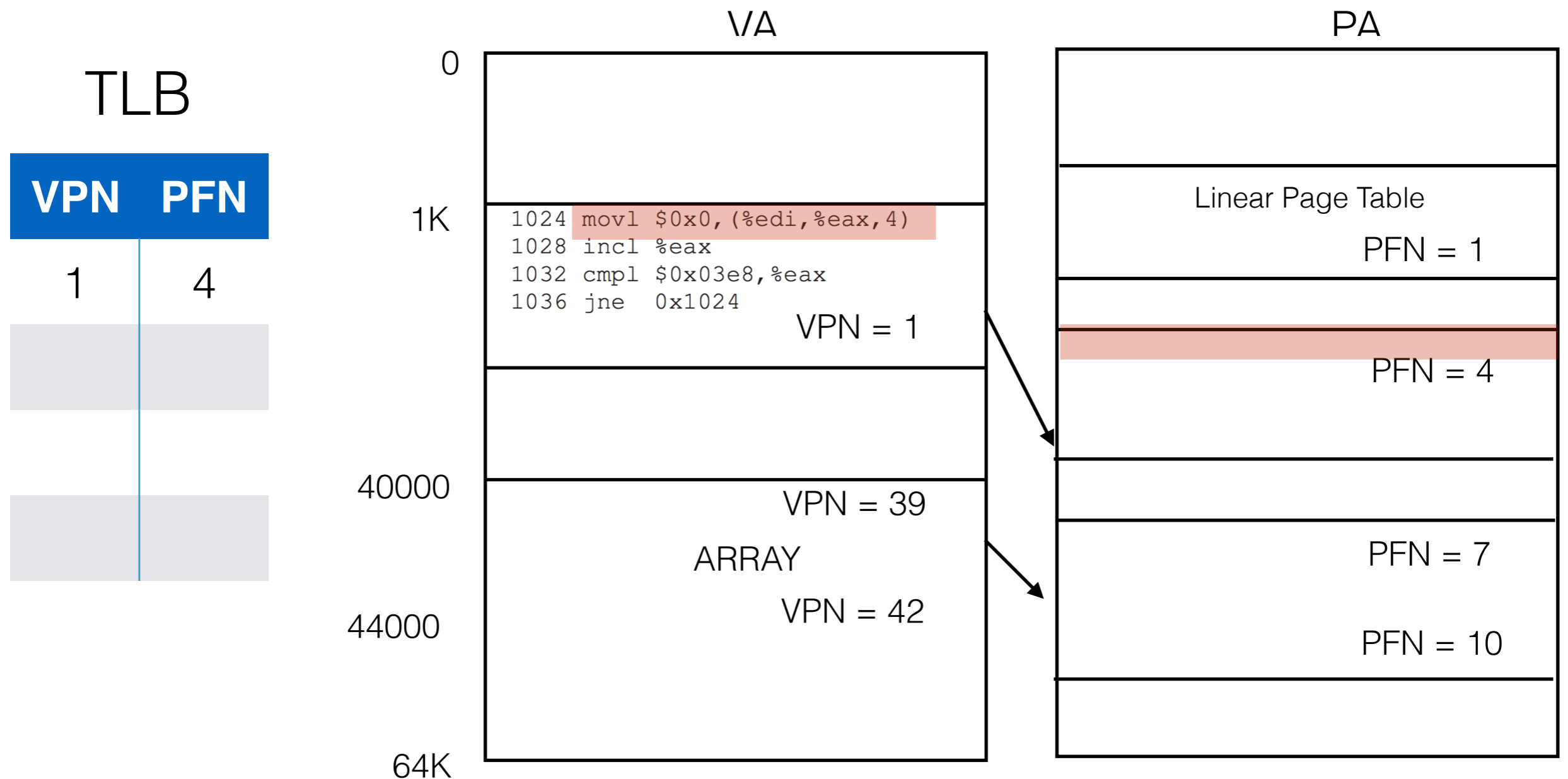
# Worked Out Example

## READ INSTRUCTION at PA(1024)



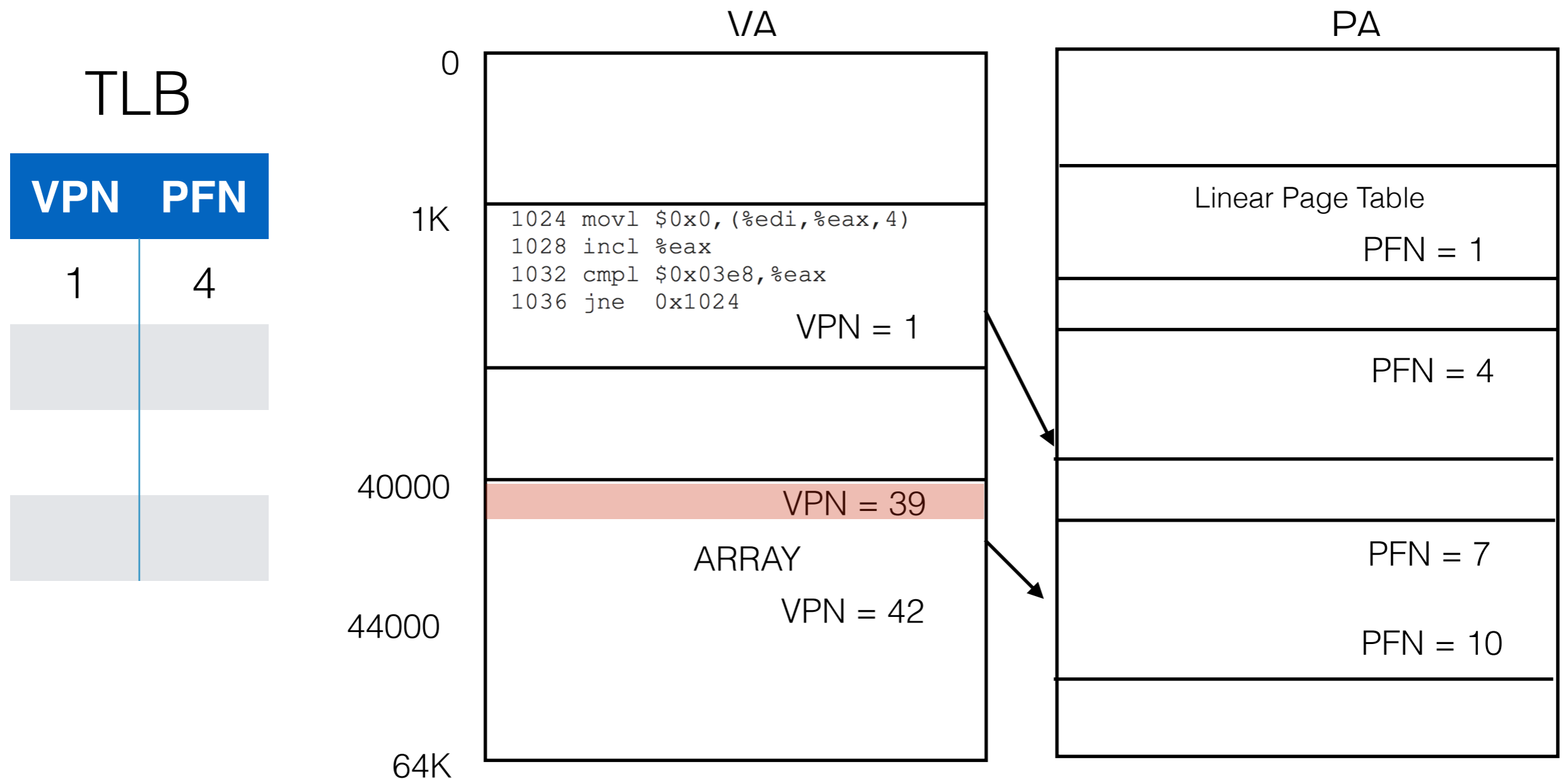
# Worked Out Example

## READ INSTRUCTION at PA(1024)



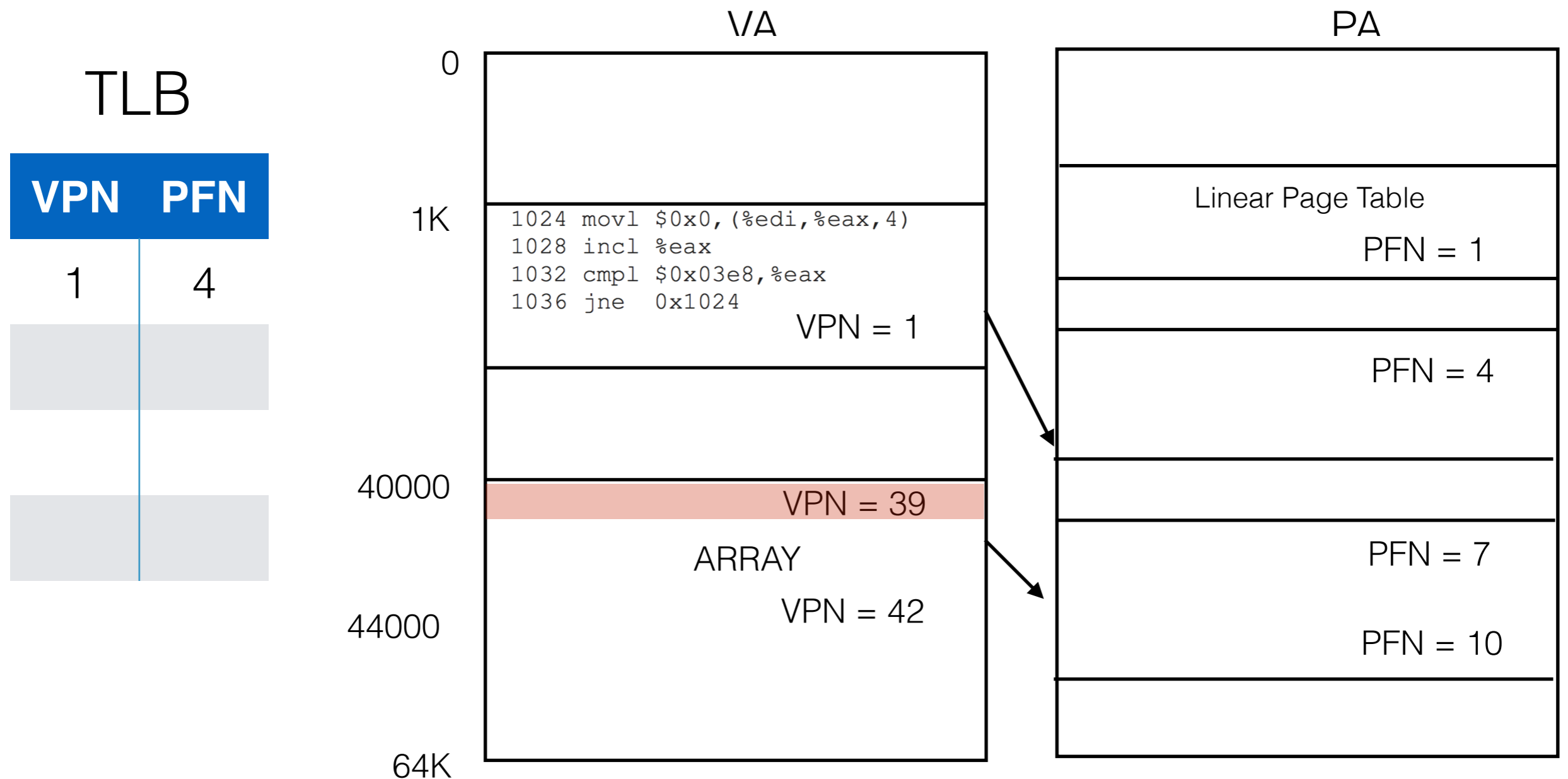
# Worked Out Example

Find  $EDI + 4 * EAX \rightarrow VA = 40000$



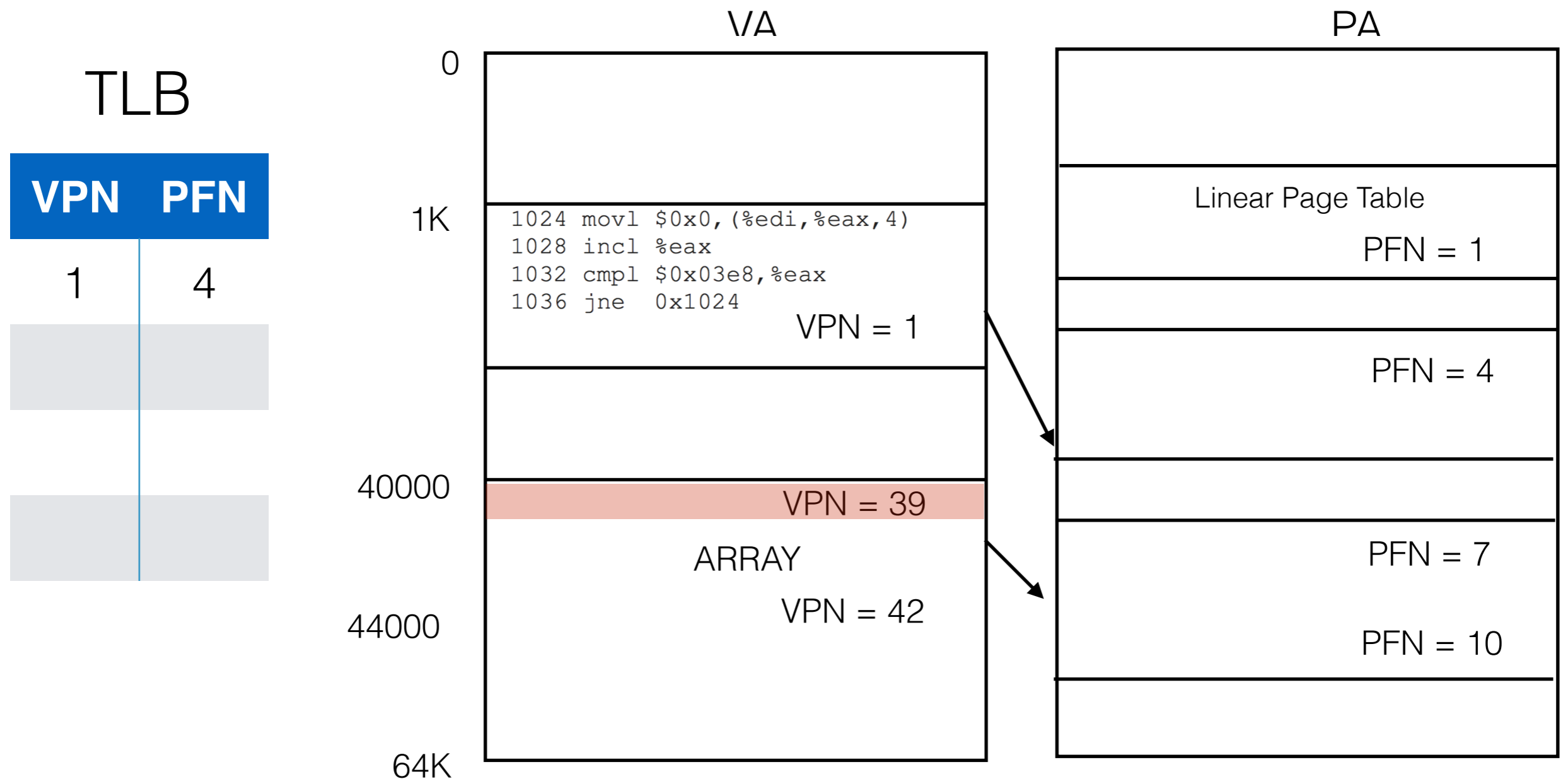
# Worked Out Example

Find VPN for VA 40000. VPN = 39



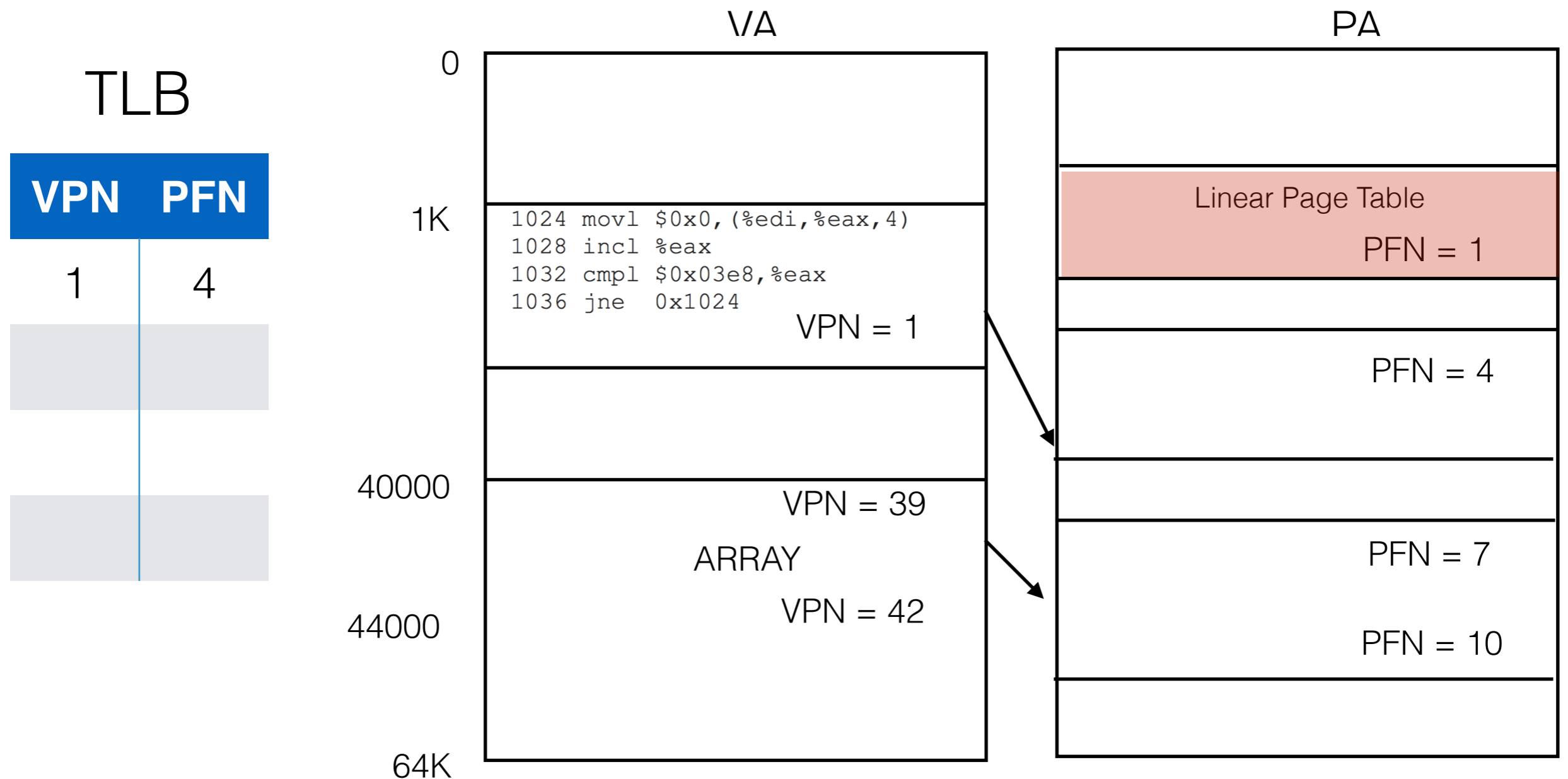
# Worked Out Example

Check TLB for VPN = 39. Miss!



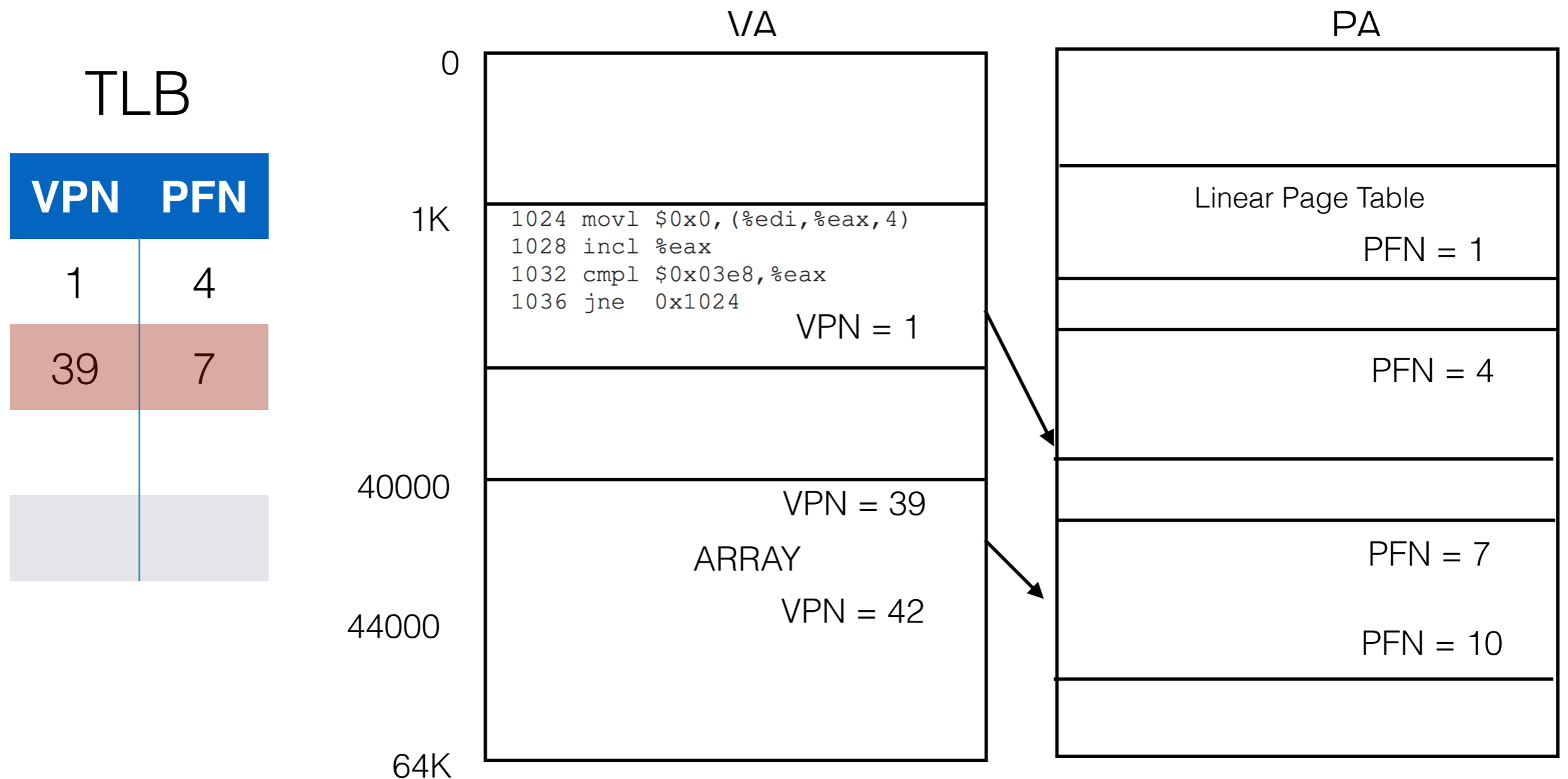
# Worked Out Example

## Get PFN for VPN = 39 from Page Table



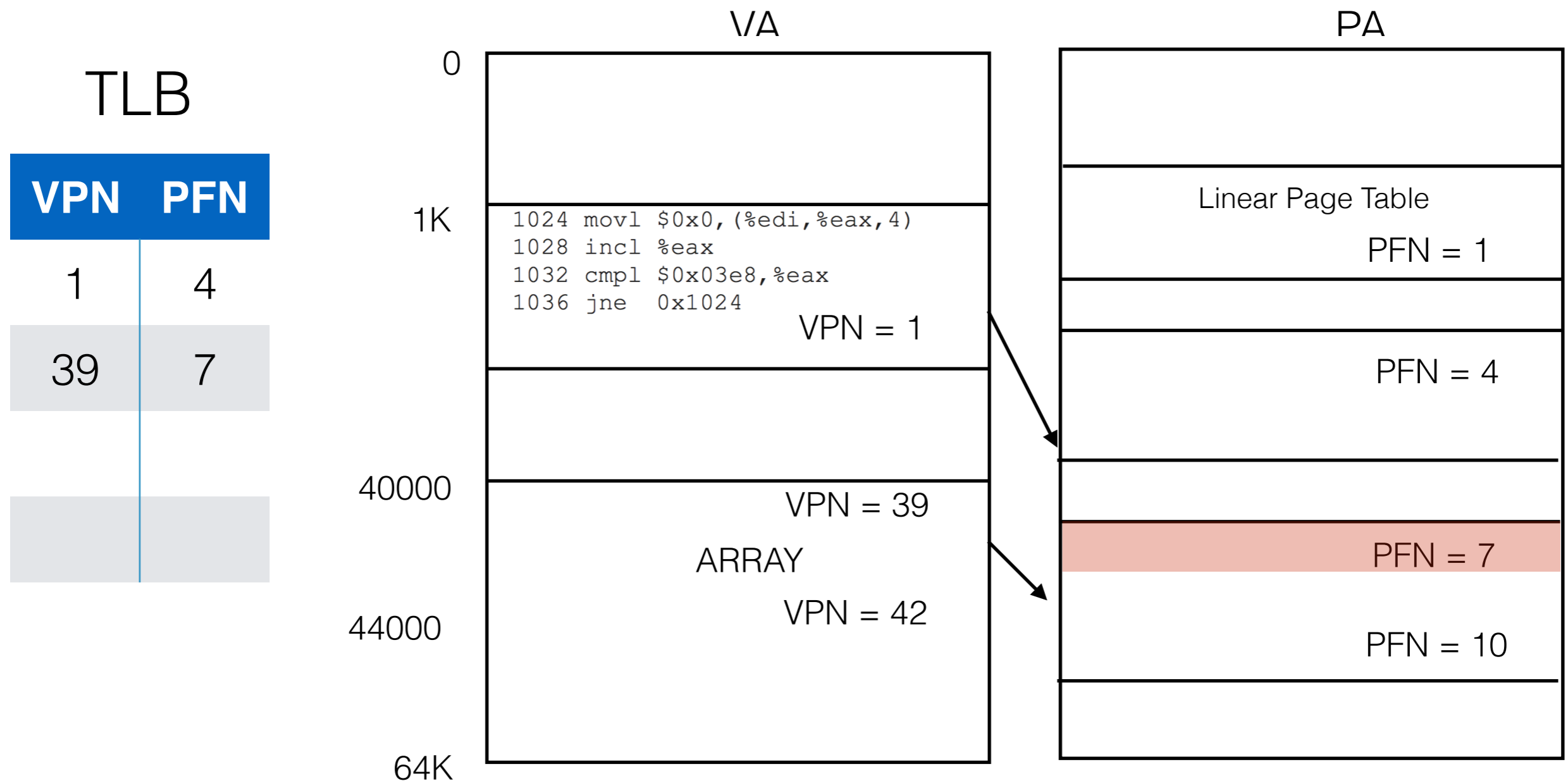
# Worked Out Example

## Store translation in TLB



# Worked Out Example

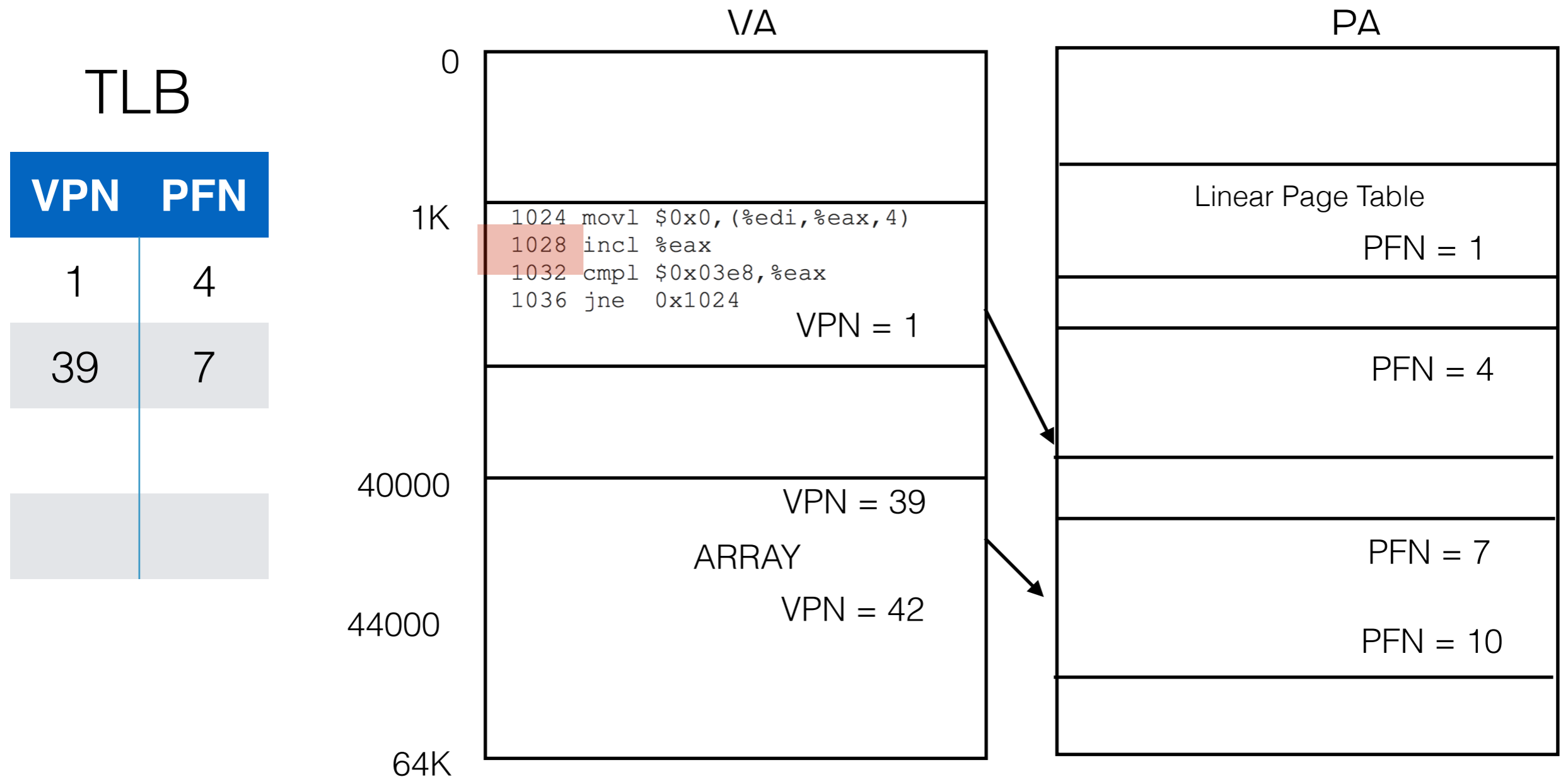
Find PFN of VPN = 39 from TLB. Add offset to get PA.





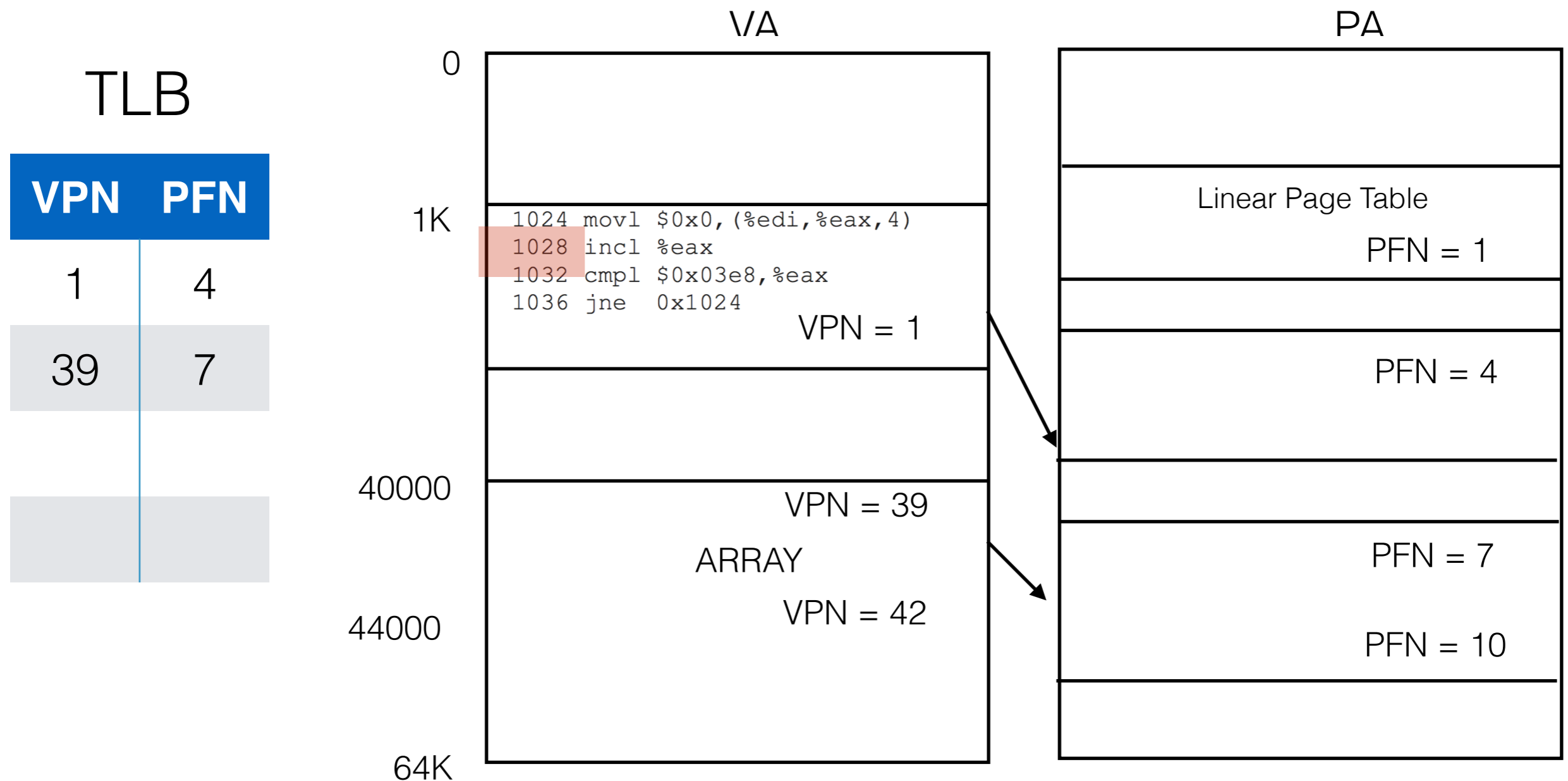
# Worked Out Example

## FIND PA FOR VA = 1028



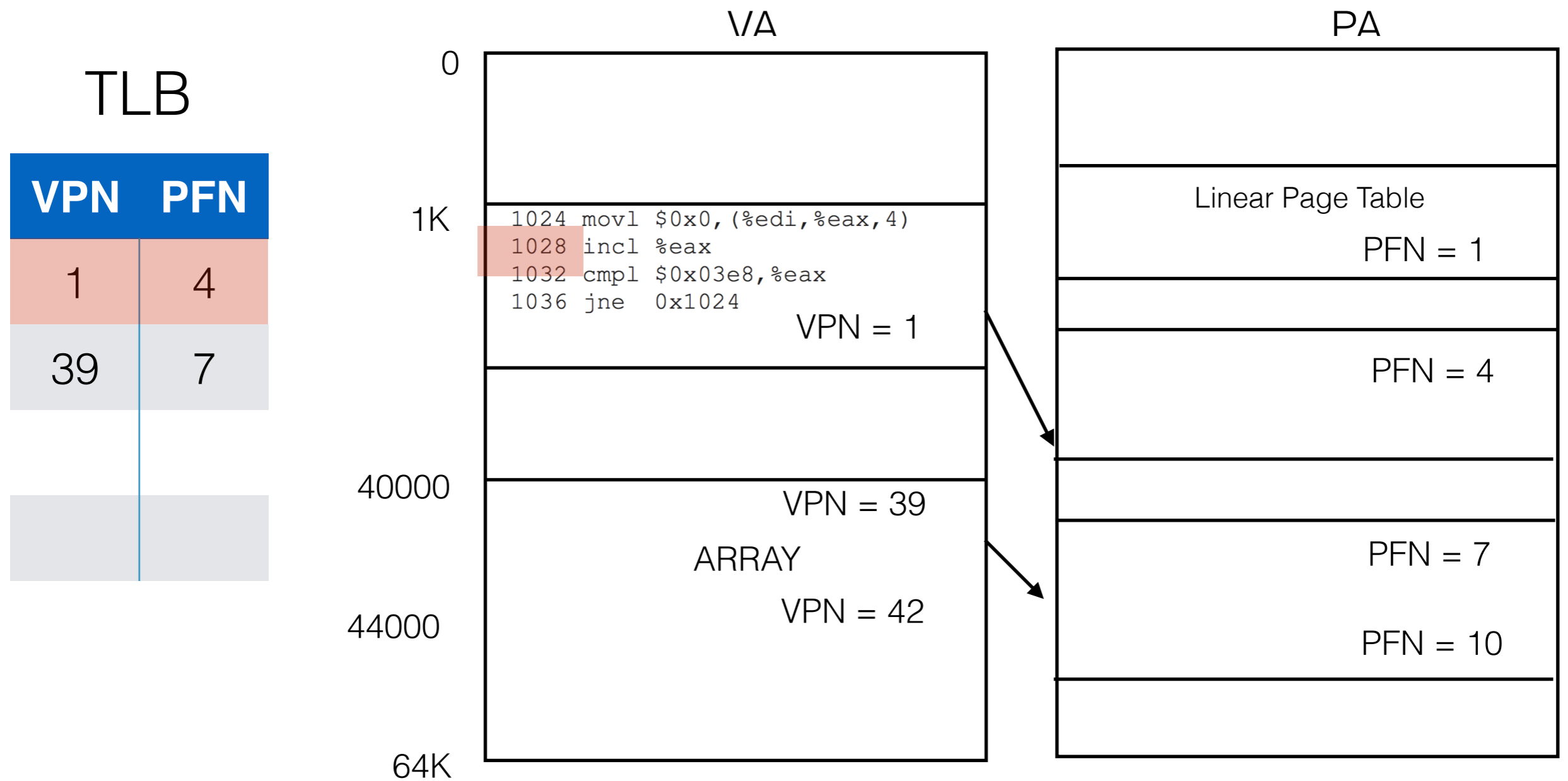
# Worked Out Example

VPN = 1. Find Translation in TLB for VPN = 1. Found!



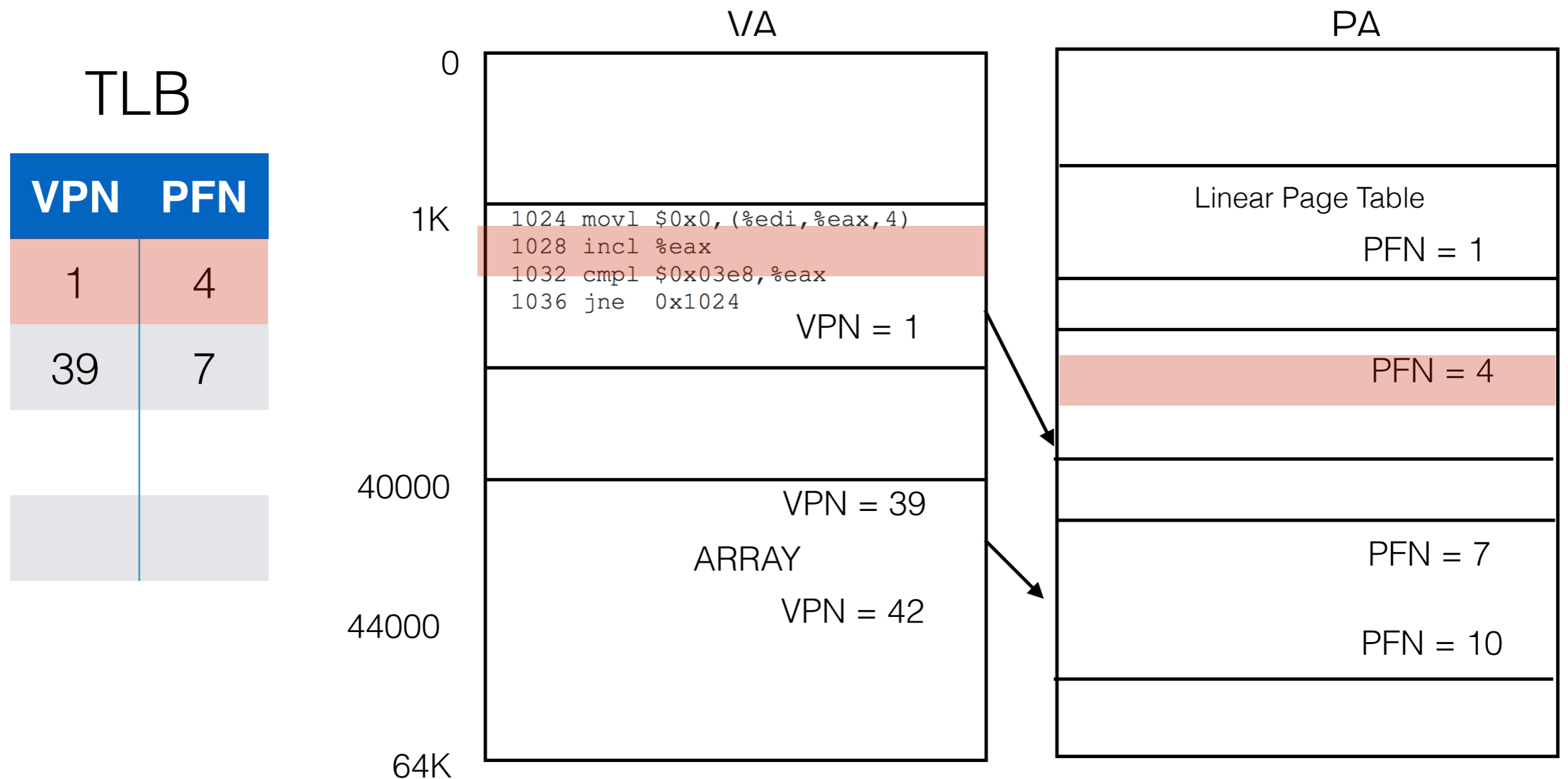
# Worked Out Example

$$\text{PFN} = \text{TLB}[1] = 4$$



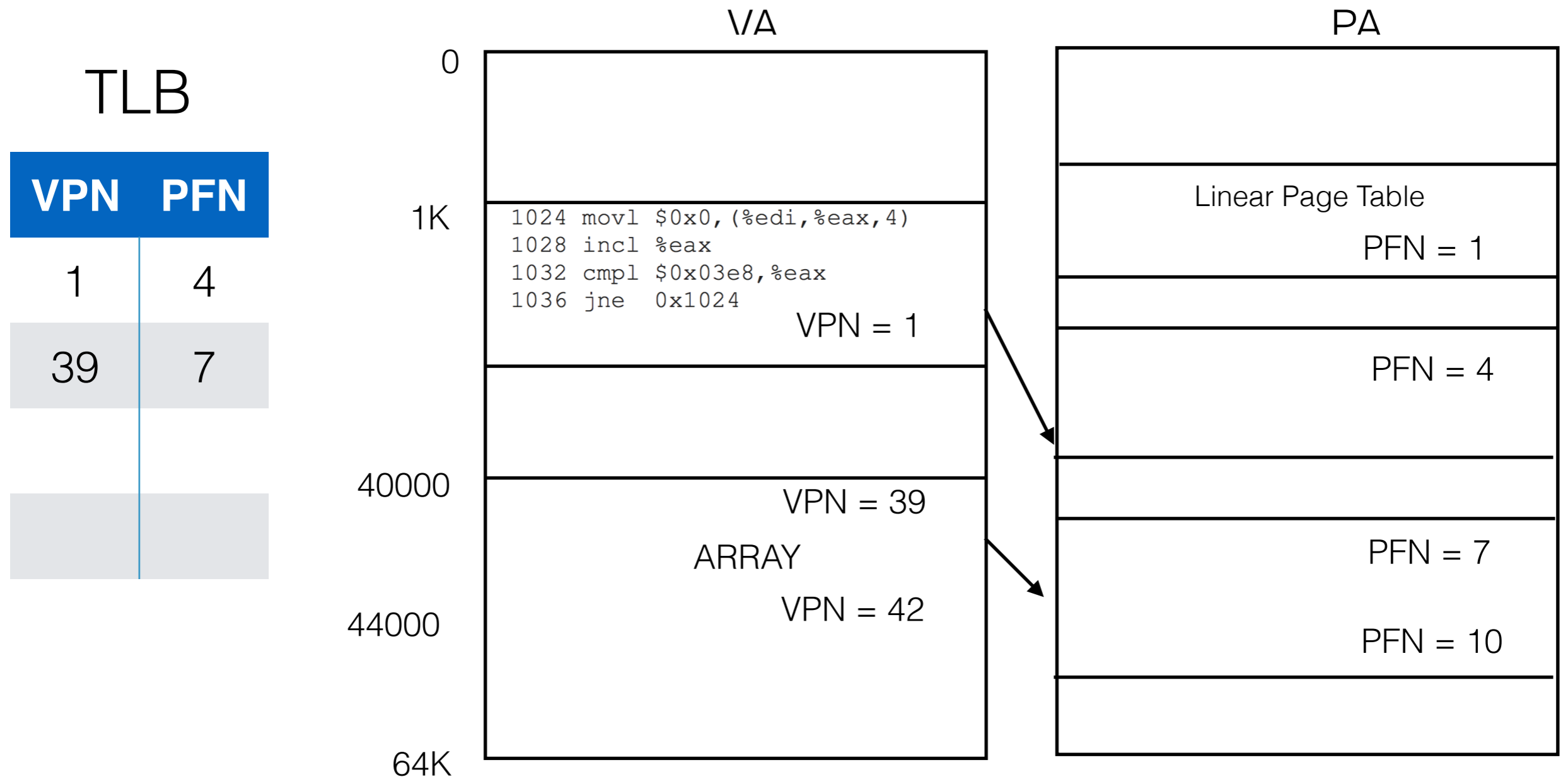
# Worked Out Example

Get PA by adding offset to PFN = 4 and execute



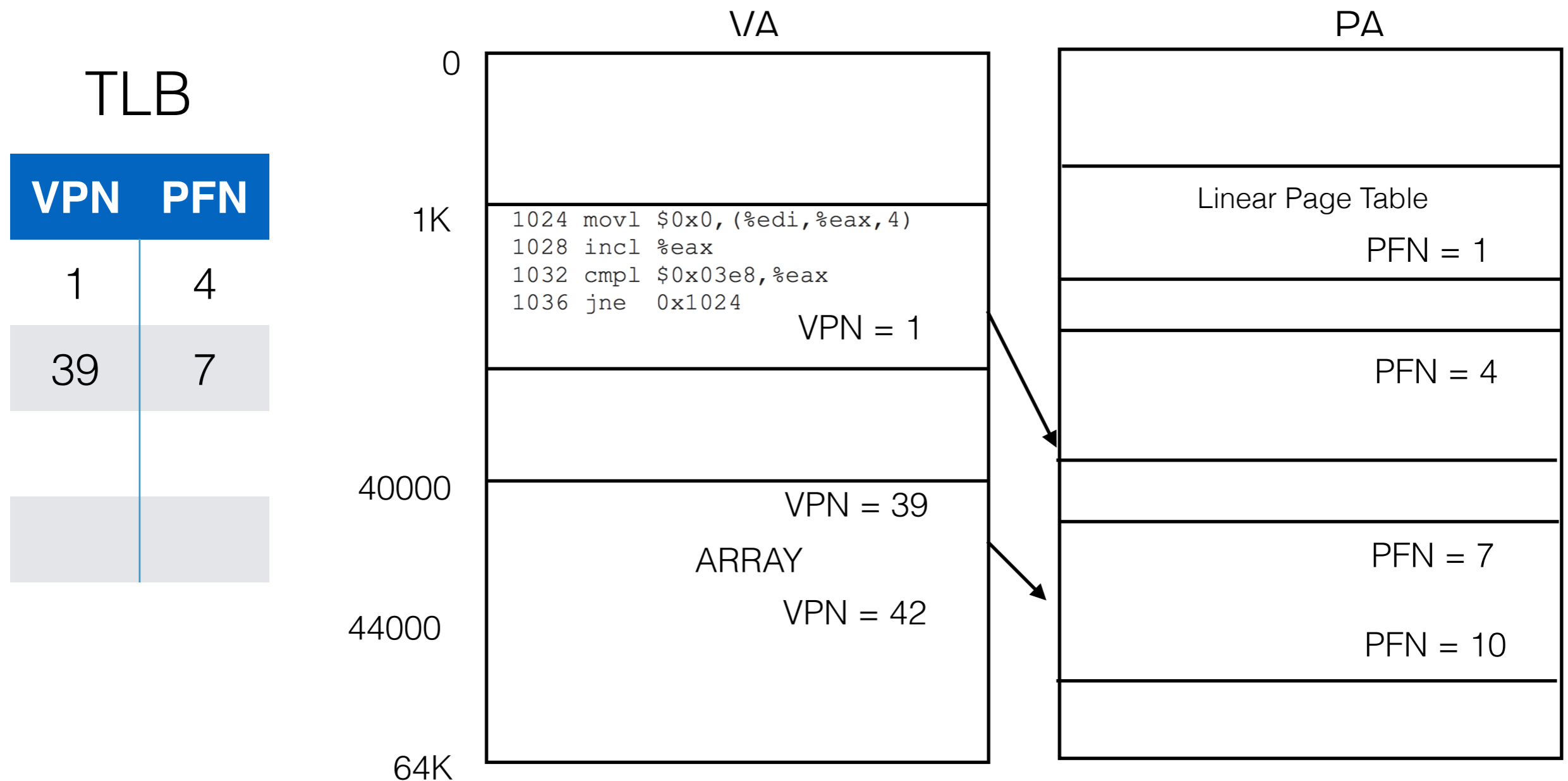
# Worked Out Example

1032, 1036, 1024, 1028,.....



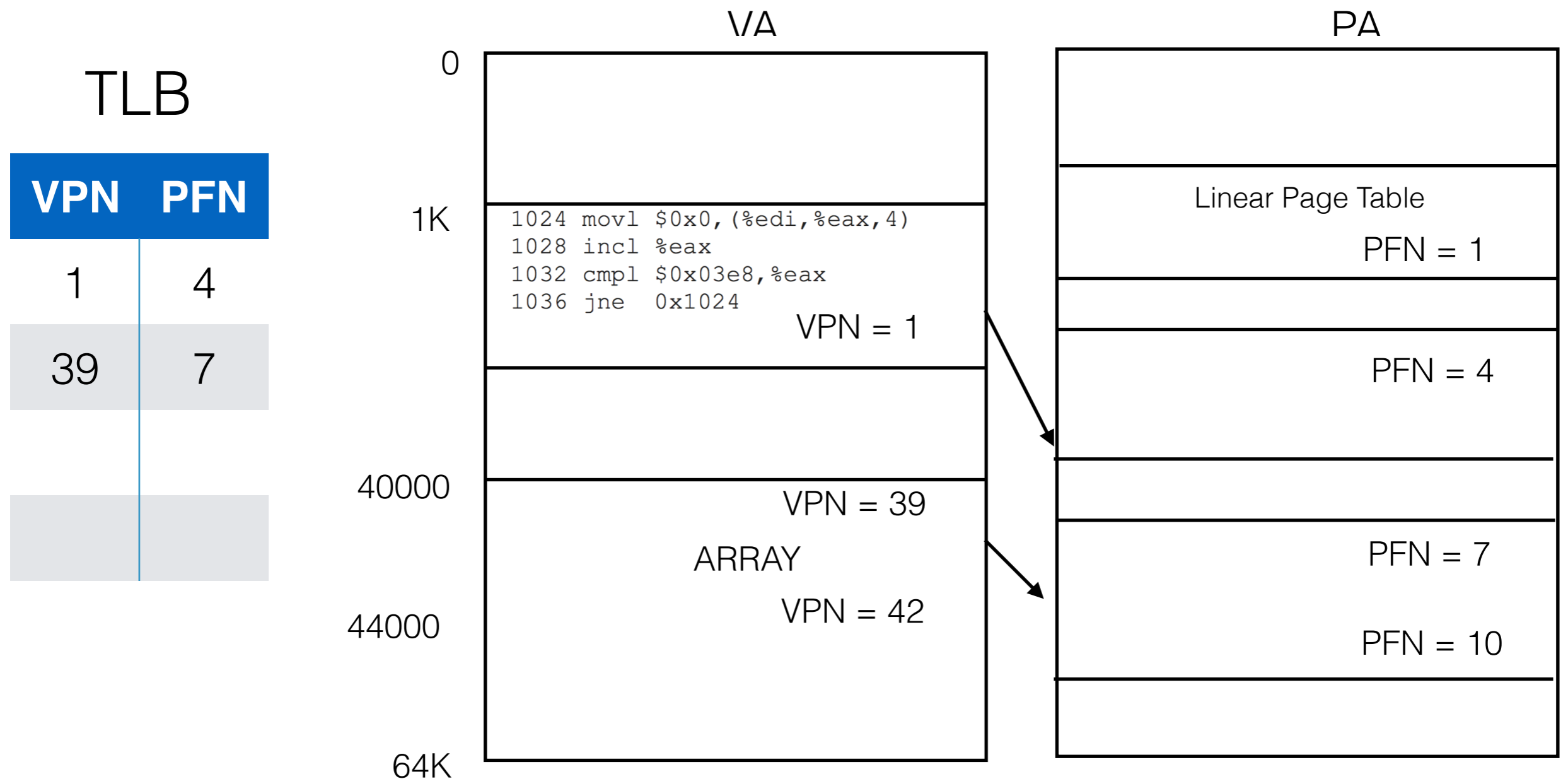
# Worked Out Example

$$EDI + 4 * EAX = 40000 + 4 * (1024/4) = 40000 + 1K \rightarrow VPN = 40$$



# Worked Out Example

## TLB miss for VPN = 40...



# Spatial and Temporal Locality

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1. Hit rate =  $\text{TLB Hit} / (\text{TLB Hit} + \text{TLB Miss})$
2. Spatial locality -> TLB has good hit rate
  1. Arrays elements are spatially close ( $\text{EDX} + 4 * \text{EAX}$ )
  2. Instructions are spatially close (1024, ...)
3. Temporal locality -> TLB has a good hit rate
  1. Loop. Re-using same instructions which exist in TLB



# Memory Cycle Rate Example

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1. Hit = 1 clock cycle
2. Miss = 30 clock cycles
3. Miss rate = 1%
4. Cycle rate =  $.99 * 1 + .01 * (30 + 1) = 1.3$  cycles

# Context Switch

---

TLB

P1 running

VPN	PFN
1	4
39	7

---

P2 running

# Context Switch

---

## TLB

P1 running

VPN	PFN
1	4
39	7
...	...
1	30

---

P2 running

What will VPN 1 be mapped to?